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(19) **United States**(12) **Patent Application Publication**
Eroz et al.(10) **Pub. No.: US 2004/0054960 A1**(43) **Pub. Date: Mar. 18, 2004**(54) **METHOD AND SYSTEM FOR PROVIDING
LOW DENSITY PARITY CHECK (LDPC)
ENCODING**(76) Inventors: **Mustafa Eroz**, Germantown, MD (US);
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on Jan. 15, 2003. Provisional application No. 60/447,641, filed on Feb. 14, 2003. Provisional application No. 60/456,220, filed on Mar. 20, 2003. Provisional application No. 60/469,356, filed on May 9, 2003. Provisional application No. 60/482,112, filed on Jun. 24, 2003. Provisional application No. 60/482,107, filed on Jun. 24, 2003.

(57) **ABSTRACT**

An approach is provided for a method of encoding structure Low Density Parity Check (LDPC) codes. Memory storing information representing a structured parity check matrix of Low Density Parity Check (LDPC) codes is accessed during the encoding process. The information is organized in tabular form, wherein each row represents occurrences of one values within a first column of a group of columns of the parity check matrix. The rows correspond to groups of columns of the parity check matrix, wherein subsequent columns within each of the groups are derived according to a predetermined operation. An LDPC coded signal is output based on the stored information representing the parity check matrix.

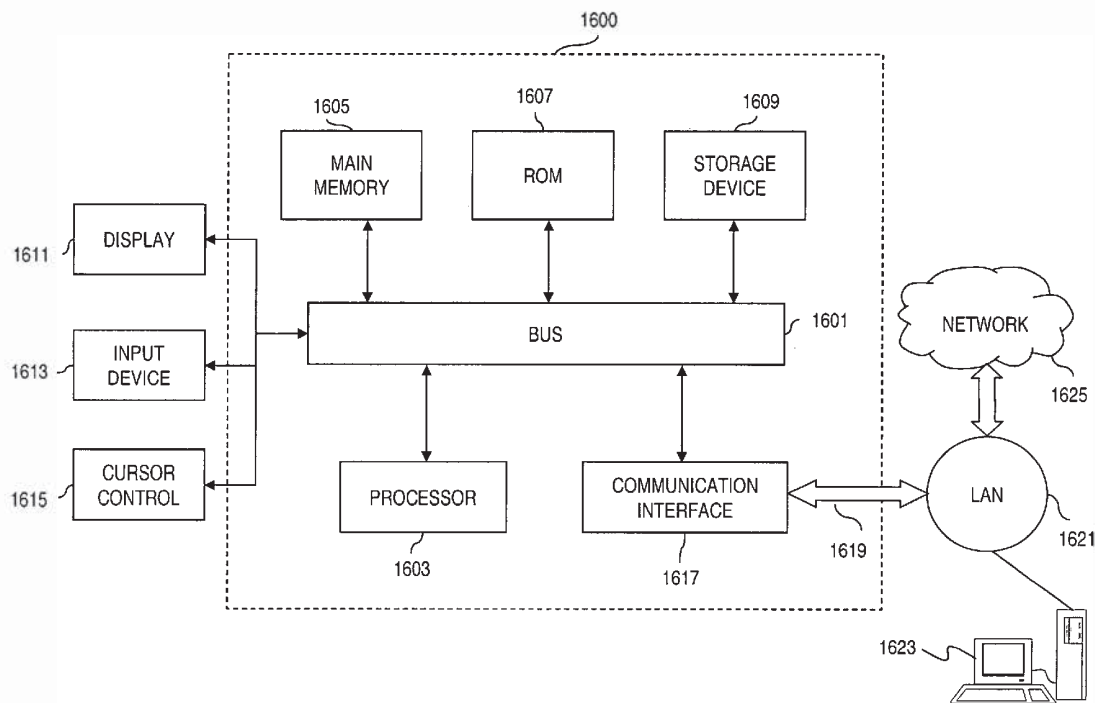


FIG. 1

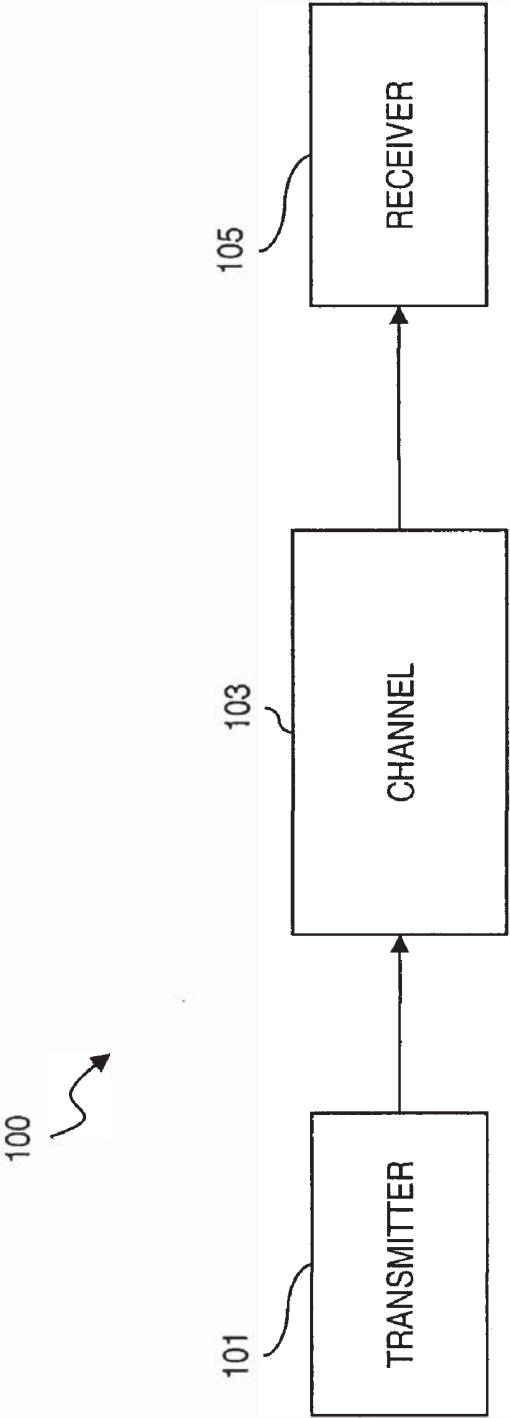


FIG. 2A

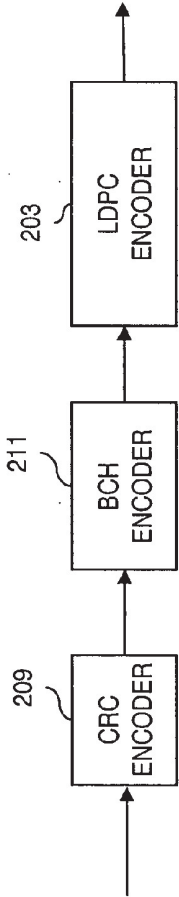
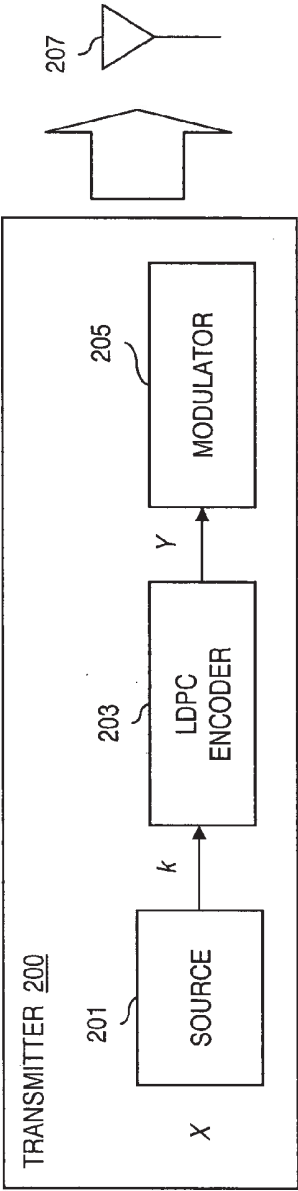


FIG. 2B

FIG. 3

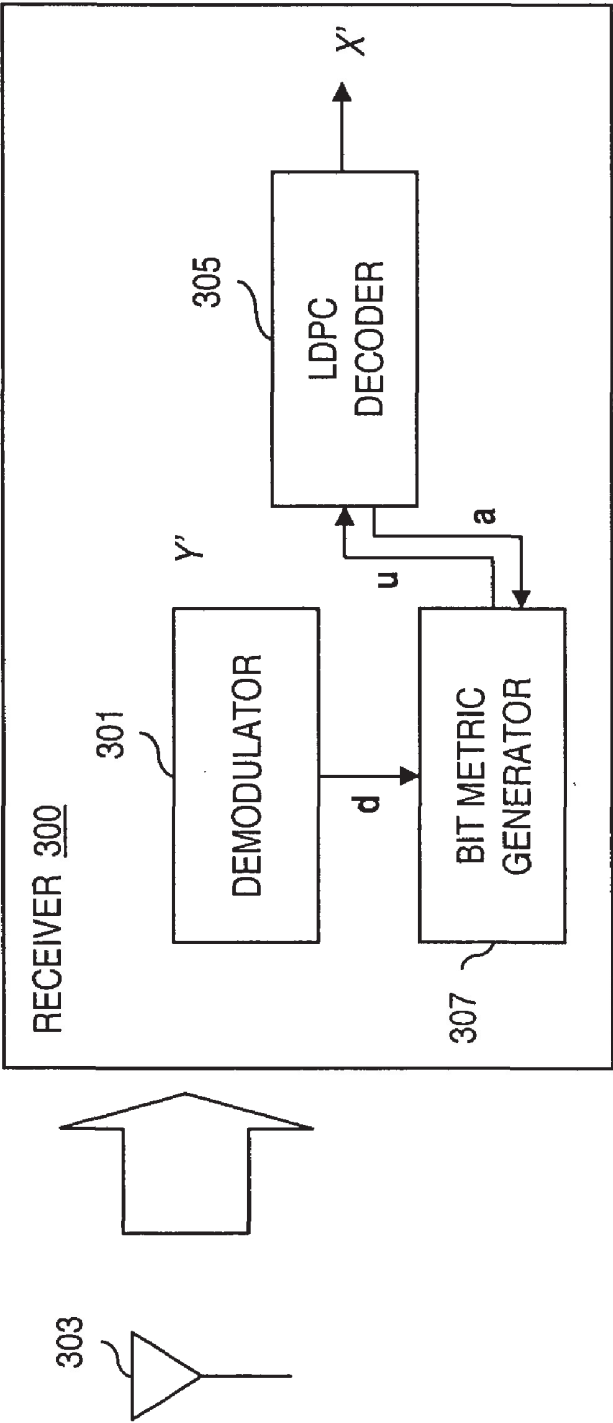


FIG. 4

$$H = \begin{matrix} & n_1 & n_2 & n_3 & n_4 & n_5 & n_6 & n_7 & n_8 \\ \begin{bmatrix} m_1 \\ m_2 \\ m_3 \\ m_4 \end{bmatrix} & \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 1 & 0 \end{bmatrix} \end{matrix}$$

FIG. 5

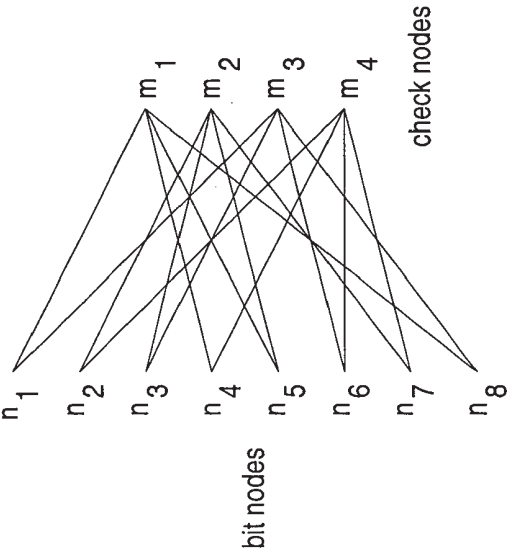


FIG. 6

$$B = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ x & x & x & x & x & x & x & x \\ x & x & x & x & x & x & x & x \\ x & x & x & x & x & x & x & x \\ x & x & x & x & x & x & x & x \\ x & x & x & x & x & x & x & x \\ x & x & x & x & x & x & x & x \\ x & x & x & x & x & x & x & x \end{bmatrix}$$

FIG. 7

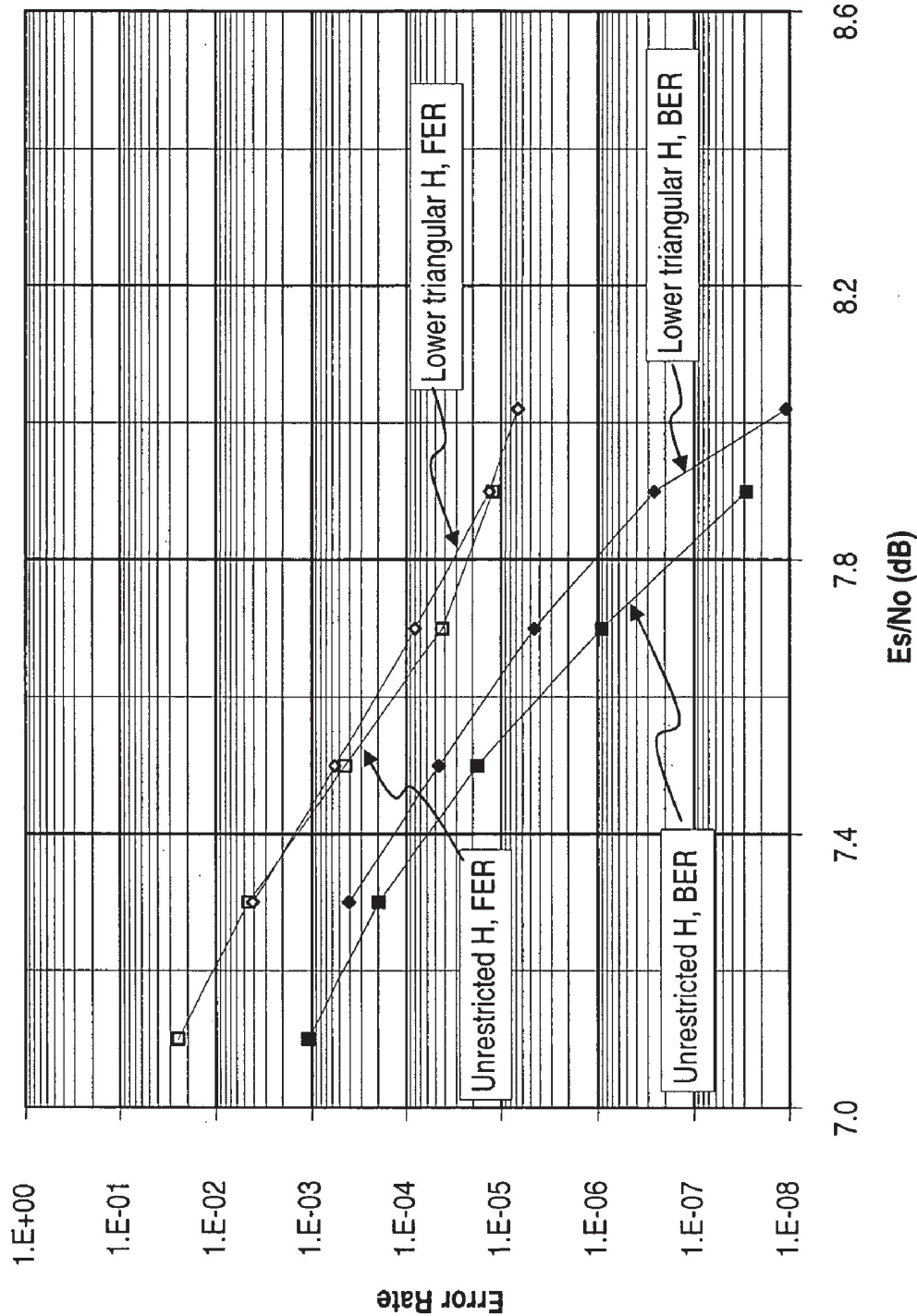


FIG. 8A

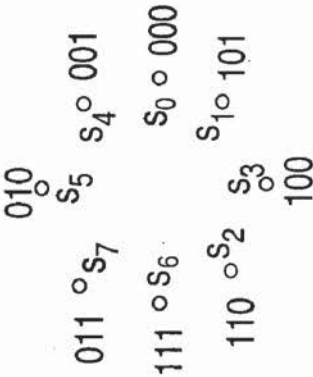
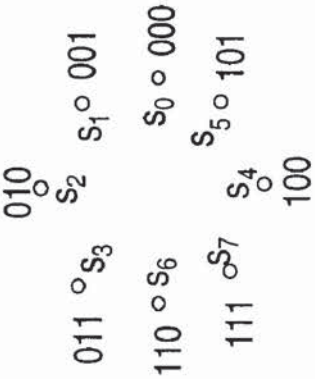
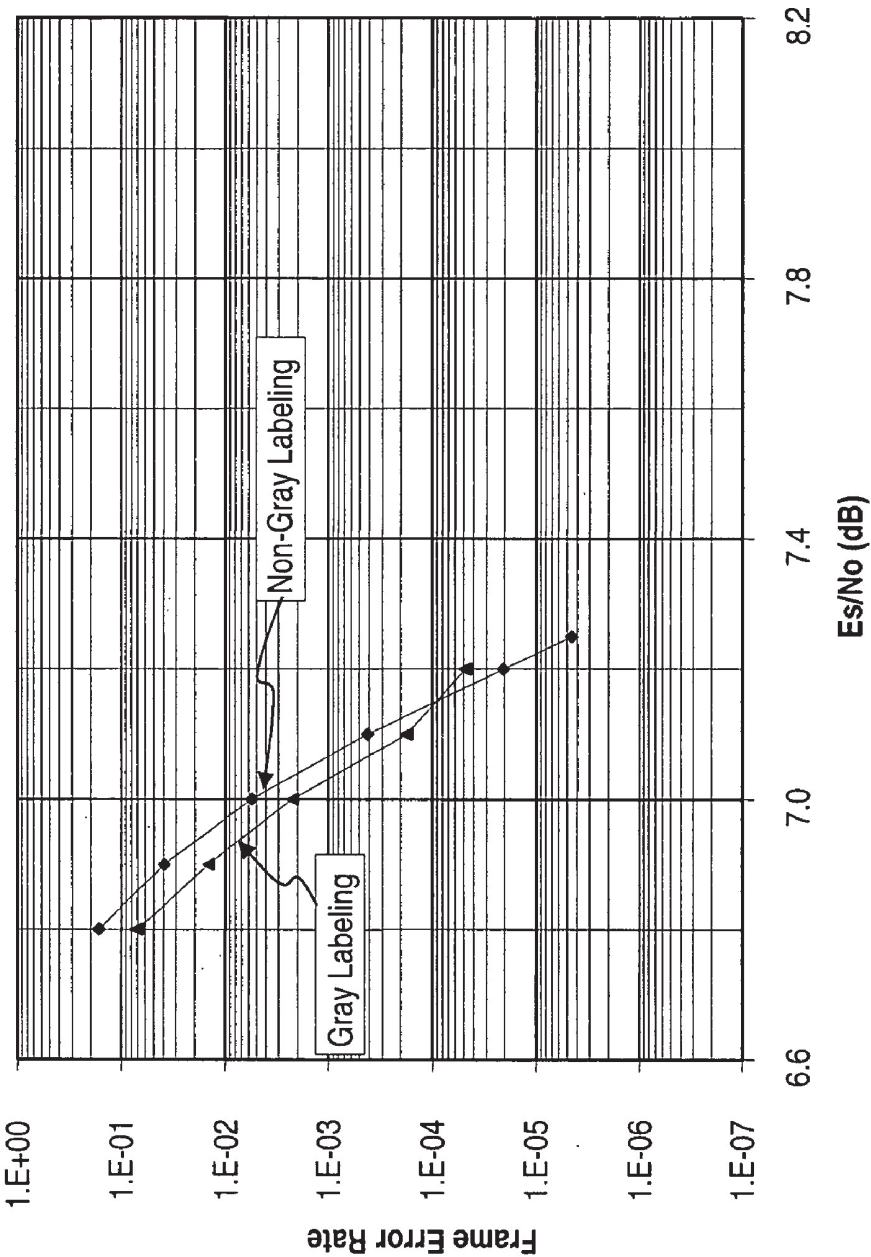
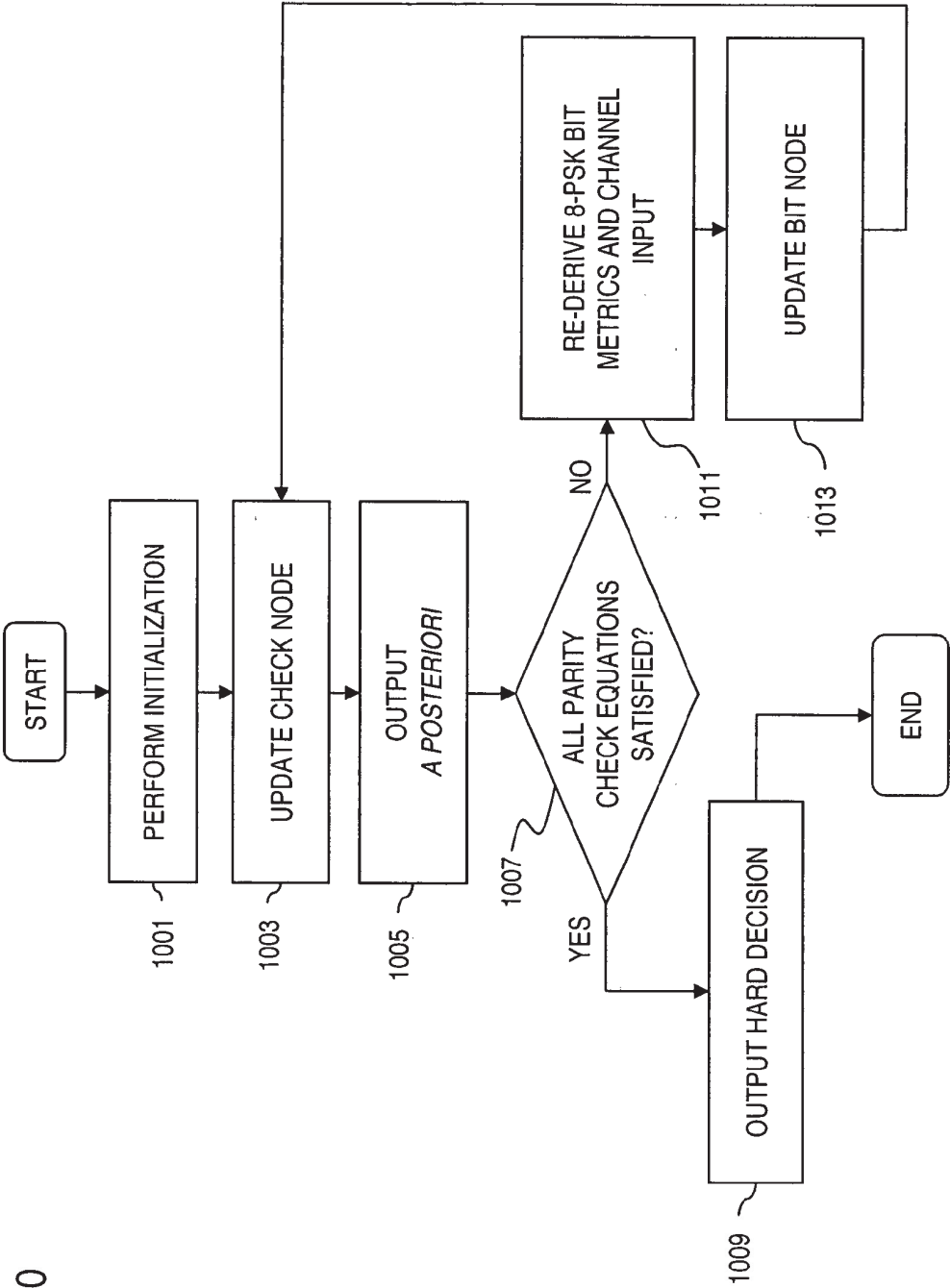


FIG. 8B

FIG. 9





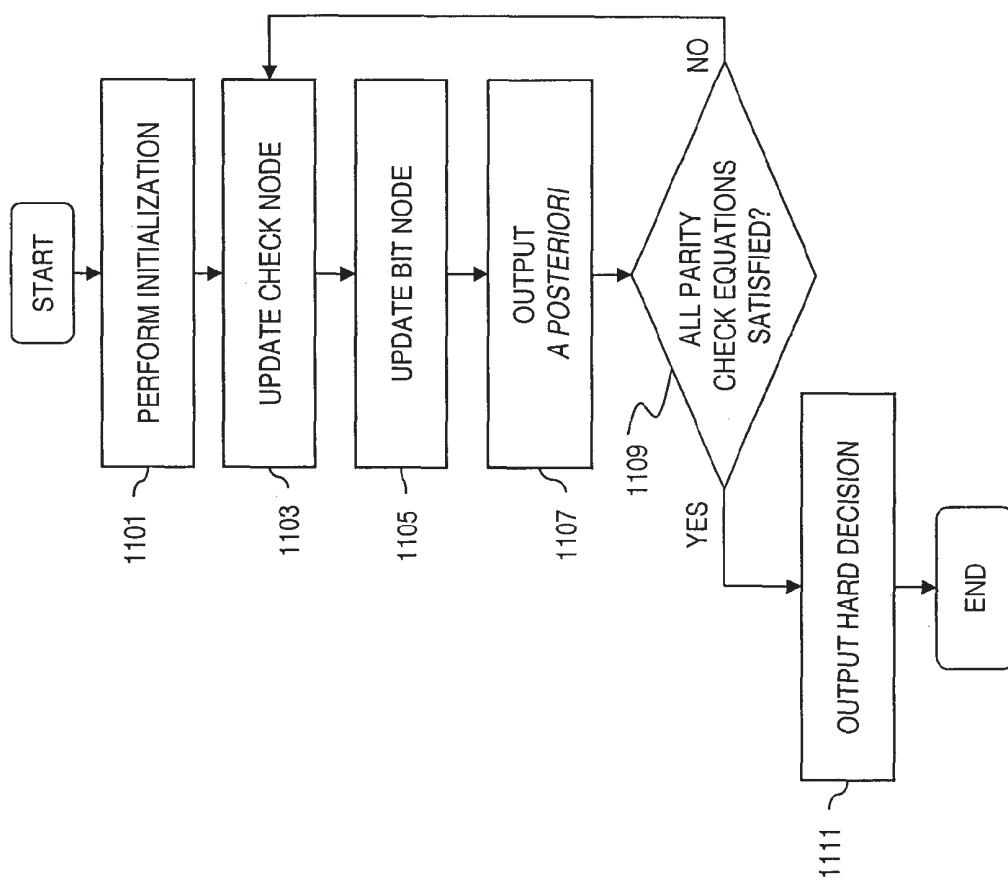


FIG. 11

FIG. 12A

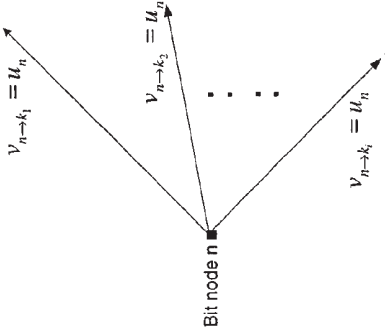


FIG. 12C

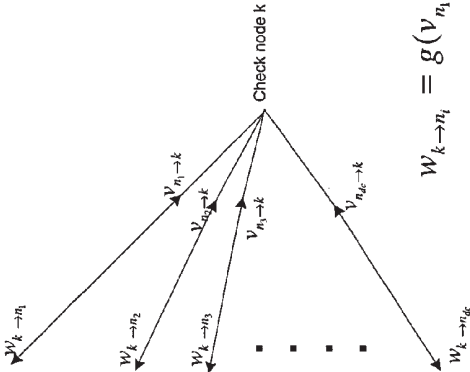
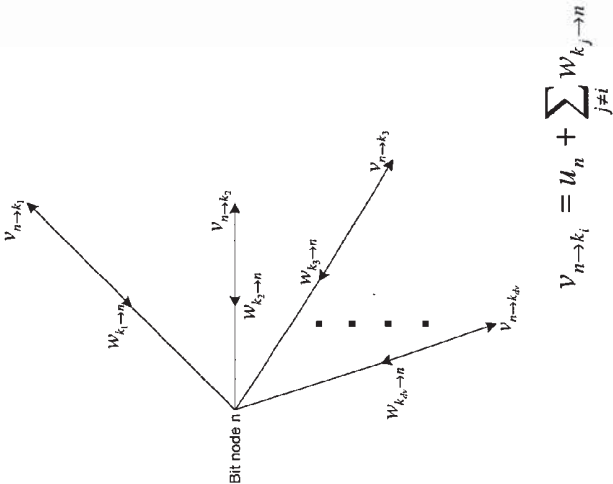


FIG. 12B

$$w_{k \rightarrow n_i} = g(v_{n_1 \rightarrow k}, v_{n_2 \rightarrow k}, \dots, v_{n_{i-1} \rightarrow k}, v_{n_{i+1} \rightarrow k}, \dots, v_{n_d \rightarrow k})$$

FIG. 13A

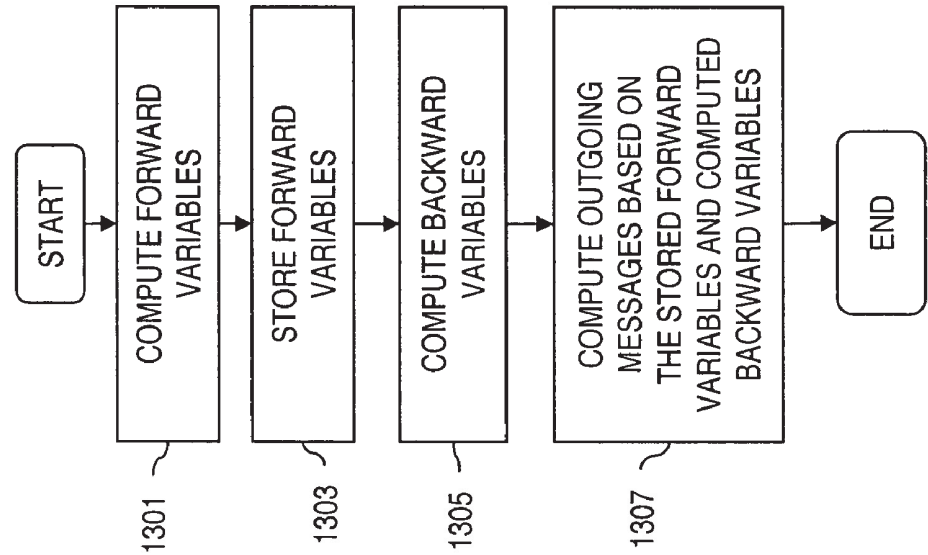


FIG. 13B

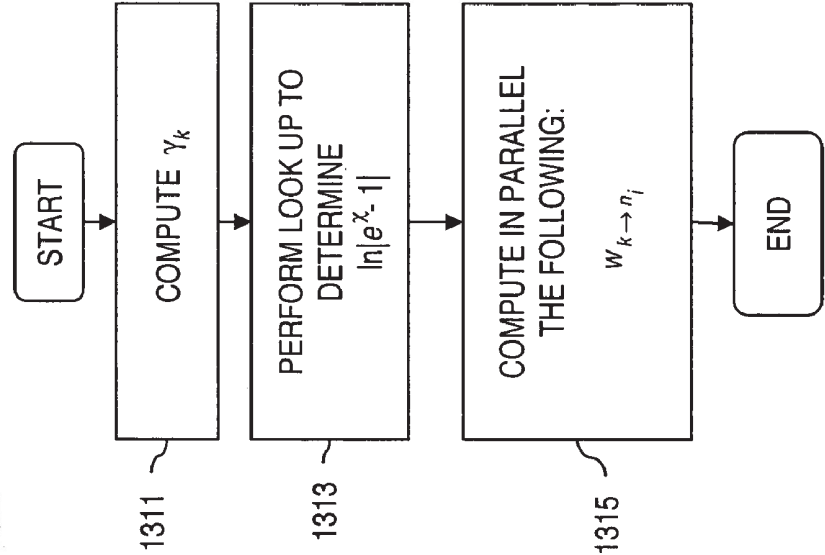


FIG. 14A

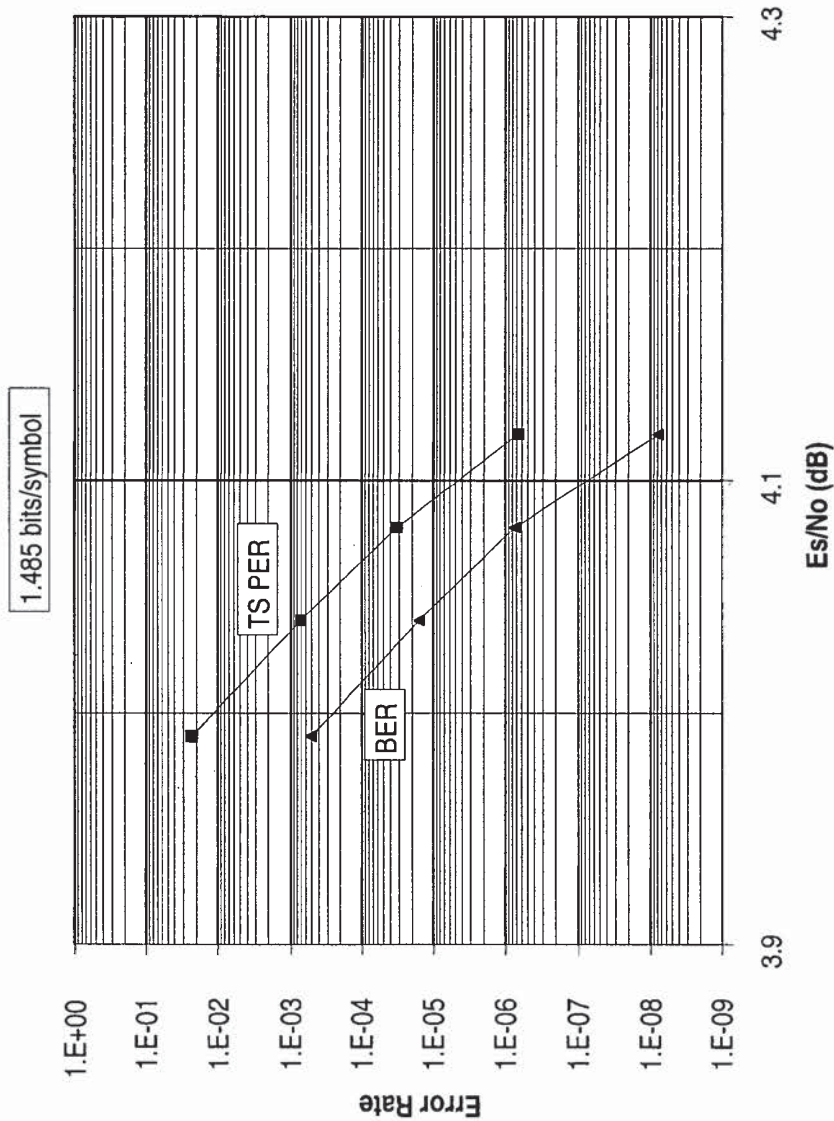


FIG. 14B

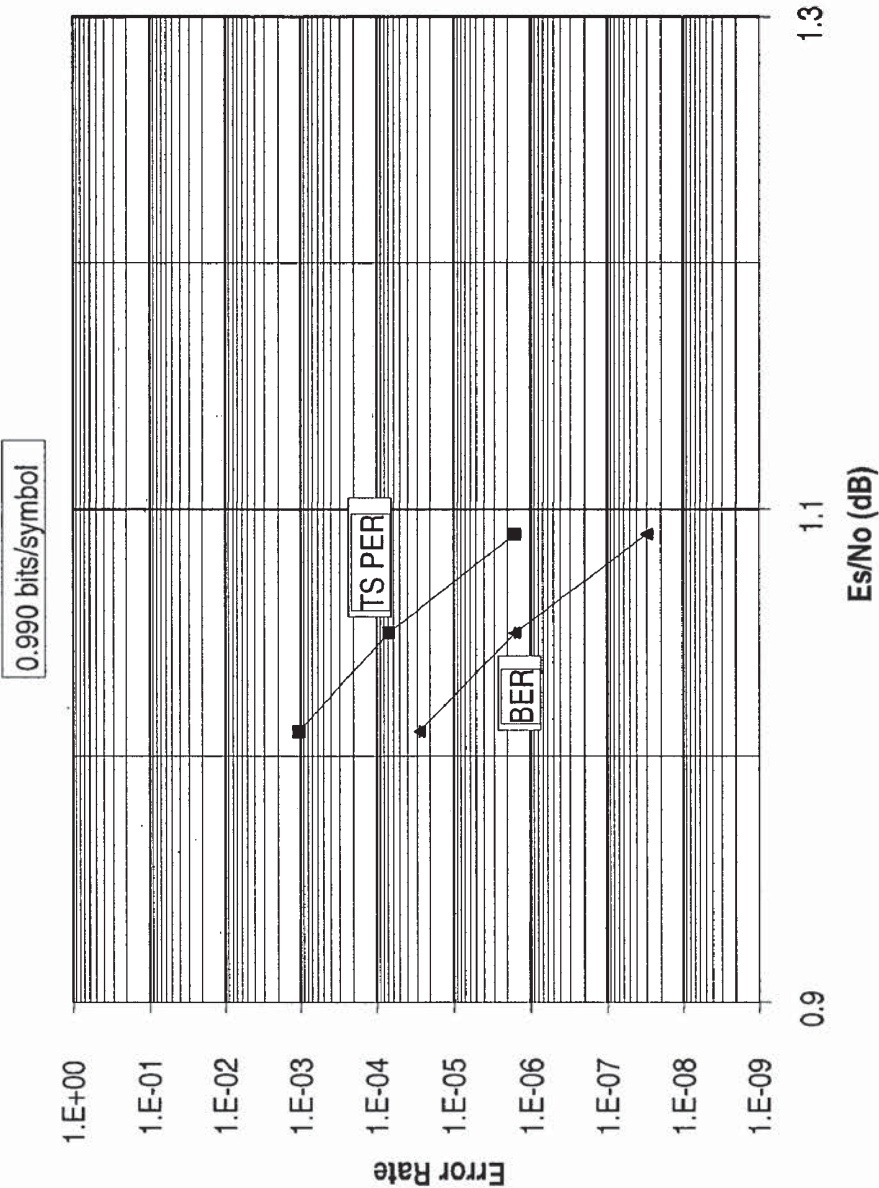


FIG. 14C

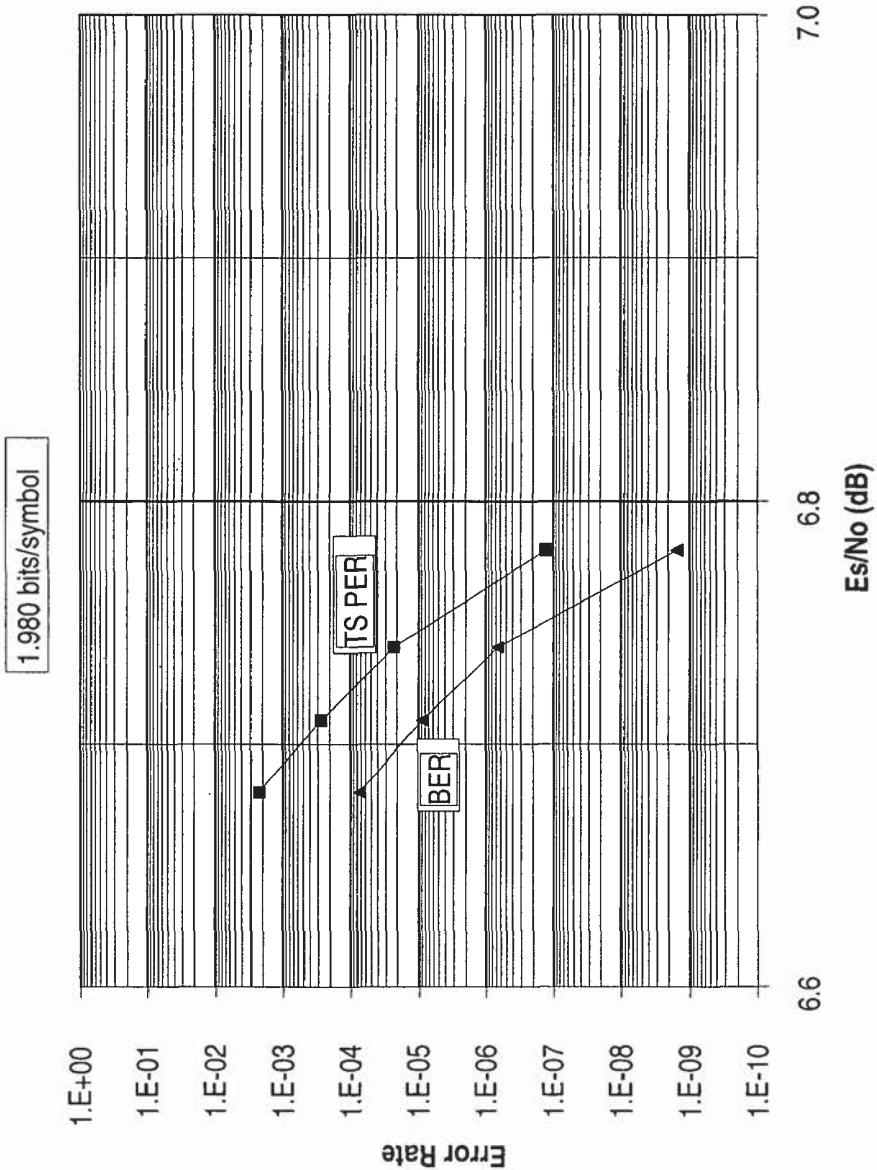
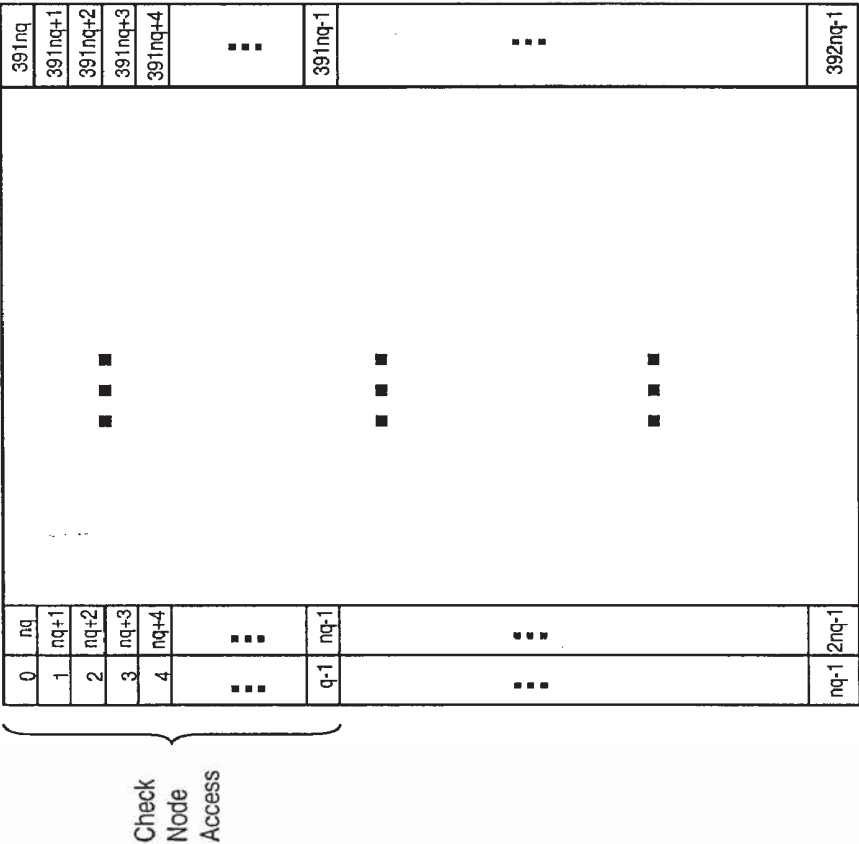


FIG. 15A



1501

FIG. 15B

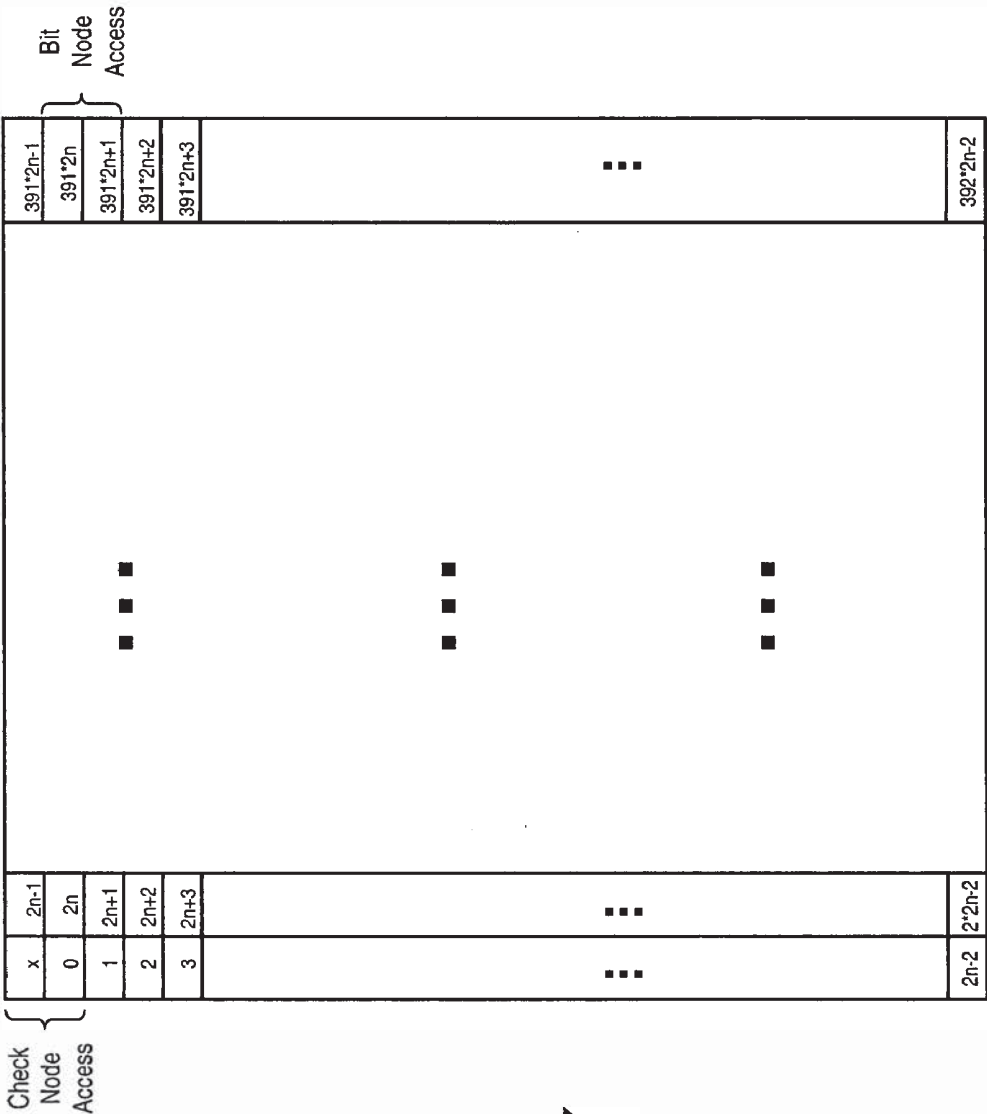
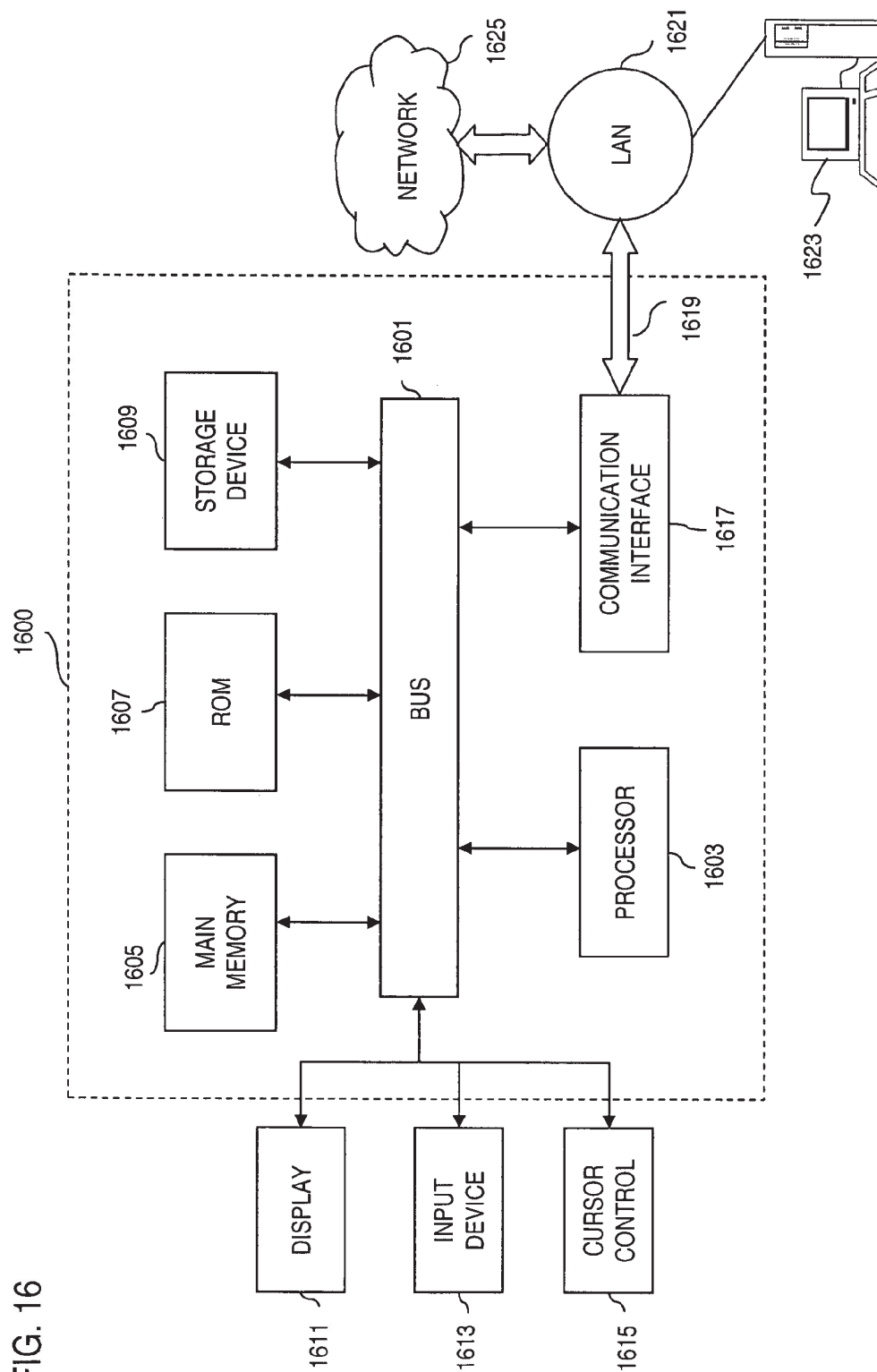


FIG. 16



METHOD AND SYSTEM FOR PROVIDING LOW DENSITY PARITY CHECK (LDPC) ENCODING

RELATED APPLICATIONS

[0001] This application is related to, and claims the benefit of the earlier filing date under 35 U.S.C. §119(e) of, U.S. Provisional Patent Application (Serial No. 60/393,457) filed Jul. 3, 2002 (Attorney Docket: PD-202095), entitled "Code Design and Implementation Improvements for Low Density Parity Check Codes," U.S. Provisional Patent Application (Serial No. 60/398,760) filed Jul. 26, 2002 (Attorney Docket: PD-202101), entitled "Code Design and Implementation Improvements for Low Density Parity Check Codes," U.S. Provisional Patent Application (Serial No. 60/403,812) filed Aug. 15, 2002 (Attorney Docket: PD-202105), entitled "Power and Bandwidth Efficient Modulation and Coding Scheme for Direct Broadcast Satellite and Broadcast Satellite Communications," U.S. Provisional Patent Application (Serial No. 60/421,505), filed Oct. 25, 2002 (Attorney Docket: PD-202101), entitled "Method and System for Generating Low Density Parity Check Codes," U.S. Provisional Patent Application (Serial No. 60/421,999), filed Oct. 29, 2002 (Attorney Docket: PD-202105), entitled "Satellite Communication System Utilizing Low Density Parity Check Codes," U.S. Provisional Patent Application (Serial No. 60/423,710), filed Nov. 4, 2002 (Attorney Docket: PD-202101), entitled "Code Design and Implementation Improvements for Low Density Parity Check Codes," U.S. Provisional Patent Application (Serial No. 60/440,199) filed Jan. 15, 2003 (Attorney Docket: PD-203009), entitled "Novel Solution to Routing Problem in Low Density Parity Check Decoders," U.S. Provisional Patent Application (Serial No. 60/447,641) filed Feb. 14, 2003 (Attorney Docket: PD-203016), entitled "Low Density Parity Check Code Encoder Design," U.S. Provisional Patent Application (Serial No. 60/456,220) filed Mar. 20, 2003 (Attorney Docket: PD-203021), entitled "Description LDPC and BCH Encoders," U.S. Provisional Patent Application filed May 9, 2003 (Attorney Docket: PD-203030), entitled "Description LDPC and BCH Encoders," U.S. Provisional Patent Application filed Jun. 24, 2003 (Attorney Docket: PD-203044), entitled "Description LDPC and BCH Encoders," and U.S. Provisional Patent Application filed Jun. 24, 2003 (Attorney Docket: PD-203059), entitled "Description LDPC and BCH Encoders"; the entireties of which are incorporated herein by reference.

FIELD OF THE INVENTION

[0002] The present invention relates to communication systems, and more particularly to coded systems.

BACKGROUND OF THE INVENTION

[0003] Communication systems employ coding to ensure reliable communication across noisy communication channels. These communication channels exhibit a fixed capacity that can be expressed in terms of bits per symbol at certain signal to noise ratio (SNR), defining a theoretical upper limit (known as the Shannon limit). As a result, coding design has aimed to achieve rates approaching this Shannon limit. One such class of codes that approach the Shannon limit is Low Density Parity Check (LDPC) codes.

[0004] Traditionally, LDPC codes have not been widely deployed because of a number of drawbacks. One drawback

is that the LDPC encoding technique is highly complex. Encoding an LDPC code using its generator matrix would require storing a very large, non-sparse matrix. Additionally, LDPC codes require large blocks to be effective; consequently, even though parity check matrices of LDPC codes are sparse, storing these matrices is problematic.

[0005] From an implementation perspective, a number of challenges are confronted. For example, storage is an important reason why LDPC codes have not become widespread in practice. Also, a key challenge in LDPC code implementation has been how to achieve the connection network between several processing engines (nodes) in the decoder. Further, the computational load in the decoding process, specifically the check node operations, poses a problem.

[0006] Therefore, there is a need for an LDPC communication system that employs simple encoding and decoding processes. There is also a need for using LDPC codes efficiently to support high data rates, without introducing greater complexity. There is also a need to improve performance of LDPC encoders and decoders. There is also a need to minimize storage requirements for implementing LDPC coding. There is a further need for a scheme that simplifies the communication between processing nodes in the LDPC decoder.

SUMMARY OF THE INVENTION

[0007] These and other needs are addressed by the present invention, wherein an approach for encoding structured Low Density Parity Check (LDPC) codes is provided. Structure of the LDPC codes is provided by restricting portion part of the parity check matrix to be lower triangular and/or satisfying other requirements such that the communication between bit nodes and check nodes of the decoder is simplified. Memory storing information representing the structured parity check matrix is accessed. The information is organized in tabular form, wherein each row represents occurrences of one values within a first column of a group of columns of the parity check matrix. The rows correspond to groups of columns of the parity check matrix, wherein subsequent columns within each of the groups are derived according to a predetermined operation (e.g., cyclic shift, addition, etc.). An LDPC coded signal based on the stored information representing the parity check matrix. According to one embodiment of the present invention, a Bose Chaudhuri Hocquenghem (BCH) encoder is utilized by the transmitter to encode an input signal using BCH codes, wherein the output LDPC coded signal corresponding to the input signal represents a code having an outer BCH code and an inner LDPC code. Further, a cyclic redundancy check (CRC) encoder is supplied to encode the input signal according to a CRC code. The approach advantageously provides expedient encoding as well as decoding of LDPC codes.

[0008] According to one aspect of an embodiment of the present invention, a method of encoding is disclosed. The method includes accessing memory storing information representing a structured parity check matrix of Low Density Parity Check (LDPC) codes. The information is organized in tabular form, wherein each row represents occurrences of one values within a first column of a group of columns of the parity check matrix. The rows correspond to groups of columns of the parity check matrix, wherein subsequent columns within each of the groups are derived according to

a predetermined operation. The method also includes outputting an LDPC coded signal based on the stored information representing the parity check matrix.

[0009] According to another aspect of an embodiment of the present invention, an encoder for generating Low Density Parity Check (LDPC) codes disclosed. The encoder includes memory storing information representing a structured parity check matrix of the LDPC codes. The information is organized in tabular form, wherein each row represents occurrences of one values within a first column of a group of columns of the parity check matrix. The rows correspond to groups of columns of the parity check matrix, wherein subsequent columns within each of the groups are derived according to a predetermined operation. The encoder also includes means for retrieving the stored information representing the parity check matrix to output an LDPC coded signal.

[0010] According to another aspect of an embodiment of the present invention, a transmitter utilizing Low Density Parity Check (LDPC) coding disclosed. The transmitter includes memory storing information representing a structured parity check matrix of the LDPC codes, the information being organized in tabular form, wherein each row represents occurrences of one values within a first column of a group of columns of the parity check matrix. The rows correspond to groups of columns of the parity check matrix, wherein subsequent columns within each of the groups are derived according to a predetermined operation. The transmitter also includes an LDPC encoder configured to access the stored information from the memory to output an LDPC coded signal.

[0011] Still other aspects, features, and advantages of the present invention are readily apparent from the following detailed description, simply by illustrating a number of particular embodiments and implementations, including the best mode contemplated for carrying out the present invention. The present invention is also capable of other and different embodiments, and its several details can be modified in various obvious respects, all without departing from the spirit and scope of the present invention. Accordingly, the drawing and description are to be regarded as illustrative in nature, and not as restrictive.

BRIEF DESCRIPTION OF THE DRAWINGS

[0012] The present invention is illustrated by way of example, and not by way of limitation, in the figures of the accompanying drawings and in which like reference numerals refer to similar elements and in which:

[0013] FIG. 1 is a diagram of a communications system configured to utilize Low Density Parity Check (LDPC) codes, according to an embodiment of the present invention;

[0014] FIGS. 2A and 2B are diagrams of exemplary LDPC encoders deployed in the transmitter of FIG. 1;

[0015] FIG. 3 is a diagram of an exemplary receiver in the system of FIG. 1;

[0016] FIG. 4 is a diagram of a sparse parity check matrix, in accordance with an embodiment of the present invention;

[0017] FIG. 5 is a diagram of a bipartite graph of an LDPC code of the matrix of FIG. 4;

[0018] FIG. 6 is a diagram of a sub-matrix of a sparse parity check matrix, wherein the sub-matrix contains parity check values restricted to the lower triangular region, according to an embodiment of the present invention;

[0019] FIG. 7 is a graph showing performance between codes utilizing unrestricted parity check matrix (H matrix) versus restricted H matrix having a sub-matrix as in FIG. 6;

[0020] FIGS. 8A and 8B are, respectively, a diagram of a non-Gray 8-PSK modulation scheme, and a Gray 8-PSK modulation, each of which can be used in the system of FIG. 1;

[0021] FIG. 9 is a graph showing performance between codes utilizing Gray labeling versus non-Gray labeling;

[0022] FIG. 10 is a flow chart of the operation of the LDPC decoder using non-Gray mapping, according to an embodiment of the present invention;

[0023] FIG. 11 is a flow chart of the operation of the LDPC decoder of FIG. 3 using Gray mapping, according to an embodiment of the present invention;

[0024] FIGS. 12A-12C are diagrams of the interactions between the check nodes and the bit nodes in a decoding process, according to an embodiment of the present invention;

[0025] FIGS. 13A and 13B are flowcharts of processes for computing outgoing messages between the check nodes and the bit nodes using, respectively, a forward-backward approach and a parallel approach, according to various embodiments of the present invention;

[0026] FIGS. 14A-14 are graphs showing simulation results of LDPC codes generated in accordance with various embodiments of the present invention;

[0027] FIGS. 15A and 15B are diagrams of the top edge and bottom edge, respectively, of memory organized to support structured access as to realize randomness in LDPC coding, according to an embodiment of the present invention; and

[0028] FIG. 16 is a diagram of a computer system that can perform the processes of encoding and decoding of LDPC codes, in accordance with embodiments of the present invention.

DESCRIPTION OF THE PREFERRED EMBODIMENT

[0029] A system, method, and software for efficiently decoding structured Low Density Parity Check (LDPC) codes are described. In the following description, for the purposes of explanation, numerous specific details are set forth in order to provide a thorough understanding of the present invention. It is apparent, however, to one skilled in the art that the present invention may be practiced without these specific details or with an equivalent arrangement. In other instances, well-known structures and devices are shown in block diagram form in order to avoid unnecessarily obscuring the present invention.

[0030] FIG. 1 is a diagram of a communications system configured to utilize Low Density Parity Check (LDPC) codes, according to an embodiment of the present invention. A digital communications system 100 includes a transmitter

101 that generates signal waveforms across a communication channel **103** to a receiver **105**. In this discrete communications system **100**, the transmitter **101** has a message source that produces a discrete set of possible messages; each of the possible messages has a corresponding signal waveform. These signal waveforms are attenuated, or otherwise altered, by communications channel **103**. To combat the noise channel **103**, LDPC, codes are utilized.

[0031] The LDPC codes that are generated by the transmitter **101** enable high speed implementation without incurring any performance loss. These structured LDPC codes output from the transmitter **101** avoid assignment of a small number of check nodes to the bit nodes already vulnerable to channel errors by virtue of the modulation scheme (e.g., 8-PSK).

[0032] Such LDPC codes have a parallelizable decoding algorithm (unlike turbo codes), which advantageously involves simple operations such as addition, comparison and table look-up. Moreover, carefully designed LDPC codes do not exhibit any sign of error floor.

[0033] According to one embodiment of the present invention, the transmitter **101** generates, using a relatively simple encoding technique, LDPC codes based on parity check matrices (which facilitate efficient memory access during decoding) to communicate with the receiver **105**. The transmitter **101** employs LDPC codes that can outperform concatenated turbo+RS (Reed-Solomon) codes, provided the block length is sufficiently large.

[0034] FIGS. 2A and 2B are diagrams of exemplary LDPC encoders deployed in the transmitter of FIG. 1. As seen in FIG. 2A, a transmitter **200** is equipped with an LDPC encoder **203** that accepts input from an information source **201** and outputs coded stream of higher redundancy suitable for error correction processing at the receiver **105**. The information source **201** generates k signals from a discrete alphabet, X . LDPC codes are specified with parity check matrices. On the other hand, encoding LDPC codes require, in general, specifying the generator matrices. Even though it is possible to obtain generator matrices from parity check matrices using Gaussian elimination, the resulting matrix is no longer sparse and storing a large generator matrix can be complex.

[0035] Encoder **203** generates signals from alphabet Y to a modulator **205** using a simple encoding technique that makes use of only the parity check matrix by imposing structure onto the parity check matrix. Specifically, a restriction is placed on the parity check matrix by constraining certain portion of the matrix to be triangular. The construction of such a parity check matrix is described more fully below in FIG. 6. Such a restriction results in negligible performance loss, and therefore, constitutes an attractive trade-off.

[0036] Modulator **205** maps the encoded messages from encoder **203** to signal waveforms that are transmitted to a transmit antenna **207**, which emits these waveforms over the communication channel **103**. Accordingly, the encoded messages are modulated and distributed to a transmit antenna **207**. The transmissions from the transmit antenna **207** propagate to a receiver, as discussed below.

[0037] FIG. 2B shows an LDPC encoder utilized with a Bose Chaudhuri Hocquenghem (BCH) encoder and a cyclic

redundancy check (CRC) encoder, according to one embodiment of the present invention. Under this scenario, the codes generated by the LDPC encoder **203**, along with the CRC encoder **209** and the BCH encoder **211**, have a concatenated outer BCH code and inner low density parity check (LDPC) code. Furthermore, error detection is achieved using cyclic redundancy check (CRC) codes. The CRC encoder **209**, in an exemplary embodiment, encodes using an 8-bit CRC code with generator polynomial $(x^5+x^4+x^3+x^2+1)(x^2+x+1)(x+1)$.

[0038] The LDPC encoder **203** systematically encodes an information block of size k_{ldpc} , $i=(i_0, i_1, \dots, i_{k_{ldpc}-1})$ onto a codeword of size n_{ldpc} , $c=(i_0, i_1, \dots, i_{k_{ldpc}-1}, p_0, p_1, \dots, p_{n_{ldpc}-k_{ldpc}-1})$. The transmission of the codeword starts in the given order from i_0 and ends with $p_{n_{ldpc}-k_{ldpc}-1}$. LDPC code parameters (n_{ldpc}, k_{ldpc}) are given in Table 1 below.

TABLE 1

LDPC Code Parameters (n_{ldpc}, k_{ldpc})		
Code Rate	LDPC Uncoded Block Length k_{ldpc}	LDPC Coded Block Length n_{ldpc}
1/2	32400	64800
2/3	43200	64800
3/4	48600	64800
4/5	51840	64800
5/6	54000	64800
3/5	38880	64800
8/9	57600	64800
9/10	58320	64800

[0039] The task of the LDPC encoder **203** is to determine $n_{ldpc}-k_{ldpc}$ parity bits $(p_0, p_1, \dots, p_{n_{ldpc}-k_{ldpc}-1})$ for every block of k_{ldpc} information bits, $(i_0, i_1, \dots, i_{k_{ldpc}-1})$. The procedure is as follows. First, the parity bits are initialized; $p_0=p_1=p_2=\dots=p_{n_{ldpc}-k_{ldpc}-1}=0$. The first information bit, i_0 , are accumulated at parity bit addresses specified in the first row of Tables 3 through 10. For example, for rate $2/3$ (Table 3), the following results:

$$\begin{aligned}
 p_0 &= p_0 \oplus i_0 \\
 p_{10491} &= p_{10491} \oplus i_0 \\
 p_{16043} &= p_{16043} \oplus i_0 \\
 p_{506} &= p_{506} \oplus i_0 \\
 p_{12826} &= p_{12826} \oplus i_0 \\
 p_{8065} &= p_{8065} \oplus i_0 \\
 p_{8226} &= p_{8226} \oplus i_0 \\
 p_{2767} &= p_{2767} \oplus i_0 \\
 p_{240} &= p_{240} \oplus i_0 \\
 p_{18673} &= p_{18673} \oplus i_0 \\
 p_{9279} &= p_{9279} \oplus i_0 \\
 p_{10579} &= p_{10579} \oplus i_0 \\
 p_{20928} &= p_{20928} \oplus i_0
 \end{aligned}$$

[0040] (All additions are in GF(2)).

[0041] Then, for the next 359 information bits, $i_m, m=1, 2, \dots, 359$, these bits are accumulated at parity bit addresses $\{x+m \bmod 360 \times q\} \bmod (n_{ldpc}-k_{ldpc})$, where x denotes the address of the parity bit accumulator corresponding to the first bit i_0 , and q is a code rate dependent constant specified

in Table 2. Continuing with the example, $q=60$ for rate $\frac{2}{3}$. By way of example, for information bit i_1 , the following operations are performed:

$$\begin{aligned} P_{60} &= P_{60} \oplus i_1 \\ P_{10551} &= P_{10551} \oplus i_1 \\ P_{16103} &= P_{16103} \oplus i_1 \\ P_{566} &= P_{566} \oplus i_1 \\ P_{12886} &= P_{12886} \oplus i_1 \\ P_{8125} &= P_{8125} \oplus i_1 \\ P_{8286} &= P_{8286} \oplus i_1 \\ P_{2827} &= P_{2827} \oplus i_1 \\ P_{300} &= P_{300} \oplus i_1 \\ P_{18733} &= P_{18733} \oplus i_1 \\ P_{9339} &= P_{9339} \oplus i_1 \\ P_{10639} &= P_{10639} \oplus i_1 \\ P_{20988} &= P_{20988} \oplus i_1 \end{aligned}$$

[0042] For the 361st information bit i_{360} , the addresses of the parity bit accumulators are given in the second row of the Tables 3 through 10. In a similar manner the addresses of the parity bit accumulators for the following 359 information bits $i_m, m=361, 362, \dots, 719$ are obtained using the formula $\{x+m \bmod 360 \times q\} \bmod (n_{ldpc} - k_{ldpc})$, where x denotes the address of the parity bit accumulator corresponding to the

information bit i_{360} , i.e., the entries in the second row of the Tables 3-10. In a similar manner, for every group of 360 new information bits, a new row from Tables 3 through 10 are used to find the addresses of the parity bit accumulators.

[0043] After all of the information bits are exhausted, the final parity bits are obtained as follows. First, the following operations are performed, starting with $i=1$

$$p_i = p_i \oplus p_{i-1}, \quad i=1, 2, \dots, n_{ldpc} - k_{ldpc} - 1.$$

[0044] Final content of $p_i, i=0, 1, \dots, n_{ldpc} - k_{ldpc} - 1$ is equal to the parity bit p_i .

TABLE 2

Code Rate	q
2/3	60
5/6	30
1/2	90
3/4	45
4/5	36
3/5	72
8/9	20
9/10	18

[0045]

TABLE 3

Address of Parity Bit Accumulators (Rate 2/3)															
0	10491	16043	506	12826	8065	8226	2767	240	18673	9279	10579	20928			
1	17819	8313	6433	6224	5120	5824	12812	17187	9940	13447	13825	18483			
2	17957	6024	8681	18628	12794	5915	14576	10970	12064	20437	4455	7151			
3	19777	6183	9972	14536	8182	17749	11341	5556	4379	17434	15477	18532			
4	4651	19689	1608	659	16707	14335	6143	3058	14618	17894	20684	5306			
5	9778	2552	12096	12369	15198	16890	4851	3109	1700	18725	1997	15882			
6	486	6111	13743	11537	5591	7433	15227	14145	1483	3887	17431	12430			
7	20647	14311	11734	4180	8110	5525	12141	15761	18661	18441	10569	8192			
8	3791	14759	15264	19918	10132	9062	10010	12786	10675	9682	19246	5454			
9	19525	9485	7777	19999	8378	9209	3163	20232	6690	16518	716	7353			
10	4588	6709	20202	10905	915	4317	11073	13576	16433	368	3508	21171			
11	14072	4033	19959	12608	631	19494	14160	8249	10223	21504	12395	4322			
12	13800	14161													
13	2948	9647													
14	14693	16027													
15	20506	11082													
16	1143	9020													
17	13501	4014													
18	1548	2190													
19	12216	21556													
20	2095	19897													
21	4189	7958													
22	15940	10048													
23	515	12614													
24	8501	8450													
25	17595	16784													
26	5913	8495													
27	16394	10423													
28	7409	6981													
29	6678	15939													
30	20344	12987													
31	2510	14588													
32	17918	6655													
33	6703	19451													
34	496	4217													
35	7290	5766													
36	10521	8925													
37	20379	11905													
38	4090	5838													
39	19082	17040													

TABLE 3-continued

Address of Parity Bit Accumulators (Rate 2/3)	
40 20233 12352	
41 19365 19546	
42 6249 19030	
43 11037 19193	
44 19760 11772	
45 19644 7428	
46 16076 3521	
47 11779 21062	
48 13062 9682	
49 8934 5217	
50 11087 3319	
51 18892 4356	
52 7894 3898	
53 5963 4360	
54 7346 11726	
55 5182 5609	
56 2412 17295	
57 9845 20494	
58 6687 1864	
59 20564 5216	
0 18226 17207	
1 9380 8266	
2 7073 3065	
3 18252 13437	
4 9161 15642	
5 10714 10153	
6 11585 9078	
7 5359 9418	
8 9024 9515	
9 1206 16354	
10 14994 1102	
11 9375 20796	
12 15964 6027	
13 14789 6452	
14 8002 18591	
15 14742 14089	
16 253 3045	
17 1274 19286	
18 14777 2044	
19 13920 9900	
20 452 7374	
21 18206 9921	
22 6131 5414	
23 10077 9726	
24 12045 5479	
25 4322 7990	
26 15616 5550	
27 15561 10661	
28 20718 7387	
29 2518 18804	
30 8984 2600	
31 6516 17909	
32 11148 98	
33 20559 3704	
34 7510 1569	
35 16000 11692	
36 9147 10303	
37 16650 191	
38 15577 18685	
39 17167 20917	
40 4256 3391	
41 20092 17219	
42 9218 5056	
43 18429 8472	
44 12093 20753	
45 16345 12748	
46 16023 11095	
47 5048 17595	
48 18995 4817	
49 16483 3536	
50 1439 16148	
51 3661 3039	
52 19010 18121	
53 8968 11793	

TABLE 3-continued

Address of Parity Bit Accumulators (Rate 2/3)	
54	13427 18003
55	5303 3083
56	531 16668
57	4771 6722
58	5695 7960
59	3589 14630

[0046]

TABLE 4

Address of Parity Bit Accumulators (Rate 5/6)	
0	4362 416 8909 4156 3216 3112 2560 2912 6405 8593 4969 6723
1	2479 1786 8978 3011 4339 9313 6397 2957 7288 5484 6031 10217
2	10175 9009 9889 3091 4985 7267 4092 8874 5671 2777 2189 8716
3	9052 4795 3924 3370 10058 1128 9996 10165 9360 4297 434 5138
4	2379 7834 4835 2327 9843 804 329 8353 7167 3070 1528 7311
5	3435 7871 348 3693 1876 6585 10340 7144 5870 2084 4052 2780
6	3917 3111 3476 1304 10331 5939 5199 1611 1991 699 8316 9960
7	6883 3237 1717 10752 7891 9764 4745 3888 10009 4176 4614 1567
8	10587 2195 1689 2968 5420 2580 2883 6496 111 6023 1024 4449
9	3786 8593 2074 3321 5057 1450 3840 5444 6572 3094 9892 1512
10	8548 1848 10372 4585 7313 6536 6379 1766 9462 2456 5606 9975
11	8204 10593 7935 3636 3882 394 5968 8561 2395 7289 9267 9978
12	7795 74 1633 9542 6867 7352 6417 7568 10623 725 2531 9115
13	7151 2482 4260 5003 10105 7419 9203 6691 8798 2092 8263 3755
14	3600 570 4527 200 9718 6771 1995 8902 5446 768 1103 6520
15	6304 7621
16	6498 9209
17	7293 6786
18	5950 1708
19	8521 1793
20	6174 7854
21	9773 1190
22	9517 10268
23	2181 9349
24	1949 5560
25	1556 555
26	8600 3827
27	5072 1057
28	7928 3542
29	3226 3762
0	7045 2420
1	9645 2641
2	2774 2452
3	5331 2031
4	9400 7503
5	1850 2338
6	10456 9774
7	1692 9276
8	10037 4038
9	3964 338
10	2640 5087
11	858 3473
12	5582 5683
13	9523 916
14	4107 1559
15	4506 3491
16	8191 4182
17	10192 6157
18	5668 3305
19	3449 1540
20	4766 2697
21	4069 6675
22	1117 1016
23	5619 3085
24	8483 8400
25	8255 394
26	6338 5042
27	6174 5119

TABLE 4-continued

Address of Parity Bit Accumulators (Rate 5/6)	
28	7203 1989
29	1781 5174
0	1464 3559
1	3376 4214
2	7238 67
3	10595 8831
4	1221 6513
5	5300 4652
6	1429 9749
7	7878 5131
8	4435 10284
9	6331 5507
10	6662 4941
11	9614 10238
12	8400 8025
13	9156 5630
14	7067 8878
15	9027 3415
16	1690 3866
17	2854 8469
18	6206 630
19	363 5453
20	4125 7008
21	1612 6702
22	9069 9226
23	5767 4060
24	3743 9237
25	7018 5572
26	8892 4536
27	853 6064
28	8069 5893
29	2051 2885
0	10691 3153
1	3602 4055
2	328 1717
3	2219 9299
4	1939 7898
5	617 206
6	8544 1374
7	10676 3240
8	6672 9489
9	3170 7457
10	7868 5731
11	6121 10732
12	4843 9132
13	580 9591
14	6267 9290
15	3009 2268
16	195 2419
17	8016 1557
18	1516 9195
19	8062 9064
20	2095 8968
21	753 7326
22	6291 3833
23	2614 7844
24	2303 646
25	2075 611
26	4687 362

TABLE 4-continued

Address of Parity Bit Accumulators (Rate 5/6)
27 8684 9940
28 4830 2065
29 7038 1363
0 1769 7837
1 3801 1689
2 10070 2359
3 3667 9918
4 1914 6920
5 4244 5669
6 10245 7821
7 7648 3944
8 3310 5488
9 6346 9666
10 7088 6122
11 1291 7827
12 10592 8945
13 3609 7120
14 9168 9112
15 6203 8052
16 3330 2895
17 4264 10563
18 10556 6496
19 8807 7645
20 1999 4530
21 9202 6818
22 3403 1734
23 2106 9023
24 6881 3883
25 3895 2171
26 4062 6424
27 3755 9536
28 4683 2131
29 7347 8027

[0047]

TABLE 5

Address of Parity Bit Accumulators (Rate 1/2)
54 9318 14392 27561 26909 10219 2534 8597
55 7263 4635 2530 28130 3033 23830 3651
56 24731 23583 26036 17299 5750 792 9169
57 5811 26154 18653 11551 15447 13685 16264
58 12610 11347 28768 2792 3174 29371 12997
59 16789 16018 21449 6165 21202 15850 3186
60 31016 21449 17618 6213 12166 8334 18212
61 22836 14213 11327 5896 718 11727 9308
62 2091 24941 29966 23634 9013 15587 5444
63 22207 3983 16904 28534 21415 27524 25912
64 25687 4501 22193 14665 14798 16158 5491
65 4520 17094 23397 4264 22370 16941 21526
66 10490 6182 32370 9597 30841 25954 2762
67 22120 22865 29870 15147 13668 14955 19235
68 6689 18408 18346 9918 25746 5443 20645
69 29982 12529 13858 4746 30370 10023 24828
70 1262 28032 29888 13063 24033 21951 7863
71 6594 29642 31451 14831 9509 9335 31552
72 1358 6454 16633 20354 24598 624 5265
73 19529 295 18011 3080 13364 8032 15323
74 11981 1510 7960 21462 9129 11370 25741
75 9276 29656 4543 30699 20646 21921 28050
76 15975 25634 5520 31119 13715 21949 19605
77 18688 4608 31755 30165 13103 10706 29224
78 21514 23117 12245 26035 31656 25631 30699
79 9674 24966 31285 29908 17042 24588 31857
80 21856 27777 29919 27000 14897 11409 7122
81 29773 23310 263 4877 28622 20545 22092
82 15605 5651 21864 3967 14419 22757 15896
83 30145 1759 10139 29223 26086 10556 5098
84 18815 16575 2936 24457 26738 6030 505

TABLE 5-continued

Address of Parity Bit Accumulators (Rate 1/2)
85 30326 22298 27562 20131 26390 6247 24791
86 928 29246 21246 12400 15311 32309 18608
87 20314 6025 26689 16302 2296 3244 19613
88 6237 11943 22851 15642 23857 15112 20947
89 26403 25168 19038 18384 8882 12719 7093
0 14567 24965
1 3908 100
2 10279 240
3 24102 764
4 12383 4173
5 13861 15918
6 21327 1046
7 5288 14579
8 28158 8069
9 16583 11098
10 16681 28363
11 13980 24725
12 32169 17989
13 10907 2767
14 21557 3818
15 26676 12422
16 7676 8754
17 14905 20232
18 15719 24646
19 31942 8589
20 19978 27197
21 27060 15071
22 6071 26649
23 10393 11176
24 9597 13370
25 7081 17677
26 1433 19513
27 26925 9014
28 19202 8900
29 18152 30647
30 20803 1737
31 11804 25221
32 31683 17783
33 29694 9345
34 12280 26611
35 6526 26122
36 26165 11241
37 7666 26962
38 16290 8480
39 11774 10120
40 30051 30426
41 1335 15424
42 6865 17742
43 31779 12489
44 32120 21001
45 14508 6996
46 979 25024
47 4554 21896
48 7989 21777
49 4972 20661
50 6612 2730
51 12742 4418
52 29194 595
53 19267 20113

[0048]

TABLE 6

Address of Parity Bit Accumulators (Rate 3/4)
0 6385 7901 14611 13389 11200 3252 5243 2504 2722 821 7374
1 11359 2698 357 13824 12772 7244 6752 15310 852 2001 11417
2 7862 7977 6321 13612 12197 14449 15137 13860 1708 6399 13444
3 1560 11804 6975 13292 3646 3812 8772 7306 5795 14327 7866
4 7626 11407 14599 9689 1628 2113 10809 9283 1230 15241 4870

TABLE 6-continued

Address of Parity Bit Accumulators (Rate 3/4)
5 1610 5699 15876 9446 12515 1400 6303 5411 14181 13925 7358
6 4059 8836 3405 7853 7992 15336 5970 10368 10278 9675 4651
7 4441 3963 9153 2109 12683 7459 12030 12221 629 15212 406
8 6007 8411 5771 3497 543 14202 875 9186 6235 13908 3563
9 3232 6625 4795 546 9781 2071 7312 3399 7250 4932 12652
10 8820 10088 11090 7069 6585 13134 10158 7183 488 7455 9238
11 1903 10818 119 215 7558 11046 10615 11545 14784 7961 15619
12 3655 8736 4917 15874 5129 2134 15944 14768 7150 2692 1469
13 8316 3820 505 8923 6757 806 7957 4216 15589 13244 2622
14 14463 4852 15733 3041 11193 12860 13673 8152 6551 15108 8758
15 3149 11981
16 13416 6906
17 13098 13352
18 2009 14460
19 7207 4314
20 3312 3945
21 4418 6248
22 2669 13975
23 7571 9023
24 14172 2967
25 7271 7138
26 6135 13670
27 7490 14559
28 8657 2466
29 8599 12834
30 3470 3152
31 13917 4365
32 6024 13730
33 10973 14182
34 2464 13167
35 5281 15049
36 1103 1849
37 2058 1069
38 9654 6095
39 14311 7667
40 15617 8146
41 4588 11218
42 13660 6243
43 8578 7874
44 11741 2686
0 1022 1264
1 12604 9965
2 8217 2707
3 3156 11793
4 354 1514
5 6978 14058
6 7922 16079
7 15087 12138
8 5053 6470
9 12687 14932
10 15458 1763
11 8121 1721
12 12431 549
13 4129 7091
14 1426 8415
15 9783 7604
16 6295 11329
17 1409 12061
18 8065 9087
19 2918 8438
20 1293 14115
21 3922 13851
22 3851 4000
23 5865 1768
24 2655 14957
25 5565 6332
26 4303 12631
27 11653 12236
28 16025 7632
29 4655 14128
30 9584 13123
31 13987 9597
32 15409 12110
33 8754 15490

TABLE 6-continued

Address of Parity Bit Accumulators (Rate 3/4)
34 7416 15325
35 2909 15549
36 2995 8257
37 9406 4791
38 11111 4854
39 2812 8521
40 8476 14717
41 7820 15360
42 1179 7939
43 2357 8678
44 7703 6216
0 3477 7067
1 3931 13845
2 7675 12899
3 1754 8187
4 7785 1400
5 9213 5891
6 2494 7703
7 2576 7902
8 4821 15682
9 10426 11935
10 1810 904
11 11332 9264
12 11312 3570
13 14916 2650
14 7679 7842
15 6089 13084
16 3938 2751
17 8509 4648
18 12204 8917
19 5749 12443
20 12613 4431
21 1344 4014
22 8488 13850
23 1730 14896
24 14942 7126
25 14983 8863
26 6578 8564
27 4947 396
28 297 12805
29 13878 6692
30 11857 11186
31 14395 11493
32 16145 12251
33 13462 7428
34 14526 13119
35 2535 11243
36 6465 12690
37 6872 9334
38 15371 14023
39 8101 10187
40 11963 4848
41 15125 6119
42 8051 14465
43 11139 5167
44 2883 14521

[0049]

TABLE 7

Address of Parity Bit Accumulators (Rate 4/5)
0 149 11212 5575 6360 12559 8108 8505 408 10026 12828
1 5237 490 10677 4998 3869 3734 3092 3509 7703 10305
2 8742 5553 2820 7085 12116 10485 564 7795 2972 2157
3 2699 4304 8350 712 2841 3250 4731 10105 517 7516
4 12067 1351 11992 12191 11267 5161 537 6166 4246 2363
5 6828 7107 2127 3724 5743 11040 10756 4073 1011 3422
6 11259 1216 9526 1466 10816 940 3744 2815 11506 11573
7 4549 11507 1118 1274 11751 5207 7854 12803 4047 6484

TABLE 7-continued

Address of Parity Bit Accumulators (Rate 4/5)
8 8430 4115 9440 413 4455 2262 7915 12402 8579 7052
9 3885 9126 5665 4505 2343 253 4707 3742 4166 1556
10 1704 8936 6775 8639 8179 7954 8234 7850 8883 8713
11 11716 4344 9087 11264 2274 8832 9147 11930 6054 5455
12 7323 3970 10329 2170 8262 3854 2087 12899 9497 11700
13 4418 1467 2490 5841 817 11453 533 11217 11962 5251
14 1541 4525 7976 3457 9536 7725 3788 2982 6307 5997
15 11484 2739 4023 12107 6516 551 2572 6628 8150 9852
16 6070 1761 4627 6534 7913 3730 11866 1813 12306 8249
17 12441 5489 8748 7837 7660 2102 11341 2936 6712 11977
18 10155 4210
19 1010 10483
20 8900 10250
21 10243 12278
22 7070 4397
23 12271 3887
24 11980 6836
25 9514 4356
26 7137 10281
27 11881 2526
28 1969 11477
29 3044 10921
30 2236 8724
31 9104 6340
32 7342 8582
33 11675 10405
34 6467 12775
35 3186 12198
0 9621 11445
1 7486 5611
2 4319 4879
3 2196 344
4 7527 6650
5 10693 2440
6 6755 2706
7 5144 5998
8 11043 8033
9 4846 4435
10 4157 9228
11 12270 6562
12 11954 7592
13 7420 2592
14 8810 9636
15 689 5430
16 920 1304
17 1253 11934
18 9559 6016
19 312 7589
20 4439 4197
21 4002 9555
22 12232 7779
23 1494 8782
24 10749 3969
25 4368 3479
26 6316 5342
27 2455 3493
28 12157 7405
29 6598 11495
30 11805 4455
31 9625 2090
32 4731 2321
33 3578 2608
34 8504 1849
35 4027 1151
0 5647 4935
1 4219 1870
2 10968 8054
3 6970 5447
4 3217 5638
5 8972 669

TABLE 7-continued

Address of Parity Bit Accumulators (Rate 4/5)
6 5618 12472
7 1457 1280
8 8868 3883
9 8866 1224
10 8371 5972
11 266 4405
12 3706 3244
13 6039 5844
14 7200 3283
15 1502 11282
16 12318 2202
17 4523 965
18 9587 7011
19 2552 2051
20 12045 10306
21 11070 5104
22 6627 6906
23 9889 2121
24 829 9701
25 2201 1819
26 6689 12925
27 2139 8757
28 12004 5948
29 8704 3191
30 8171 10933
31 6297 7116
32 616 7146
33 5142 9761
34 10377 8138
35 7616 5811
0 7285 9863
1 7764 10867
2 12343 9019
3 4414 8331
4 3464 642
5 6960 2039
6 786 3021
7 710 2086
8 7423 5601
9 8120 4885
10 12385 11990
11 9739 10034
12 424 10162
13 1347 7597
14 1450 112
15 7965 8478
16 8945 7397
17 6590 8316
18 6838 9011
19 6174 9410
20 255 113
21 6197 5835
22 12902 3844
23 4377 3505
24 5478 8672
25 4453 2132
26 9724 1380
27 12131 11526
28 12323 9511
29 8231 1752
30 497 9022
31 9288 3080
32 2481 7515
33 2696 268
34 4023 12341
35 7108 5553

[0050]

TABLE 8

Address of Parity Bit Accumulators (Rate 3/5)

22422	10282	11626	19997	11161	2922	3122	99	5625	17064	8270	179
25087	16218	17015	828	20041	25656	4186	11629	22599	17305	22515	6463
11049	22853	25706	14388	5500	19245	8732	2177	13555	11346	17265	3069
16581	22225	12563	19717	23577	11555	25496	6853	25403	5218	15925	21766
16529	14487	7643	10715	17442	11119	5679	14155	24213	21000	1116	15620
5340	8636	16693	1434	5635	6516	9482	20189	1066	15013	25361	14243
18506	22236	20912	8952	5421	15691	6126	21595	500	6904	13059	6802
8433	4694	5524	14216	3685	19721	25420	9937	23813	9047	25651	16826
21500	24814	6344	17382	7064	13929	4004	16552	12818	8720	5286	2206
22517	2429	19065	2921	21611	1873	7507	5661	23006	23128	20543	19777
1770	4636	20900	14931	9247	12340	11008	12966	4471	2731	16445	791
6635	14556	18865	22421	22124	12697	9803	25485	7744	18254	11313	9004
19982	23963	18912	7206	12500	4382	20067	6177	21007	1195	23547	24837
756	11158	14646	20534	3647	17728	11676	11843	12937	4402	8261	22944
9306	24009	10012	11081	3746	24325	8060	19826	842	8836	2898	5019
7575	7455	25244	4736	14400	22981	5543	8006	24203	13053	1120	5128
3482	9270	13059	15825	7453	23747	3656	24585	16542	17507	22462	14670
15627	15290	4198	22748	5842	13395	23918	16985	14929	3726	25350	24157
24896	16365	16423	13461	16615	8107	24741	3604	25904	8716	9604	20365
3729	17245	18448	9862	20831	25326	20517	24618	13282	5099	14183	8804
16455	17646	15376	18194	25528	1777	6066	21855	14372	12517	4488	17490
1400	8135	23375	20879	8476	4084	12936	25536	22309	16582	6402	24360
25119	23586	128	4761	10443	22536	8607	9752	25446	15053	1856	4040
377	21160	13474	5451	17170	5938	10256	11972	24210	17833	22047	16108
13075	9648	24546	13150	23867	7309	19798	2988	16858	4825	23950	15125
20526	3553	11525	23366	2452	17626	19265	20172	18060	24593	13255	1552
18839	21132	20119	15214	14705	7096	10174	5663	18651	19700	12524	14033
4127	2971	17499	16287	22368	21463	7943	18880	5567	8047	23363	6797
10651	24471	14325	4081	7258	4949	7044	1078	797	22910	20474	4318
21374	13231	22985	5056	3821	23718	14178	9978	19030	23594	8895	25358
6199	22056	7749	13310	3999	23697	16445	22636	5225	22437	24153	9442
7978	12177	2893	20778	3175	8645	11863	24623	10311	25767	17057	3691
20473	11294	9914	22815	2574	8439	3699	5431	24840	21908	16088	18244
8208	5755	19059	8541	24924	6454	11234	10492	16406	10831	11436	9649
16264	11275	24953	2347	12667	19190	7257	7174	24819	2938	2522	11749
3627	5969	13862	1538	23176	6353	2855	17720	2472	7428	573	15036
0	18539	18661									
1	10502	3002									
2	9368	10761									
3	12299	7828									
4	15048	13362									
5	18444	24640									
6	20775	19175									
7	18970	10971									
8	5329	19982									
9	11296	18655									
10	15046	20659									
11	7300	22140									
12	22029	14477									
13	11129	742									
14	13254	13813									
15	19234	13273									
16	6079	21122									
17	22782	5828									
18	19775	4247									
19	1660	19413									
20	4403	3649									
21	13371	25851									
22	22770	21784									
23	10757	14131									
24	16071	21617									
25	6393	3725									
26	597	19968									
27	5743	8084									
28	6770	9548									
29	4285	17542									
30	13568	22599									
31	1786	4617									
32	23238	11648									
33	19627	2030									
34	13601	13458									
35	13740	17328									

TABLE 8-continued

Address of Parity Bit Accumulators (Rate 3/5)		
36	25012	13944
37	22513	6687
38	4934	12587
39	21197	5133
40	22705	6938
41	7534	24633
42	24400	12797
43	21911	25712
44	12039	1140
45	24306	1021
46	14012	20747
47	11265	15219
48	4670	15531
49	9417	14359
50	2415	6504
51	24964	24690
52	14443	8816
53	6926	1291
54	6209	20806
55	13915	4079
56	24410	13196
57	13505	6117
58	9869	8220
59	1570	6044
60	25780	17387
61	20671	24913
62	24558	20591
63	12402	3702
64	8314	1357
65	20071	14616
66	17014	3688
67	19837	946
68	15195	12136
69	7758	22808
70	3564	2925
71	3434	7769

[0051]

TABLE 9

Address of Parity Bit Accumulators (Rate 8/9)		
0	6235	2848 3222
1	5800	3492 5348
2	2757	927 90
3	6961	4516 4739
4	1172	3237 6264
5	1927	2425 3683
6	3714	6309 2495
7	3070	6342 7154
8	2428	613 3761
9	2906	264 5927
10	1716	1950 4273
11	4613	6179 3491
12	4865	3286 6005
13	1343	5923 3529
14	4589	4035 2132
15	1579	3920 6737
16	1644	1191 5998
17	1482	2381 4620
18	6791	6014 6596
19	2738	5918 3786
0	5156	6166
1	1504	4356
2	130	1904
3	6027	3187
4	6718	759
5	6240	2870
6	2343	1311
7	1039	5465

TABLE 9-continued

Address of Parity Bit Accumulators (Rate 8/9)		
8	6617	2513
9	1588	5222
10	6561	535
11	4765	2054
12	5966	6892
13	1969	3869
14	3571	2420
15	4632	981
16	3215	4163
17	973	3117
18	3802	6198
19	3794	3948
0	3196	6126
1	573	1909
2	850	4034
3	5622	1601
4	6005	524
5	5251	5783
6	172	2032
7	1875	2475
8	497	1291
9	2566	3430
10	1249	740
11	2944	1948
12	6528	2899
13	2243	3616
14	867	3733
15	1374	4702
16	4698	2285

TABLE 9-continued

Address of Parity Bit Accumulators (Rate 8/9)
17 4760 3917
18 1859 4058
19 6141 3527
0 2148 5066
1 1306 145
2 2319 871
3 3463 1061
4 5554 6647
5 5837 339
6 5821 4932
7 6356 4756
8 3930 418
9 211 3094
10 1007 4928
11 3584 1235
12 6982 2869
13 1612 1013
14 953 4964
15 4555 4410
16 4925 4842
17 5778 600
18 6509 2417
19 1260 4903
0 3369 3031
1 3557 3224
2 3028 583
3 3258 440
4 6226 6655
5 4895 1094
6 1481 6847
7 4433 1932
8 2107 1649
9 2119 2065
10 4003 6388
11 6720 3622
12 3694 4521
13 1164 7050
14 1965 3613
15 4331 66
16 2970 1796
17 4652 3218
18 1762 4777
19 5736 1399
0 970 2572
1 2062 6599
2 4597 4870
3 1228 6913
4 4159 1037
5 2916 2362
6 395 1226
7 6911 4548
8 4618 2241
9 4120 4280
10 5825 474
11 2154 5558
12 3793 5471
13 5707 1595
14 1403 325
15 6601 5183
16 6369 4569
17 4846 896
18 7092 6184
19 6764 7127
0 6358 1951
1 3117 6960
2 2710 7062
3 1133 3604
4 3694 657
5 1355 110
6 3329 6736
7 2505 3407
8 2462 4806
9 4216 214
10 5348 5619

TABLE 9-continued

Address of Parity Bit Accumulators (Rate 8/9)
11 6627 6243
12 2644 5073
13 4212 5088
14 3463 3889
15 5306 478
16 4320 6121
17 3961 1125
18 5699 1195
19 6511 792
0 3934 2778
1 3238 6587
2 1111 6596
3 1457 6226
4 1446 3885
5 3907 4043
6 6839 2873
7 1733 5615
8 5202 4269
9 3024 4722
10 5445 6372
11 370 1828
12 4695 1600
13 680 2074
14 1801 6690
15 2669 1377
16 2463 1681
17 5972 5171
18 5728 4284
19 1696 1459

[0052]

TABLE 10

Address of Parity Bit Accumulators (Rate 9/10)
0 5611 2563 2900
1 5220 3143 4813
2 2481 834 81
3 6265 4064 4265
4 1055 2914 5638
5 1734 2182 3315
6 3342 5678 2246
7 2185 552 3385
8 2615 236 5334
9 1546 1755 3846
10 4154 5561 3142
11 4382 2957 5400
12 1209 5329 3179
13 1421 3528 6063
14 1480 1072 5398
15 3843 1777 4369
16 1334 2145 4163
17 2368 5055 260
0 6118 5405
1 2994 4370
2 3405 1669
3 4640 5550
4 1354 3921
5 117 1713
6 5425 2866
7 6047 683
8 5616 2582
9 2108 1179
10 933 4921
11 5953 2261
12 1430 4699
13 5905 480
14 4289 1846
15 5374 6208
16 1775 3476

TABLE 10-continued

Address of Parity Bit Accumulators (Rate 9/10)
17 3216 2178
0 4165 884
1 2896 3744
2 874 2801
3 3423 5579
4 3404 3552
5 2876 5515
6 516 1719
7 765 3631
8 5059 1441
9 5629 598
10 5405 473
11 4724 5210
12 155 1832
13 1689 2229
14 449 1164
15 2308 3088
16 1122 669
17 2268 5758
0 5878 2609
1 782 3359
2 1231 4231
3 4225 2052
4 4286 3517
5 5531 3184
6 1935 4560
7 1174 131
8 3115 956
9 3129 1088
10 5238 4440
11 5722 4280
12 3540 375
13 191 2782
14 906 4432
15 3225 1111
16 6296 2583
17 1457 903
0 855 4475
1 4097 3970
2 4433 4361
3 5198 541
4 1146 4426
5 3202 2902
6 2724 525
7 1083 4124
8 2326 6003
9 5605 5990
10 4376 1579
11 4407 984
12 1332 6163
13 5359 3975
14 1907 1854
15 3601 5748
16 6056 3266
17 3322 4085
0 1768 3244
1 2149 144
2 1589 4291
3 5154 1252
4 1855 5939
5 4820 2706
6 1475 3360
7 4266 693
8 4156 2018
9 2103 752
10 3710 3853
11 5123 931
12 6146 3323
13 1939 5002
14 5140 1437
15 1263 293
16 5949 4665
17 4548 6380
0 3171 4690

TABLE 10-continued

Address of Parity Bit Accumulators (Rate 9/10)
1 5204 2114
2 6384 5565
3 5722 1757
4 2805 6264
5 1202 2616
6 1018 3244
7 4018 5289
8 2257 3067
9 2483 3073
10 1196 5329
11 649 3918
12 3791 4581
13 5028 3803
14 3119 3506
15 4779 431
16 3888 5510
17 4387 4084
0 5836 1692
1 5126 1078
2 5721 6165
3 3540 2499
4 2225 6348
5 1044 1484
6 6323 4042
7 1313 5603
8 1303 3496
9 3516 3639
10 5161 2293
11 4682 3845
12 3045 643
13 2818 2616
14 3267 649
15 6236 593
16 646 2948
17 4213 1442
0 5779 1596
1 2403 1237
2 2217 1514
3 5609 716
4 5155 3858
5 1517 1312
6 2554 3158
7 5280 2643
8 4990 1353
9 5648 1170
10 1152 4366
11 3561 5368
12 3581 1411
13 5647 4661
14 1542 5401
15 5078 2687
16 316 1755
17 3392 1991

[0053] As regards the BCH encoder 211, the BCH code parameters are enumerated in Table 11.

TABLE 11

LDPC Code Rate	BCH Uncoded Block		BCH Error Correction (bits)
	Length k_{bch}	BCH Coded Block Length n_{bch}	
1/2	32208	32400	12
2/3	43040	43200	10
3/4	48408	48600	12
4/5	51648	51840	12
5/6	53840	54000	10

TABLE 11-continued

LDPC Code Rate	BCH Uncoded Block Length k_{bch}	BCH Coded Block Length n_{bch}	BCH Error Correction (bits)
3/5	38688	38880	12
8/9	57472	57600	8
9/10	58192	58320	8

[0054] It is noted that in the above table, $n_{\text{bch}} = k_{\text{ldpc}}$.

[0055] The generator polynomial of the t error correcting BCH encoder 211 is obtained by multiplying the first t polynomials in the following list of Table 12:

TABLE 12

$g1(x)$	$1 + x^2 + x^3 + x^5 + x^{16}$
$g2(x)$	$1 + x + x^4 + x^5 + x^6 + x^8 + x^{16}$
$g3(x)$	$1 + x^2 + x^3 + x^4 + x^5 + x^7 + x^8 + x^9 + x^{10} + x^{11} + x^{16}$
$g4(x)$	$1 + x^2 + x^4 + x^6 + x^9 + x^{11} + x^{12} + x^{14} + x^{16}$
$g5(x)$	$1 + x + x^2 + x^3 + x^5 + x^8 + x^9 + x^{10} + x^{11} + x^{12} + x^{16}$
$g6(x)$	$1 + x^2 + x^4 + x^5 + x^7 + x^8 + x^9 + x^{10} + x^{12} + x^{13} + x^{14} + x^{15} + x^{16}$
$g7(x)$	$1 + x^2 + x^5 + x^6 + x^8 + x^9 + x^{10} + x^{11} + x^{13} + x^{15} + x^{16}$
$g8(x)$	$1 + x + x^2 + x^5 + x^6 + x^8 + x^9 + x^{12} + x^{13} + x^{14} + x^{16}$
$g9(x)$	$1 + x^5 + x^7 + x^9 + x^{10} + x^{11} + x^{16}$
$g10(x)$	$1 + x + x^2 + x^5 + x^7 + x^9 + x^{10} + x^{12} + x^{13} + x^{14} + x^{16}$
$g11(x)$	$1 + x^2 + x^3 + x^5 + x^9 + x^{11} + x^{12} + x^{13} + x^{16}$
$g12(x)$	$1 + x + x^5 + x^6 + x^7 + x^9 + x^{11} + x^{12} + x^{16}$

[0056] BCH encoding of information bits $m = (m_{k_{\text{bch}}-1}, m_{k_{\text{bch}}-2}, \dots, m_1, m_0)$ onto a codeword $c = (m_{k_{\text{bch}}-1}, m_{k_{\text{bch}}-2}, \dots, m_1, m_0, d_{n_{\text{bch}}-k_{\text{bch}}-1}, d_{n_{\text{bch}}-k_{\text{bch}}-2}, \dots, d_1, d_0)$ is achieved as follows. The message polynomial $m(x) = m_{k_{\text{bch}}-1}x^{k_{\text{bch}}-1} + m_{k_{\text{bch}}-2}x^{k_{\text{bch}}-2} + \dots + m_1x + m_0$ is multiplied by $x^{n_{\text{bch}}-k_{\text{bch}}}$. Next, $x^{n_{\text{bch}}-k_{\text{bch}}}m(x)$ divided by $g(x)$. With $d(x) = d_{n_{\text{bch}}-k_{\text{bch}}-1}x^{n_{\text{bch}}-k_{\text{bch}}-1} + \dots + d_1x + d_0$ as the remainder, the codeword polynomial is set as follows: $c(x) = x^{n_{\text{bch}}-k_{\text{bch}}}m(x) + d(x)$.

[0057] The above LDPC codes, in an exemplary embodiment, can be used to variety of digital video applications, such as MPEG (Motion Pictures Expert Group) packet transmission.

[0058] FIG. 3 is a diagram of an exemplary receiver in the system of FIG. 1. At the receiving side, a receiver 300 includes a demodulator 301 that performs demodulation of received signals from transmitter 200. These signals are received at a receive antenna 303 for demodulation. After demodulation, the received signals are forwarded to a decoder 305, which attempts to reconstruct the original source messages by generating messages, X' , in conjunction with a bit metric generator 307. With non-Gray mapping, the bit metric generator 307 exchanges probability information with the decoder 305 back and forth (iteratively) during the decoding process, which is detailed in FIG. 10. Alternatively, if Gray mapping is used (according to one embodiment of the present invention), one pass of the bit metric generator is sufficient, in which further attempts of bit metric generation after each LDPC decoder iteration are likely to yield limited performance improvement; this approach is more fully described with respect to FIG. 11. To appreciate the advantages offered by the present invention, it is instructive to examine how LDPC codes are generated, as discussed in FIG. 4.

[0059] FIG. 4 is a diagram of a sparse parity check matrix, in accordance with an embodiment of the present invention. LDPC codes are long, linear block codes with sparse parity check matrix $H_{(n-k) \times n}$. Typically the block length, n , ranges from thousands to tens of thousands of bits. For example, a parity check matrix for an LDPC code of length $n=8$ and rate $\frac{1}{2}$ is shown in FIG. 4. The same code can be equivalently represented by the bipartite graph, per FIG. 5.

[0060] FIG. 5 is a diagram of a bipartite graph of an LDPC code of the matrix of FIG. 4. Parity check equations imply that for each check node, the sum (over GF (Galois Field)(2)) of all adjacent bit nodes is equal to zero. As seen in the figure, bit nodes occupy the left side of the graph and are associated with one or more check nodes, according to a predetermined relationship. For example, corresponding to check node m_1 , the following expression exists $n_1 + n_4 + n_5 + n_8 = 0$ with respect to the bit nodes.

[0061] Returning the receiver 303, the LDPC decoder 305 is considered a message passing decoder, whereby the decoder 305 aims to find the values of bit nodes. To accomplish this task, bit nodes and check nodes iteratively communicate with each other. The nature of this communication is described below.

[0062] From check nodes to bit nodes, each check node provides to an adjacent bit node an estimate ("opinion") regarding the value of that bit node based on the information coming from other adjacent bit nodes. For instance, in the above example if the sum of n_4 , n_5 and n_8 "looks like" 0 to m_1 , then m_1 would indicate to n_1 that the value of n_1 is believed to be 0 (since $n_1 + n_4 + n_5 + n_8 = 0$); otherwise m_1 indicate to n_1 that the value of n_1 is believed to be 1. Additionally, for soft decision decoding, a reliability measure is added.

[0063] From bit nodes to check nodes, each bit node relays to an adjacent check node an estimate about its own value based on the feedback coming from its other adjacent check nodes. In the above example n_1 has only two adjacent check nodes m_1 and m_3 . If the feedback coming from m_3 to n_1 indicates that the value of n_1 is probably 0, then n_1 would notify m_1 that an estimate of n_1 's own value is 0. For the case in which the bit node has more than two adjacent check nodes, the bit node performs a majority vote (soft decision) on the feedback coming from its other adjacent check nodes before reporting that decision to the check node it communicates. The above process is repeated until all bit nodes are considered to be correct (i.e., all parity check equations are satisfied) or until a predetermined maximum number of iterations is reached, whereby a decoding failure is declared.

[0064] FIG. 6 is a diagram of a sub-matrix of a sparse parity check matrix, wherein the sub-matrix contains parity check values restricted to the lower triangular region, according to an embodiment of the present invention. As described previously, the encoder 203 (of FIGS. 2A and 2B) can employ a simple encoding technique by restricting the values of the lower triangular area of the parity check matrix. According to an embodiment of the present invention, the restriction imposed on the parity check matrix is of the form:

$$H_{(n-k) \times n} = [A_{(n-k) \times k} B_{(n-k) \times (n-k)}]$$

[0065] , where B is lower triangular.

[0066] Any information block $i=(i_0, i_1, \dots, i_{k-1})$ is encoded to a codeword $c=(i_0, i_1, \dots, i_{k-1}, p_0, p_1, \dots, p_{n-k-1})$ using $Hc^T=0$, and recursively solving for parity bits; for example,

$$a_{00}i_0 + a_{01}i_1 + \dots + a_{0,k-1}i_{k-1} + p_0 = 0 \Rightarrow \text{Solve } p_0$$

$$a_{10}i_0 + a_{11}i_1 + \dots + a_{1,k-1}i_{k-1} + b_{10}p_0 + p_1 = 0 \Rightarrow \text{Solve } p_1$$

[0067] and similarly for $p_2, p_3, \dots, p_{n-k-1}$.

[0068] FIG. 7 is a graph showing performance between codes utilizing unrestricted parity check matrix (H matrix) versus restricted H matrix of FIG. 6. The graph shows the performance comparison between two LDPC codes: one with a general parity check matrix and the other with a parity check matrix restricted to be lower triangular to simplify encoding. The modulation scheme, for this simulation, is 8-PSK. The performance loss is within 0.1 dB. Therefore, the performance loss is negligible based on the restriction of the lower triangular H matrices, while the gain in simplicity of the encoding technique is significant. Accordingly, any parity check matrix that is equivalent to a lower triangular or upper triangular under row and/or column permutation can be utilized for the same purpose.

[0069] FIGS. 8A and 8B are, respectively, a diagram of a non-Gray 8-PSK modulation scheme, and a Gray 8-PSK modulation, each of which can be used in the system of FIG. 1. The non-Gray 8-PSK scheme of FIG. 8A can be utilized in the receiver of FIG. 3 to provide a system that requires very low Frame Erasure Rate (FER). This requirement can also be satisfied by using a Gray 8-PSK scheme, as shown in FIG. 8B, in conjunction with an outer code, such as Bose, Chaudhuri, and Hocquenghem (BCH), Hamming, or Reed-Solomon (RS) code.

[0070] Under this scheme, there is no need to iterate between the LDPC decoder 305 (FIG. 3) and the bit metric generator 307, which may employ 8-PSK modulation. In the absence of an outer code, the LDPC decoder 305 using Gray labeling exhibit an earlier error floor, as shown in FIG. 9 below.

[0071] FIG. 9 is a graph showing performance between codes utilizing Gray labeling versus non-Gray labeling of FIGS. 8A and 8B. The error floor stems from the fact that assuming correct feedback from LDPC decoder 305, regeneration of 8-PSK bit metrics is more accurate with non-Gray labeling since the two 8-PSK symbols with known two bits are further apart with non-Gray labeling. This can be equivalently seen as operating at higher Signal-to-Noise Ratio (SNR). Therefore, even though error asymptotes of the same LDPC code using Gray or non-Gray labeling have the same slope (i.e., parallel to each other), the one with non-Gray labeling passes through lower FER at any SNR.

[0072] On the other hand, for systems that do not require very low FER, Gray labeling without any iteration between LDPC decoder 305 and 8-PSK bit metric generator 307 may be more suitable because re-generating 8-PSK bit metrics before every LDPC decoder iteration causes additional complexity. Moreover, when Gray labeling is used, re-generating 8-PSK bit metrics before every LDPC decoder iteration yields only very slight performance improvement. As mentioned previously, Gray labeling without iteration may be used for systems that require very low FER, provided an outer code is implemented.

[0073] The choice between Gray labeling and non-Gray labeling depends also on the characteristics of the LDPC code. Typically, the higher bit or check node degrees, the better it is for Gray labeling, because for higher node degrees, the initial feedback from LDPC decoder 305 to 8-PSK (or similar higher order modulation) bit metric generator 307 deteriorates more with non-Gray labeling.

[0074] When 8-PSK (or similar higher order) modulation is utilized with a binary decoder, it is recognized that the three (or more) bits of a symbol are not received "equally noisy". For example with Gray 8-PSK labeling, the third bit of a symbol is considered more noisy to the decoder than the other two bits. Therefore, the LDPC code design does not assign a small number of edges to those bit nodes represented by "more noisy" third bits of 8-PSK symbol so that those bits are not penalized twice.

[0075] FIG. 10 is a flow chart of the operation of the LDPC decoder using non-Gray mapping, according to an embodiment of the present invention. Under this approach, the LDPC decoder and bit metric generator iterate one after the other. In this example, 8-PSK modulation is utilized; however, the same principles apply to other higher modulation schemes as well. Under this scenario, it is assumed that the demodulator 301 outputs a distance vector, d , denoting the distances between received noisy symbol points and 8-PSK symbol points to the bit metric generator 307, whereby the vector components are as follows:

$$d_i = \frac{E_s}{N_0} \{ (r_x - s_{i,x})^2 + (r_y - s_{i,y})^2 \} \quad i = 0, 1, \dots, 7.$$

[0076] The 8-PSK bit metric generator 307 communicates with the LDPC decoder 305 to exchange a priori probability information and a posteriori probability information, which respectively are represented as u , and a . That is, the vectors u and a respectively represent a priori and a posteriori probabilities of log likelihood ratios of coded bits.

[0077] The 8-PSK bit metric generator 307 generates the a priori likelihood ratios for each group of three bits as follows. First, extrinsic information on coded bits is obtained:

$$e_j = a_j - u_j, j=0, 1, 2.$$

[0078] Next, 8-PSK symbol probabilities, p_i $i=0, 1, \dots, 7$, are determined.

$$^*y_j = -f(0, e_j) \quad j=0, 1, 2 \text{ where } f(a, b) = \max(a, b) + LUT_f(a, b)$$

[0079] with

$$LUT_f(a, b) = \ln(1 + e^{-|a-b|})$$

[0080]

$^*x_j = y_j + e_j$	$j = 0, 1, 2$
$^*p_0 = x_0 + x_1 + x_2$	$p_4 = y_0 + x_1 + x_2$
$p_1 = x_0 + x_1 + y_2$	$p_5 = y_0 + x_1 + y_2$
$p_2 = x_0 + y_1 + x_2$	$p_6 = y_0 + y_1 + x_2$
$p_3 = x_0 + y_1 + y_2$	$p_7 = y_0 + y_1 + y_2$

[0081] Next, the bit metric generator 307 determines a priori log likelihood ratios of the coded bits as input to LDPC decoder 305, as follows:

$$\begin{aligned} u_0 &= f(d_0 + p_0, d_1 + p_1, d_2 + p_2, d_3 + p_3) - f(d_4 + p_4, d_5 + p_5, d_6 + p_6, \\ &\quad d_7 + p_7) - e_0 \\ u_1 &= f(d_0 + p_0, d_1 + p_1, d_4 + p_4, d_5 + p_5) - f(d_2 + p_2, d_3 + p_3, d_6 + p_6, \\ &\quad d_7 + p_7) - e_1 \\ u_2 &= f(d_0 + p_0, d_2 + p_2, d_4 + p_4, d_5 + p_5) - f(d_1 + p_1, d_3 + p_3, d_6 + p_6, \\ &\quad d_7 + p_7) - e_2 \end{aligned}$$

[0082] It is noted that the function $f(\cdot)$ with more than two variables can be evaluated recursively; e.g. $f(a, b, c) = f(f(a, b), c)$.

[0083] The operation of the LDPC decoder 305 utilizing non-Gray mapping is now described. In step 1001, the LDPC decoder 305 initializes log likelihood ratios of coded bits, v , before the first iteration according to the following (and as shown in FIG. 12A):

$$\begin{aligned} v_{n \rightarrow k_i} &= u_n, \\ n &= 0, 1, \dots, N-1, \\ i &= 1, 2, \dots, \text{deg}(\text{bit node } n) \end{aligned}$$

[0084] Here, $v_{n \rightarrow k_i}$ denotes the message that goes from bit node n to its adjacent check node k_i , u_n denotes the demodulator output for the bit n and N is the codeword size.

[0085] In step 1003, a check node, k , is updated, whereby the input v yields the output w . As seen in FIG. 12B, the incoming messages to the check node k from its d_c adjacent bit nodes are denoted by $v_{n_1 \rightarrow k}, v_{n_2 \rightarrow k}, \dots, v_{n_{d_c} \rightarrow k}$. The goal is to compute the outgoing messages from the check node k back to d_c adjacent bit nodes. These messages are denoted by

$$w_{k \rightarrow n_1}, w_{k \rightarrow n_2}, \dots, w_{k \rightarrow n_{d_c}},$$

[0086] where

$$w_{k \rightarrow n_1} = g(v_{n_1 \rightarrow k}, v_{n_2 \rightarrow k}, \dots, v_{n_{d_c-1} \rightarrow k}, v_{n_{d_c} \rightarrow k}, k).$$

[0087] The function $g(\cdot)$ is defined as follows:

$$g(a, b) = \text{sign}(a) \times \text{sign}(b) \times \{\min(|a|, |b|)\} + \text{LUT}_g(a, b),$$

[0088] where $\text{LUT}_g(a, b) = \ln(1 + e^{-|a+b|}) - \ln(1 + e^{-|a-b|})$. Similar to function f function g with more than two variables can be evaluated recursively.

[0089] Next, the decoder 305, per step 1205, outputs a posteriori probability information (FIG. 12C), such that:

$$a_n = u_n + \sum_j w_{k_j \rightarrow n}.$$

[0090] Per step 1007, it is determined whether all the parity check equations are satisfied. If these parity check equations are not satisfied, then the decoder 305, as in step 1009, re-derives 8-PSK bit metrics and channel input u_n . Next, the bit node is updated, as in step 1011. As shown in FIG. 14C, the incoming messages to the bit node n from its d_v adjacent check nodes are denoted by $w_{k_1 \rightarrow n}, w_{k_2 \rightarrow n}, \dots, w_{k_{d_v} \rightarrow n}$. The outgoing messages from the bit node n are computed back to d_v adjacent check nodes; such messages are denoted by $v_{n \rightarrow k_1}, v_{n \rightarrow k_2}, \dots, v_{n \rightarrow k_{d_v}}$, and computed as follows:

$$v_{n \rightarrow k_i} = u_n + \sum_{j \neq i} w_{k_j \rightarrow n}$$

[0091] In step 1013, the decoder 305 outputs the hard decision (in the case that all parity check equations are satisfied):

$$\hat{c}_n = \begin{cases} 0, & a_n \geq 0 \\ 1, & a_n < 0 \end{cases} \quad \text{Stop if } H\hat{c}^T = 0$$

[0092] The above approach is appropriate when non-Gray labeling is utilized. However, when Gray labeling is implemented, the process of FIG. 11 is executed.

[0093] FIG. 11 is a flow chart of the operation of the LDPC decoder of FIG. 3 using Gray mapping, according to an embodiment of the present invention. When Gray labeling is used, bit metrics are advantageously generated only once before the LDPC decoder, as re-generating bit metrics after every LDPC decoder iteration may yield nominal performance improvement. As with steps 1001 and 1003 of FIG. 10, initialization of the log likelihood ratios of coded bits, v , are performed, and the check node is updated, per steps 1101 and 1103. Next, the bit node n is updated, as in step 1105. Thereafter, the decoder outputs the a posteriori probability information (step 1107). In step 1109, a determination is made whether all of the parity check equations are satisfied; if so, the decoder outputs the hard decision (step 1111). Otherwise, steps 1103-1107 are repeated.

[0094] FIG. 13A is a flowchart of process for computing outgoing messages between the check nodes and the bit nodes using a forward-backward approach, according to an embodiment of the present invention. For a check node with d_c adjacent edges, the computation of $d_c(d_c-1)$ and numerous $g(\cdot, \cdot)$ functions are performed. However, the forward-backward approach reduces the complexity of the computation to $3(d_c-2)$, in which d_c-1 variables are stored.

[0095] Referring to FIG. 12B, the incoming messages to the check node k from d_c adjacent bit nodes are denoted by $v_{n_1 \rightarrow k}, v_{n_2 \rightarrow k}, \dots, v_{n_{d_c} \rightarrow k}$. It is desired that the outgoing messages are computed from the check node k back to d_c adjacent bit nodes; these outgoing messages are denoted by $w_{k \rightarrow n_1}, w_{k \rightarrow n_2}, \dots, w_{k \rightarrow n_{d_c}}$.

[0096] Under the forward-backward approach to computing these outgoing messages, forward variables, $f_1, f_2, \dots, f_{d_c}$, are defined as follows:

$$\begin{aligned} f_1 &= v_{1 \rightarrow k} \\ f_2 &= g(f_1, v_{2 \rightarrow k}) \\ f_3 &= g(f_2, v_{3 \rightarrow k}) \\ f_{d_c} &= g(f_{d_c-1}, v_{d_c \rightarrow k}) \end{aligned}$$

[0097] In step 1301, these forward variables are computed, and stored, per step 1303.

[0098] Similarly, backward variables, $b_1, b_2, \dots, b_{dc}$, are defined by the following:

$$\begin{aligned} b_{dc} &= v_{dc \rightarrow k} \\ b_{dc-1} &= g(b_{dc}, v_{dc-1 \rightarrow k}) \\ b_1 &= g(b_2, v_{1 \rightarrow k}) \end{aligned}$$

[0099] In step 1305, these backward variables are then computed. Thereafter, the outgoing messages are computed, as in step 1307, based on the stored forward variables and the computed backward variables. The outgoing messages are computed as follows:

$$\begin{aligned} w_{k \rightarrow 1} &= b_2 \\ w_{k \rightarrow i} &= g(f_{i-1}, b_{i+1}) \quad i=2, 3, \dots, dc-1 \\ w_{k \rightarrow dc} &= f_{dc-1} \end{aligned}$$

[0100] Under this approach, only the forward variables, $f_2, f_3, \dots, f_{dc}$, are required to be stored. As the backward variables b_i are computed, the outgoing messages, $w_{k \rightarrow i}$, are simultaneously computed, thereby negating the need for storage of the backward variables.

[0101] The computation load can be further enhanced by a parallel approach, as next discussed.

[0102] FIG. 13B is a flowchart of process for computing outgoing messages between the check nodes and the bit nodes using a parallel approach, according to an embodiment of the present invention. For a check node k with inputs $v_{n_1 \rightarrow k}, v_{n_2 \rightarrow k}, \dots, v_{n_{dc} \rightarrow k}$ from d_c adjacent bit nodes, the following parameter is computed, as in step 1311:

$$\gamma_k = g(v_{n_1 \rightarrow k}, v_{n_2 \rightarrow k}, \dots, v_{n_{dc} \rightarrow k}).$$

[0103] It is noted that the $g(\dots)$ function can also be expressed as follows:

$$g(a, b) = \ln \frac{1 + e^{a+b}}{e^a + e^b}.$$

[0104] Exploiting the recursive nature of the $g(\dots)$ function, the following expression results:

$$\gamma_k = \ln \frac{1 + e^{g(v_{n_1 \rightarrow k}, \dots, v_{n_{i-1} \rightarrow k}, v_{n_{i+1} \rightarrow k}, \dots, v_{n_{dc} \rightarrow k}) + v_{n_i \rightarrow k}}}{e^{g(v_{n_1 \rightarrow k}, \dots, v_{n_{i-1} \rightarrow k}, v_{n_{i+1} \rightarrow k}, \dots, v_{n_{dc} \rightarrow k})} + e^{v_{n_i \rightarrow k}}} = \ln \frac{1 + e^{w_{k \rightarrow n_i} + v_{n_i \rightarrow k}}}{e^{w_{k \rightarrow n_i}} + e^{v_{n_i \rightarrow k}}}.$$

[0105] Accordingly, $w_{k \rightarrow n_i}$ can be solved in the following manner:

$$w_{k \rightarrow n_i} = \ln \frac{e^{v_{n_i \rightarrow k} + \gamma_k} - 1}{e^{v_{n_i \rightarrow k}} - 1} - \gamma_k$$

[0106] The $\ln(\dots)$ term of the above equation can be obtained using a look-up table LUT_x that represents the function $\ln[e^x - 1]$ (step 1313). Unlike the other look-up tables LUT_f or LUT_g , the table LUT_x would likely requires as many entries as the number of quantization levels. Once γ_k is obtained, the calculation of $w_{k \rightarrow n_i}$ for all n_i can occur in parallel using the above equation, per step 1315.

[0107] The computational latency of γ_k is advantageously $\log_2(d_c)$.

[0108] FIGS. 14A-14C are graphs showing simulation results of LDPC codes generated in accordance with various embodiments of the present invention. In particular, FIGS. 14A-14C show the performance of LDPC codes with higher order modulation and code rates of $\frac{3}{4}$ (QPSK, 1.485 bits/symbol), $\frac{2}{3}$ (8-PSK, 1.980 bits/symbol), and $\frac{5}{6}$ (8-PSK, 2.474 bits/symbol).

[0109] Two general approaches exist to realize the interconnections between check nodes and bit nodes: (1) a fully parallel approach, and (2) a partially parallel approach. In fully parallel architecture, all of the nodes and their interconnections are physically implemented. The advantage of this architecture is speed.

[0110] The fully parallel architecture, however, may involve greater complexity in realizing all of the nodes and their connections. Therefore with fully parallel architecture, a smaller block size may be required to reduce the complexity. In that case, for the same clock frequency, a proportional reduction in throughput and some degradation in FER versus E_s/N_0 performance may result.

[0111] The second approach to implementing LDPC codes is to physically realize only a subset of the total number of the nodes and use only these limited number of "physical" nodes to process all of the "functional" nodes of the code. Even though the LDPC decoder operations can be made extremely simple and can be performed in parallel, the further challenge in the design is how the communication is established between "randomly" distributed bit nodes and check nodes. The decoder 305 (of FIG. 3), according to one embodiment of the present invention, addresses this problem by accessing memory in a structured way, as to realize a seemingly random code. This approach is explained with respect to FIGS. 15A and 15B.

[0112] FIGS. 15A and 15B are diagrams of the top edge and bottom edge, respectively, of memory organized to support structured access as to realize randomness in LDPC coding, according to an embodiment of the present invention. Structured access can be achieved without compromising the performance of a truly random code by focusing on the generation of the parity check matrix. In general, a parity check matrix can be specified by the connections of the check nodes with the bit nodes. For example, the bit nodes can be divided into groups of a fixed size, which for illustrative purposes is 392. Additionally, assuming the check nodes connected to the first bit node of degree 3, for instance, are numbered as a, b and c , then the check nodes connected to the second bit node are numbered as $a+p, b+p$ and $c+p$, the check nodes connected to the third bit node are numbered as $a+2p, b+2p$ and $c+2p$ etc.; where p =(number of check nodes)/392. For the next group of 392 bit nodes, the check nodes connected to the first bit node are different from a, b, c so that with a suitable choice of p , all the check nodes have the same degree. A random search is performed over the free constants such that the resulting LDPC code is cycle-4 and cycle-6 free. Because of the structural characteristics of the parity check matrix of the present invention, the edge information can be stored to permit concurrent access to a group of relevant edge values during decoding.

[0113] In other words, the approach of the present invention facilitates memory access during check node and bit

node processing. The values of the edges in the bipartite graph can be stored in a storage medium, such as random access memory (RAM). It is noted that for a truly random LDPC code during check node and bit node processing, the values of the edges would need to be accessed one by one in a random fashion. However, such a conventional access scheme would be too slow for a high data rate application. The RAM of FIGS. 15A and 15B are organized in a manner, whereby a large group of relevant edges can be fetched in one clock cycle; accordingly, these values are placed "together" in memory, according to a predetermined scheme or arrangement. It is observed that, in actuality, even with a truly random code, for a group of check nodes (and respectively bit nodes), the relevant edges can be placed next to one another in RAM, but then the relevant edges adjacent to a group of bit nodes (respectively check nodes) will be randomly scattered in RAM. Therefore, the "togetherness," under the present invention, stems from the design of the parity check matrices themselves. That is, the check matrix design ensures that the relevant edges for a group of bit nodes and check nodes are simultaneously placed together in RAM.

[0114] As seen in FIGS. 15A and 15B, each box contains the value of an edge, which is multiple bits (e.g., 6). Edge RAM, according to one embodiment of the present invention, is divided into two parts: top edge RAM 1501 (FIG. 15A) and bottom edge RAM 1503 (FIG. 15B). Bottom edge RAM 1503 contains the edges between bit nodes of degree 2, for example, and check nodes. Top edge RAM contains the edges between bit nodes of degree greater than 2 and check nodes. Therefore, for every check node, 2 adjacent edges are stored in the bottom RAM 1503, and the rest of the edges are stored in the top edge RAM 1501. For example, the size of the top edge RAM 1501 and bottom edge RAM 1503 for various code rates are given in Table 14:

TABLE 14

	1/2	2/3	3/4	5/6
TopEdge RAM	400 × 392	440 × 392	504 × 392	520 × 392
Bottom Edge RAM	160 × 392	110 × 392	72 × 392	52 × 392

[0115] Based on Table 14, an edge RAM of size 576×392 is sufficient to store the edge metrics for all the code rates of 1/2, 2/3, 3/4, and 5/6.

[0116] As noted, under this exemplary scenario, a group of 392 bit nodes and 392 check nodes are selected for processing at a time. For 392 check node processing, $q=d_c-2$ consecutive rows are accessed from the top edge RAM 1501, and 2 consecutive rows from the bottom edge RAM 1503. The value of d_c depends on the specific code, for example $d_c=7$ for rate 1/2, $d_c=10$ for rate 2/3, $d_c=16$ for rate 3/4 and $d_c=22$ for rate 5/6 for the above course other values of d_c for other codes are possible. In this instance, $q+2$ is the degree of each check node.

[0117] For bit node processing, if the group of 392 bit nodes has degree 2, their edges are located in 2 consecutive rows of the bottom edge RAM 1503. If the bit nodes have degree $d>2$, their edges are located in some d rows of the top edge RAM 1501. The address of these d rows can be stored

in non-volatile memory, such as Read-Only Memory (ROM). The edges in one of the rows correspond to the first edges of 392 bit nodes, the edges in another row correspond to the second edges of 392 bit nodes, etc. Moreover for each row, the column index of the edge that belongs to the first bit node in the group of 392 can also be stored in ROM. The edges that correspond to the second, third, etc. bit nodes follow the starting column index in a "wrapped around" fashion. For example, if the j^{th} edge in the row belongs to the first bit node, then the $(j+1)^{\text{st}}$ edge belongs to the second bit node, $(j+2)^{\text{nd}}$ edge belongs to the third bit node, . . . , and $(j-1)^{\text{st}}$ edge belongs to the 392th bit node.

[0118] With the organization shown in FIGS. 15A and 15B, speed of memory access is greatly enhanced during LDPC coding.

[0119] FIG. 16 illustrates a computer system upon which an embodiment according to the present invention can be implemented. The computer system 1600 includes a bus 1601 or other communication mechanism for communicating information, and a processor 1603 coupled to the bus 1601 for processing information. The computer system 1600 also includes main memory 1605, such as a random access memory (RAM) or other dynamic storage device, coupled to the bus 1601 for storing information and instructions to be executed by the processor 1603. Main memory 1605 can also be used for storing temporary variables or other intermediate information during execution of instructions to be executed by the processor 1603. The computer system 1600 further includes a read only memory (ROM) 1607 or other static storage device coupled to the bus 1601 for storing static information and instructions for the processor 1603. A storage device 1609, such as a magnetic disk or optical disk, is additionally coupled to the bus 1601 for storing information and instructions.

[0120] The computer system 1600 may be coupled via the bus 1601 to a display 1611, such as a cathode ray tube (CRT), liquid crystal display, active matrix display, or plasma display, for displaying information to a computer user. An input device 1613, such as a keyboard including alphanumeric and other keys, is coupled to the bus 1601 for communicating information and command selections to the processor 1603. Another type of user input device is cursor control 1615, such as a mouse, a trackball, or cursor direction keys for communicating direction information and command selections to the processor 1603 and for controlling cursor movement on the display 1611.

[0121] According to one embodiment of the invention, generation of LDPC codes is provided by the computer system 1600 in response to the processor 1603 executing an arrangement of instructions contained in main memory 1605. Such instructions can be read into main memory 1605 from another computer-readable medium, such as the storage device 1609. Execution of the arrangement of instructions contained in main memory 1605 causes the processor 1603 to perform the process steps described herein. One or more processors in a multi-processing arrangement may also be employed to execute the instructions contained in main memory 1605. In alternative embodiments, hard-wired circuitry may be used in place of or in combination with software instructions to implement the embodiment of the present invention. Thus, embodiments of the present invention are not limited to any specific combination of hardware circuitry and software.

[0122] The computer system 1600 also includes a communication interface 1617 coupled to bus 1601. The communication interface 1617 provides a two-way data communication coupling to a network link 1619 connected to a local network 1621. For example, the communication interface 1617 may be a digital subscriber line (DSL) card or modem, an integrated services digital network (ISDN) card, a cable modem, or a telephone modem to provide a data communication connection to a corresponding type of telephone line. As another example, communication interface 1617 may be a local area network (LAN) card (e.g. for Ethernet™ or an Asynchronous Transfer Model (ATM) network) to provide a data communication connection to a compatible LAN. Wireless links can also be implemented. In any such implementation, communication interface 1617 sends and receives electrical, electromagnetic, or optical signals that carry digital data streams representing various types of information. Further, the communication interface 1617 can include peripheral interface devices, such as a Universal Serial Bus (USB) interface, a PCMCIA (Personal Computer Memory Card International Association) interface, etc.

[0123] The network link 1619 typically provides data communication through one or more networks to other data devices. For example, the network link 1619 may provide a connection through local network 1621 to a host computer 1623, which has connectivity to a network 1625 (e.g. a wide area network (WAN) or the global packet data communication network now commonly referred to as the "Internet") or to data equipment operated by service provider. The local network 1621 and network 1625 both use electrical, electromagnetic, or optical signals to convey information and instructions. The signals through the various networks and the signals on network link 1619 and through communication interface 1617, which communicate digital data with computer system 1600, are exemplary forms of carrier waves bearing the information and instructions.

[0124] The computer system 1600 can send messages and receive data, including program code, through the network(s), network link 1619, and communication interface 1617. In the Internet example, a server (not shown) might transmit requested code belonging to an application program for implementing an embodiment of the present invention through the network 1625, local network 1621 and communication interface 1617. The processor 1603 may execute the transmitted code while being received and/or store the code in storage device 1609, or other non-volatile storage for later execution. In this manner, computer system 1600 may obtain application code in the form of a carrier wave.

[0125] The term "computer-readable medium" as used herein refers to any medium that participates in providing instructions to the processor 1603 for execution. Such a medium may take many forms, including but not limited to non-volatile media, volatile media, and transmission media. Non-volatile media include, for example, optical or magnetic disks, such as storage device 1609. Volatile media include dynamic memory, such as main memory 1605. Transmission media include coaxial cables, copper wire and fiber optics, including the wires that comprise bus 1601. Transmission media can also take the form of acoustic, optical, or electromagnetic waves, such as those generated during radio frequency (RF) and infrared (IR) data communications. Common forms of computer-readable media

include, for example, a floppy disk, a flexible disk, hard disk, magnetic tape, any other magnetic medium, a CD-ROM, CDRW, DVD, any other optical medium, punch cards, paper tape, optical mark sheets, any other physical medium with patterns of holes or other optically recognizable indicia, a RAM, a PROM, and EPROM, a FLASH-EPROM, any other memory chip or cartridge, a carrier wave, or any other medium from which a computer can read.

[0126] Various forms of computer-readable media may be involved in providing instructions to a processor for execution. For example, the instructions for carrying out at least part of the present invention may initially be borne on a magnetic disk of a remote computer. In such a scenario, the remote computer loads the instructions into main memory and sends the instructions over a telephone line using a modem. A modem of a local computer system receives the data on the telephone line and uses an infrared transmitter to convert the data to an infrared signal and transmit the infrared signal to a portable computing device, such as a personal digital assistance (PDA) and a laptop. An infrared detector on the portable computing device receives the information and instructions borne by the infrared signal and places the data on a bus. The bus conveys the data to main memory, from which a processor retrieves and executes the instructions. The instructions received by main memory may optionally be stored on storage device either before or after execution by processor.

[0127] Accordingly, the various embodiments of the present invention provide an approach for encoding structured Low Density Parity Check (LDPC) codes. Structure of the LDPC codes is provided by restricting portion part of the parity check matrix to be lower triangular and/or satisfying other requirements such that the communication between bit nodes and check nodes of the decoder is simplified. Memory storing information representing the structured parity check matrix is accessed. The information is organized in tabular form, wherein each row represents occurrences of one values within a first column of a group of columns of the parity check matrix. The rows correspond to groups of columns of the parity check matrix, wherein subsequent columns within each of the groups are derived according to a predetermined operation (e.g., cyclic shift, addition, etc.). An LDPC coded signal based on the stored information representing the parity check matrix. According to one embodiment of the present invention, a Bose Chaudhuri Hocquenghem (BCH) encoder is utilized by the transmitter to encode an input signal using BCH codes, wherein the output LDPC coded signal corresponding to the input signal represents a code having an outer BCH code and an inner LDPC code. Further, a cyclic redundancy check (CRC) encoder is supplied to encode the input signal according to a CRC code. The above approach advantageously yields reduced complexity without sacrificing performance.

[0128] While the present invention has been described in connection with a number of embodiments and implementations, the present invention is not so limited but covers various obvious modifications and equivalent arrangements, which fall within the purview of the appended claims.

What is claimed is:

1. A method of encoding, comprising:

accessing memory storing information representing a structured parity check matrix of Low Density Parity

Check (LDPC) codes, the information being organized in tabular form, wherein each row represents occurrences of one values within a first column of a group of columns of the parity check matrix, the rows correspond to groups of columns of the parity check matrix, wherein subsequent columns within each of the groups are derived according to a predetermined operation; and

outputting an LDPC coded signal based on the stored information representing the parity check matrix.

2. A method according to claim 1, wherein the predetermined operation specifies one of the steps of:

performing a cyclic shift on the first column of each of the group; and

adding a constant to the first column of each of the group, the constant being dependent on code rate of the LDPC code.

3. A method according to claim 1, wherein the parity bits are determined sequentially, the method further comprising:

determining an i^{th} parity bit by adding the $(i-1)^{\text{th}}$ parity bit and the j^{th} information bit if the j^{th} entry in the i^{th} row of the parity check matrix is 1.

4. A method according to claim 1, further comprising:

initializing parity bit accumulators zero;

accumulating the first information bit in the j^{th} group of M information bits in the i^{th} parity bit accumulator if the i^{th} entry in $(jM)^{\text{th}}$ column of the parity check matrix is 1, where $j=0,1,2,3, \dots, k_{\text{ldpc}}/M-1$;

accumulating the remaining $(M-1)$ information bits $m=jM+1, jM+2, jM+3, \dots, (j+1)M-1$ of the j^{th} group in the parity bit accumulators according to

$\{x \bmod M \times q\} \bmod (n_{\text{ldpc}} - k_{\text{ldpc}})$, wherein x denotes the address of the parity bit accumulator corresponding to the first bit, jM , in the group, and q is a code rate dependent constant; and

after all of the information bits are exhausted, performing operations, starting with $i=1$ according to $p_i = p_i \oplus p_{i-1}$,

$i=1,2, \dots, n_{\text{ldpc}} - k_{\text{ldpc}} - 1$, wherein final content of p_i , $i=0,1, \dots, n_{\text{ldpc}} - k_{\text{ldpc}} - 1$ is equal to the parity bit p_i .

5. A method according to claim 4, wherein $M=360$.

6. A method according to claim 4, wherein the code dependent constant q is 60, 30, 90, 45, 36, 72, 20, and 18 for code rates $\frac{2}{3}$, $\frac{5}{6}$, $\frac{1}{2}$, $\frac{3}{4}$, $\frac{4}{5}$, $\frac{3}{5}$, $\frac{8}{9}$, and $\frac{9}{10}$, respectively.

7. A method according to claim 1, further comprising:

modulating the LDPC coded signal according to a signal constellation that includes one of 8-PSK (Phase Shift Keying), 16-QAM (Quadrature Amplitude Modulation), QPSK (Quadrature Phase Shift Keying), 16-APSK (Amplitude Phase Shift Keying) and 32-APSK.

8. A method according to claim 1, further comprising:

encoding an input signal according to Bose Chaudhuri Hocquenghem (BCH) codes, wherein the output LDPC coded signal corresponding to the input signal represents a code having an outer BCH code and an inner LDPC code.

9. A method according to claim 8, wherein the number of redundant BCH bits is $n_{\text{BCH}} - k_{\text{BCH}} = 16 \times t$, wherein t represents error correcting capability of the BCH code.

10. A method according to claim 8, wherein the error correction capability of the BCH code is 12 bits when used in concatenation with rate $\frac{1}{2}$, $\frac{3}{4}$, $\frac{4}{5}$ and $\frac{3}{5}$ LDPC codes, is 10 bits when used in concatenation with rate $\frac{2}{3}$ and $\frac{5}{6}$ LDPC codes, and is 8 bits when used in concatenation with rate $\frac{8}{9}$ and $\frac{9}{10}$ LDPC codes.

11. A method according to claim 1, wherein the row indices of 1's in the column index $j \times 360 (j=0,1,2,3, \dots,$

$$(j=0, 1, 2, 3, \dots, \frac{k_{\text{ldpc}}}{360} - 1)$$

of the parity check matrix are given at the j^{th} row according to one of Tables 1-8:

TABLE 1

Address of Parity Bit Accumulators (Rate 2/3)

0	10491	16043	506	12826	8065	8226	2767	240	18673	9279	10579	20928
1	17819	8313	6433	6224	5120	5824	12812	17187	9940	13447	13825	18483
2	17957	6024	8681	18628	12794	5915	14576	10970	12064	20437	4455	7151
3	19777	6183	9972	14536	8182	17749	11341	5556	4379	17434	15477	18532
4	4651	19689	1608	659	16707	14335	6143	3058	14618	17894	20684	5306
5	9778	2552	12096	12369	15198	16890	4851	3109	1700	18725	1997	15882
6	486	6111	13743	11537	5591	7433	15227	14145	1483	3887	17431	12430
7	20647	14311	11734	4180	8110	5525	12141	15761	18661	18441	10569	8192
8	3791	14759	15264	19918	10132	9062	10010	12786	10675	9682	19246	5454
9	19525	9485	7777	19999	8378	9209	3163	20232	6690	16518	716	7353
10	4588	6709	20202	10905	915	4317	11073	13576	16433	368	3508	21171
11	14072	4033	19959	12608	631	19494	14160	8249	10223	21504	12395	4322
12	13800	14161										
13	2948	9647										
14	14693	16027										
15	20506	11082										
16	1143	9020										
17	13501	4014										
18	1548	2190										
19	12216	21556										
20	2095	19897										

TABLE 1-continued

Address of Parity Bit Accumulators (Rate 2/3)	
21 4189 7958	
22 15940 10048	
23 515 12614	
24 8501 8450	
25 17595 16784	
26 5913 8495	
27 16394 10423	
28 7409 6981	
29 6678 15939	
30 20344 12987	
31 2510 14588	
32 17918 6655	
33 6703 19451	
34 496 4217	
35 7290 5766	
36 10521 8925	
37 20379 11905	
38 4090 5838	
39 19082 17040	
40 20233 12352	
41 19365 19546	
42 6249 19030	
43 11037 19193	
44 19760 11772	
45 19644 7428	
46 16076 3521	
47 11779 21062	
48 13062 9682	
49 8934 5217	
50 11087 3319	
51 18892 4356	
52 7894 3898	
53 5963 4360	
54 7346 11726	
55 5182 5609	
56 2412 17295	
57 9845 20494	
58 6687 1864	
59 20564 5216	
0 18226 17207	
1 9380 8266	
2 7073 3065	
3 18252 13437	
4 9161 15642	
5 10714 10153	
6 11585 9078	
7 5359 9418	
8 9024 9515	
9 1206 16354	
10 14994 1102	
11 9375 20796	
12 15964 6027	
13 14789 6452	
14 8002 18591	
15 14742 14089	
16 253 3045	
17 1274 19286	
18 14777 2044	
19 13920 9900	
20 452 7374	
21 18206 9921	
22 6131 5414	
23 10077 9726	
24 12045 5479	
25 4322 7990	
26 15616 5550	
27 15561 10661	
28 20718 7387	
29 2518 18804	
30 8984 2600	
31 6516 17909	
32 11148 98	
33 20559 3704	
34 7510 1569	

TABLE 1-continued

Address of Parity Bit Accumulators (Rate 2/3)	
35	16000 11692
36	9147 10303
37	16650 191
38	15577 18685
39	17167 20917
40	4256 3391
41	20092 17219
42	9218 5056
43	18429 8472
44	12093 20753
45	16345 12748
46	16023 11095
47	5048 17595
48	18995 4817
49	16483 3536
50	1439 16148
51	3661 3039
52	19010 18121
53	8968 11793
54	13427 18003
55	5303 3083
56	531 16668
57	4771 6722
58	5695 7960
59	3589 14630

TABLE 2

Address of Parity Bit Accumulators (Rate 5/6)	
0	4362 416 8909 4156 3216 3112 2560 2912 6405 8593 4969 6723
1	2479 1786 8978 3011 4339 9313 6397 2957 7288 5484 6031 10217
2	10175 9009 9889 3091 4985 7267 4092 8874 5671 2777 2189 8716
3	9052 4795 3924 3370 10058 1128 9996 10165 9360 4297 434 5138
4	2379 7834 4835 2327 9843 804 329 8353 7167 3070 1528 7311
5	3435 7871 348 3693 1876 6585 10340 7144 5870 2084 4052 2780
6	3917 3111 3476 1304 10331 5939 5199 1611 1991 699 8316 9960
7	6883 3237 1717 10752 7891 9764 4745 3888 10009 4176 4614 1567
8	10587 2195 1689 2968 5420 2580 2883 6496 111 6023 1024 4449
9	3786 8593 2074 3321 5057 1450 3840 5444 6572 3094 9892 1512
10	8548 1848 10372 4585 7313 6536 6379 1766 9462 2456 5606 9975
11	8204 10593 7935 3636 3882 394 5968 8561 2395 7289 9267 9978
12	7795 74 1633 9542 6867 7352 6417 7568 10623 725 2531 9115
13	7151 2482 4260 5003 10105 7419 9203 6691 8798 2092 8263 3755
14	3600 570 4527 200 9718 6771 1995 8902 5446 768 1103 6520
15	6304 7621
16	6498 9209
17	7293 6786
18	5950 1708
19	8521 1793
20	6174 7854
21	9773 1190
22	9517 10268
23	2181 9349
24	1949 5560
25	1556 555
26	8600 3827
27	5072 1057
28	7928 3542
29	3226 3762
0	7045 2420
1	9645 2641
2	2774 2452
3	5331 2031
4	9400 7503
5	1850 2338
6	10456 9774
7	1692 9276
8	10037 4038
9	3964 338

TABLE 2-continued

Address of Parity Bit Accumulators (Rate 5/6)	
10	2640 5087
11	858 3473
12	5582 5683
13	9523 916
14	4107 1559
15	4506 3491
16	8191 4182
17	10192 6157
18	5668 3305
19	3449 1540
20	4766 2697
21	4069 6675
22	1117 1016
23	5619 3085
24	8483 8400
25	8255 394
26	6338 5042
27	6174 5119
28	7203 1989
29	1781 5174
0	1464 3559
1	3376 4214
2	7238 67
3	10595 8831
4	1221 6513
5	5300 4652
6	1429 9749
7	7878 5131
8	4435 10284
9	6331 5507
10	6662 4941
11	9614 10238
12	8400 8025
13	9156 5630
14	7067 8878
15	9027 3415
16	1690 3866
17	2854 8469
18	6206 630
19	363 5453

TABLE 2-continued

Address of Parity Bit Accumulators (Rate 5/6)
20 4125 7008
21 1612 6702
22 9069 9226
23 5767 4060
24 3743 9237
25 7018 5572
26 8892 4536
27 853 6064
28 8069 5893
29 2051 2885
0 10691 3153
1 3602 4055
2 328 1717
3 2219 9299
4 1939 7898
5 617 206
6 8544 1374
7 10676 3240
8 6672 9489
9 3170 7457
10 7868 5731
11 6121 10732
12 4843 9132
13 580 9591
14 6267 9290
15 3009 2268
16 195 2419
17 8016 1557
18 1516 9195
19 8062 9064
20 2095 8968
21 753 7326
22 6291 3833
23 2614 7844
24 2303 646
25 2075 611
26 4687 362
27 8684 9940
28 4830 2065
29 7038 1363
0 1769 7837
1 3801 1689
2 10070 2359
3 3667 9918
4 1914 6920
5 4244 5669
6 10245 7821
7 7648 3944
8 3310 5488
9 6346 9666
10 7088 6122
11 1291 7827
12 10592 8945
13 3609 7120
14 9168 9112
15 6203 8052
16 3330 2895
17 4264 10563
18 10556 6496
19 8807 7645
20 1999 4530
21 9202 6818
22 3403 1734
23 2106 9023
24 6881 3883
25 3895 2171
26 4062 6424
27 3755 9536
28 4683 2131
29 7347 8027

TABLE 3

Address of Parity Bit Accumulators (Rate 1/2)
54 9318 14392 27561 26909 10219 2534 8597
55 7263 4635 2530 28130 3033 23830 3651
56 24731 23583 26036 17299 5750 792 9169
57 5811 26154 18653 11551 15447 13685 16264
58 12610 11347 28768 2792 3174 29371 12997
59 16789 16018 21449 6165 21202 15850 3186
60 31016 21449 17618 6213 12166 8334 18212
61 22836 14213 11327 5896 718 11727 9308
62 2091 24941 29966 23634 9013 15587 5444
63 22207 3983 16904 28534 21415 27524 25912
64 25687 4501 22193 14665 14798 16158 5491
65 4520 17094 23397 4264 22370 16941 21526
66 10490 6182 32370 9597 30841 25954 2762
67 22120 22865 29870 15147 13668 14955 19235
68 6689 18408 18346 9918 25746 5443 20645
69 29982 12529 13858 4746 30370 10023 24828
70 1262 28032 29888 13063 24033 21951 7863
71 6594 29642 31451 14831 9509 9335 31552
72 1358 6454 16633 20354 24598 624 5265
73 19529 295 18011 3080 13364 8032 15323
74 11981 1510 7960 21462 9129 11370 25741
75 9276 29656 4543 30699 20646 21921 28050
76 15975 25634 5520 31119 13715 21949 19605
77 18688 4608 31755 30165 13103 10706 29224
78 21514 23117 12245 26035 31656 25631 30699
79 9674 24966 31285 29908 17042 24588 31857
80 21856 27777 29919 27000 14897 11409 7122
81 29773 23310 263 4877 28622 20545 22092
82 15605 5651 21864 3967 14419 22757 15896
83 30145 1759 10139 29223 26086 10556 5098
84 18815 16575 2936 24457 26738 6030 505
85 30326 22298 27562 20131 26390 6247 24791
86 928 29246 21246 12400 15311 32309 18608
87 20314 6025 26689 16302 2296 3244 19613
88 6237 11943 22851 15642 23857 15112 20947
89 26403 25168 19038 18384 8882 12719 7093
0 14567 24965
1 3908 100
2 10279 240
3 24102 764
4 12383 4173
5 13861 15918
6 21327 1046
7 5288 14579
8 28158 8069
9 16583 11098
10 16681 28363
11 13980 24725
12 32169 17989
13 10907 2767
14 21557 3818
15 26676 12422
16 7676 8754
17 14905 20232
18 15719 24646
19 31942 8589
20 19978 27197
21 27060 15071
22 6071 26649
23 10393 11176
24 9597 13370
25 7081 17677
26 1433 19513
27 26925 9014
28 19202 8900
29 18152 30647
30 20803 1737
31 11804 25221
32 31683 17783
33 29694 9345
34 12280 26611
35 6526 26122
36 26165 11241
37 7666 26962

TABLE 3-continued

Address of Parity Bit Accumulators (Rate 1/2)
38 16290 8480
39 11774 10120
40 30051 30426
41 1335 15424
42 6865 17742
43 31779 12489
44 32120 21001
45 14508 6996
46 979 25024
47 4554 21896
48 7989 21777
49 4972 20661
50 6612 2730
51 12742 4418
52 29194 595
53 19267 20113

TABLE 4

Address of Parity Bit Accumulators (Rate 3/4)
0 6385 7901 14611 13389 11200 3252 5243 2504 2722 821 7374
1 11359 2698 357 13824 12772 7244 6752 15310 852 2001 11417
2 7862 7977 6321 13612 12197 14449 15137 13860 1708 6399 13444
3 1560 11804 6975 13292 3646 3812 8772 7306 5795 14327 7866
4 7626 11407 14599 9689 1628 2113 10809 9283 1230 15241 4870
5 1610 5699 15876 9446 12515 1400 6303 5411 14181 13925 7358
6 4059 8836 3405 7853 7992 15336 5970 10368 10278 9675 4651
7 4441 3963 9153 2109 12683 7459 12030 12221 629 15212 406
8 6007 8411 5771 3497 543 14202 875 9186 6235 13908 3563
9 3232 6625 4795 546 9781 2071 7312 3399 7250 4932 12652
10 8820 10088 11090 7069 6585 13134 10158 7183 488 7455 9238
11 1903 10818 119 215 7558 11046 10615 11545 14784 7961 15619
12 3655 8736 4917 15874 5129 2134 15944 14768 7150 2692 1469
13 8316 3820 505 8923 6757 806 7957 4216 15589 13244 2622
14 14463 4852 15733 3041 11193 12860 13673 8152 6551 15108 8758
15 3149 11981
16 13416 6906
17 13098 13352
18 2009 14460
19 7207 4314
20 3312 3945
21 4418 6248
22 2669 13975
23 7571 9023
24 14172 2967
25 7271 7138
26 6135 13670
27 7490 14559
28 8657 2466
29 8599 12834
30 3470 3152
31 13917 4365
32 6024 13730
33 10973 14182
34 2464 13167
35 5281 15049
36 1103 1849
37 2058 1069
38 9654 6095
39 14311 7667
40 15617 8146
41 4588 11218
42 13660 6243
43 8578 7874
44 11741 2686
0 1022 1264
1 12604 9965
2 8217 2707
3 3156 11793
4 354 1514

TABLE 4-continued

Address of Parity Bit Accumulators (Rate 3/4)
5 6978 14058
6 7922 16079
7 15087 12138
8 5053 6470
9 12687 14932
10 15458 1763
11 8121 1721
12 12431 549
13 4129 7091
14 1426 8415
15 9783 7604
16 6295 11329
17 1409 12061
18 8065 9087
19 2918 8438
20 1293 14115
21 3922 13851
22 3851 4000
23 5865 1768
24 2655 14957
25 5565 6332
26 4303 12631
27 11653 12236
28 16025 7632
29 4655 14128
30 9584 13123
31 13987 9597
32 15409 12110
33 8754 15490
34 7416 15325
35 2909 15549
36 2995 8257
37 9406 4791
38 11111 4854
39 2812 8521
40 8476 14717
41 7820 15360
42 1179 7939
43 2357 8678
44 7703 6216
0 3477 7067
1 3931 13845
2 7675 12899
3 1754 8187
4 7785 1400
5 9213 5891
6 2494 7703
7 2576 7902
8 4821 15682
9 10426 11935
10 1810 904
11 11332 9264
12 11312 3570
13 14916 2650
14 7679 7842
15 6089 13084
16 3938 2751
17 8509 4648
18 12204 8917
19 5749 12443
20 12613 4431
21 1344 4014
22 8488 13850
23 1730 14896
24 14942 7126
25 14983 8863
26 6578 8564
27 4947 396
28 297 12805
29 13878 6692
30 11857 11186
31 14395 11493
32 16145 12251
33 13462 7428

TABLE 4-continued

Address of Parity Bit Accumulators (Rate 3/4)
34 14526 13119
35 2535 11243
36 6465 12690
37 6872 9334
38 15371 14023
39 8101 10187
40 11963 4848
41 15125 6119
42 8051 14465
43 11139 5167
44 2883 14521

TABLE 5

Address of Parity Bit Accumulators (Rate 4/5)
0 149 11212 5575 6360 12559 8108 8505 408 10026 12828
1 5237 490 10677 4998 3869 3734 3092 3509 7703 10305
2 8742 5553 2820 7085 12116 10485 564 7795 2972 2157
3 2699 4304 8350 712 2841 3250 4731 10105 517 7516
4 12067 1351 11992 12191 11267 5161 537 6166 4246 2363
5 6828 7107 2127 3724 5743 11040 10756 4073 1011 3422
6 11259 1216 9526 1466 10816 940 3744 2815 11506 11573
7 4549 11507 1118 1274 11751 5207 7854 12803 4047 6484
8 8430 4115 9440 413 4455 2262 7915 12402 8579 7052
9 3885 9126 5665 4505 2343 253 4707 3742 4166 1556
10 1704 8936 6775 8639 8179 7954 8234 7850 8883 8713
11 11716 4344 9087 11264 2274 8832 9147 11930 6054 5455
12 7323 3970 10329 2170 8262 3854 2087 12899 9497 11700
13 4418 1467 2490 5841 817 11453 533 11217 11962 5251
14 1541 4525 7976 3457 9536 7725 3788 2982 6307 5997
15 11484 2739 4023 12107 6516 551 2572 6628 8150 9852
16 6070 1761 4627 6534 7913 3730 11866 1813 12306 8249
17 12441 5489 8748 7837 7660 2102 11341 2936 6712 11977
18 10155 4210
19 1010 10483
20 8900 10250
21 10243 12278
22 7070 4397
23 12271 3887
24 11980 6836
25 9514 4356
26 7137 10281
27 11881 2526
28 1969 11477
29 3044 10921
30 2236 8724
31 9104 6340
32 7342 8582
33 11675 10405
34 6467 12775
35 3186 12198
0 9621 11445
1 7486 5611
2 4319 4879
3 2196 344
4 7527 6650
5 10693 2440
6 6755 2706
7 5144 5998
8 11043 8033
9 4846 4435
10 4157 9228
11 12270 6562
12 11954 7592
13 7420 2592
14 8810 9636
15 689 5430
16 920 1304
17 1253 11934
18 9559 6016

TABLE 5-continued

Address of Parity Bit Accumulators (Rate 4/5)
19 312 7589
20 4439 4197
21 4002 9555
22 12232 7779
23 1494 8782
24 10749 3969
25 4368 3479
26 6316 5342
27 2455 3493
28 12157 7405
29 6598 11495
30 11805 4455
31 9625 2090
32 4731 2321
33 3578 2608
34 8504 1849
35 4027 1151
0 5647 4935
1 4219 1870
2 10968 8054
3 6970 5447
4 3217 5638
5 8972 669
6 5618 12472
7 1457 1280
8 8868 3883
9 8866 1224
10 8371 5972
11 266 4405
12 3706 3244
13 6039 5844
14 7200 3283
15 1502 11282
16 12318 2202
17 4523 965
18 9587 7011
19 2552 2051
20 12045 10306
21 11070 5104
22 6627 6906
23 9889 2121
24 829 9701
25 2201 1819
26 6689 12925
27 2139 8757
28 12004 5948
29 8704 3191
30 8171 10933
31 6297 7116
32 616 7146
33 5142 9761
34 10377 8138
35 7616 5811
0 7285 9863
1 7764 10867
2 12343 9019
3 4414 8331
4 3464 642
5 6960 2039
6 786 3021
7 710 2086
8 7423 5601
9 8120 4885
10 12385 11990
11 9739 10034
12 424 10162
13 1347 7597
14 1450 112
15 7965 8478
16 8945 7397
17 6590 8316
18 6838 9011
19 6174 9410
20 255 113

TABLE 5-continued

Address of Parity Bit Accumulators (Rate 4/5)
21 6197 5835
22 12902 3844
23 4377 3505
24 5478 8672
25 4453 2132
26 9724 1380
27 12131 11526
28 12323 9511
29 8231 1752

TABLE 5-continued

Address of Parity Bit Accumulators (Rate 4/5)
30 497 9022
31 9288 3080
32 2481 7515
33 2696 268
34 4023 12341
35 7108 5553

TABLE 6

Address of Parity Bit Accumulators (Rate 3/5)
22422 10282 11626 19997 11161 2922 3122 99 5625 17064 8270 179
25087 16218 17015 828 20041 25656 4186 11629 22599 17305 22515 6463
11049 22853 25706 14388 5500 19245 8732 2177 13555 11346 17265 3069
16581 22225 12563 19717 23577 11555 25496 6853 25403 5218 15925 21766
16529 14487 7643 10715 17442 11119 5679 14155 24213 21000 1116 15620
5340 8636 16693 1434 5635 6516 9482 20189 1066 15013 25361 14243
18506 22236 20912 8952 5421 15691 6126 21595 500 6904 13059 6802
8433 4694 5524 14216 3685 19721 25420 9937 23813 9047 25651 16826
21500 24814 6344 17382 7064 13929 4004 16552 12818 8720 5286 2206
22517 2429 19065 2921 21611 1873 7507 5661 23006 23128 20543 19777
1770 4636 20900 14931 9247 12340 11008 12966 4471 2731 16445 791
6635 14556 18865 22421 22124 12697 9803 25485 7744 18254 11313 9004
19982 23963 18912 7206 12500 4382 20067 6177 21007 1195 23547 24837
756 11158 14646 20534 3647 17728 11676 11843 12937 4402 8261 22944
9306 24009 10012 11081 3746 24325 8060 19826 842 8836 2898 5019
7575 7455 25244 4736 14400 22981 5543 8006 24203 13053 1120 5128
3482 9270 13059 15825 7453 23747 3656 24585 16542 17507 22462 14670
15627 15290 4198 22748 5842 13395 23918 16985 14929 3726 25350 24157
24896 16365 16423 13461 16615 8107 24741 3604 25904 8716 9604 20365
3729 17245 18448 9862 20831 25326 20517 24618 13282 5099 14183 8804
16455 17646 15376 18194 25528 1777 6066 21855 14372 12517 4488 17490
1400 8135 23375 20879 8476 4084 12936 25536 22309 16582 6402 24360
25119 23586 128 4761 10443 22536 8607 9752 25446 15053 1856 4040
377 21160 13474 5451 17170 5938 10256 11972 24210 17833 22047 16108
13075 9648 24546 13150 23867 7309 19798 2988 16858 4825 23950 15125
20526 3553 11525 23366 2452 17626 19265 20172 18060 24593 13255 1552
18839 21132 20119 15214 14705 7096 10174 5663 18651 19700 12524 14033
4127 2971 17499 16287 22368 21463 7943 18880 5567 8047 23363 6797
10651 24471 14325 4081 7258 4949 7044 1078 797 22910 20474 4318
21374 13231 22985 5056 3821 23718 14178 9978 19030 23594 8895 25358
6199 22056 7749 13310 3999 23697 16445 22636 5225 22437 24153 9442
7978 12177 2893 20778 3175 8645 11863 24623 10311 25767 17057 3691
20473 11294 9914 22815 2574 8439 3699 5431 24840 21908 16088 18244
8208 5755 19059 8541 24924 6454 11234 10492 16406 10831 11436 9649
16264 11275 24953 2347 12667 19190 7257 7174 24819 2938 2522 11749
3627 5969 13862 1538 23176 6353 2855 17720 2472 7428 573 15036
0 18539 18661
1 10502 3002
2 9368 10761
3 12299 7828
4 15048 13362
5 18444 24640
6 20775 19175
7 18970 10971
8 5329 19982
9 11296 18655
10 15046 20659
11 7300 22140
12 22029 14477
13 11129 742
14 13254 13813
15 19234 13273

TABLE 6-continued

Address of Parity Bit Accumulators (Rate 3/5)	
16	6079 21122
17	22782 5828
18	19775 4247
19	1660 19413
20	4403 3649
21	13371 25851
22	22770 21784
23	10757 14131
24	16071 21617
25	6393 3725
26	597 19968
27	5743 8084
28	6770 9548
29	4285 17542
30	13568 22599
31	1786 4617
32	23238 11648
33	19627 2030
34	13601 13458
35	13740 17328
36	25012 13944
37	22513 6687
38	4934 12587
39	21197 5133
40	22705 6938
41	7534 24633
42	24400 12797
43	21911 25712
44	12039 1140
45	24306 1021
46	14012 20747
47	11265 15219
48	4670 15531
49	9417 14359
50	2415 6504
51	24964 24690
52	14443 8816
53	6926 1291
54	6209 20806
55	13915 4079
56	24410 13196
57	13505 6117
58	9869 8220
59	1570 6044
60	25780 17387
61	20671 24913
62	24558 20591
63	12402 3702
64	8314 1357
65	20071 14616
66	17014 3688
67	19837 946
68	15195 12136
69	7758 22808
70	3564 2925
71	3434 7769

TABLE 7

Address of Parity Bit Accumulators (Rate 8/9)	
0	6235 2848 3222
1	5800 3492 5348
2	2757 927 90
3	6961 4516 4739
4	1172 3237 6264
5	1927 2425 3683
6	3714 6309 2495
7	3070 6342 7154
8	2428 613 3761

TABLE 7-continued

Address of Parity Bit Accumulators (Rate 8/9)	
9	2906 264 5927
10	1716 1950 4273
11	4613 6179 3491
12	4865 3286 6005
13	1343 5923 3529
14	4589 4035 2132
15	1579 3920 6737
16	1644 1191 5998
17	1482 2381 4620

TABLE 7-continued

Address of Parity Bit Accumulators (Rate 8/9)
18 6791 6014 6596
19 2738 5918 3786
0 5156 6166
1 1504 4356
2 130 1904
3 6027 3187
4 6718 759
5 6240 2870
6 2343 1311
7 1039 5465
8 6617 2513
9 1588 5222
10 6561 535
11 4765 2054
12 5966 6892
13 1969 3869
14 3571 2420
15 4632 981
16 3215 4163
17 973 3117
18 3802 6198
19 3794 3948
0 3196 6126
1 573 1909
2 850 4034
3 5622 1601
4 6005 524
5 5251 5783
6 172 2032
7 1875 2475
8 497 1291
9 2566 3430
10 1249 740
11 2944 1948
12 6528 2899
13 2243 3616
14 867 3733
15 1374 4702
16 4698 2285
17 4760 3917
18 1859 4058
19 6141 3527
0 2148 5066
1 1306 145
2 2319 871
3 3463 1061
4 5554 6647
5 5837 339
6 5821 4932
7 6356 4756
8 3930 418
9 211 3094
10 1007 4928
11 3584 1235
12 6982 2869
13 1612 1013
14 953 4964
15 4555 4410
16 4925 4842
17 5778 600
18 6509 2417
19 1260 4903
0 3369 3031
1 3557 3224
2 3028 583
3 3258 440
4 6226 6655
5 4895 1094
6 1481 6847
7 4433 1932
8 2107 1649
9 2119 2065
10 4003 6388
11 6720 3622

TABLE 7-continued

Address of Parity Bit Accumulators (Rate 8/9)
12 3694 4521
13 1164 7050
14 1965 3613
15 4331 66
16 2970 1796
17 4652 3218
18 1762 4777
19 5736 1399
0 970 2572
1 2062 6599
2 4597 4870
3 1228 6913
4 4159 1037
5 2916 2362
6 395 1226
7 6911 4548
8 4618 2241
9 4120 4280
10 5825 474
11 2154 5558
12 3793 5471
13 5707 1595
14 1403 325
15 6601 5183
16 6369 4569
17 4846 896
18 7092 6184
19 6764 7127
0 6358 1951
1 3117 6960
2 2710 7062
3 1133 3604
4 3694 657
5 1355 110
6 3329 6736
7 2505 3407
8 2462 4806
9 4216 214
10 5348 5619
11 6627 6243
12 2644 5073
13 4212 5088
14 3463 3889
15 5306 478
16 4320 6121
17 3961 1125
18 5699 1195
19 6511 792
0 3934 2778
1 3238 6587
2 1111 6596
3 1457 6226
4 1446 3885
5 3907 4043
6 6839 2873
7 1733 5615
8 5202 4269
9 3024 4722
10 5445 6372
11 370 1828
12 4695 1600
13 680 2074
14 1801 6690
15 2669 1377
16 2463 1681
17 5972 5171
18 5728 4284
19 1696 1459

TABLE 8

Address of Parity Bit Accumulators (Rate 9/10)
0 5611 2563 2900
1 5220 3143 4813
2 2481 834 81
3 6265 4064 4265
4 1055 2914 5638
5 1734 2182 3315
6 3342 5678 2246
7 2185 552 3385
8 2615 236 5334
9 1546 1755 3846
10 4154 5561 3142
11 4382 2957 5400
12 1209 5329 3179
13 1421 3528 6063
14 1480 1072 5398
15 3843 1777 4369
16 1334 2145 4163
17 2368 5055 260
0 6118 5405
1 2994 4370
2 3405 1669
3 4640 5550
4 1354 3921
5 117 1713
6 5425 2866
7 6047 683
8 5616 2582
9 2108 1179
10 933 4921
11 5953 2261
12 1430 4699
13 5905 480
14 4289 1846
15 5374 6208
16 1775 3476
17 3216 2178
0 4165 884
1 2896 3744
2 874 2801
3 3423 5579
4 3404 3552
5 2876 5515
6 516 1719
7 765 3631
8 5059 1441
9 5629 598
10 5405 473
11 4724 5210
12 155 1832
13 1689 2229
14 449 1164
15 2308 3088
16 1122 669
17 2268 5758
0 5878 2609
1 782 3359
2 1231 4231
3 4225 2052
4 4286 3517
5 5531 3184
6 1935 4560
7 1174 131
8 3115 956
9 3129 1088
10 5238 4440
11 5722 4280
12 3540 375
13 191 2782
14 906 4432
15 3225 1111
16 6296 2583
17 1457 903
0 855 4475
1 4097 3970

TABLE 8-continued

Address of Parity Bit Accumulators (Rate 9/10)
2 4433 4361
3 5198 541
4 1146 4426
5 3202 2902
6 2724 525
7 1083 4124
8 2326 6003
9 5605 5990
10 4376 1579
11 4407 984
12 1332 6163
13 5359 3975
14 1907 1854
15 3601 5748
16 6056 3266
17 3322 4085
0 1768 3244
1 2149 144
2 1589 4291
3 5154 1252
4 1855 5939
5 4820 2706
6 1475 3360
7 4266 693
8 4156 2018
9 2103 752
10 3710 3853
11 5123 931
12 6146 3323
13 1939 5002
14 5140 1437
15 1263 293
16 5949 4665
17 4548 6380
0 3171 4690
1 5204 2114
2 6384 5565
3 5722 1757
4 2805 6264
5 1202 2616
6 1018 3244
7 4018 5289
8 2257 3067
9 2483 3073
10 1196 5329
11 649 3918
12 3791 4581
13 5028 3803
14 3119 3506
15 4779 431
16 3888 5510
17 4387 4084
0 5836 1692
1 5126 1078
2 5721 6165
3 3540 2499
4 2225 6348
5 1044 1484
6 6323 4042
7 1313 5603
8 1303 3496
9 3516 3639
10 5161 2293
11 4682 3845
12 3045 643
13 2818 2616
14 3267 649
15 6236 593
16 646 2948
17 4213 1442
0 5779 1596
1 2403 1237
2 2217 1514
3 5609 716

TABLE 8-continued

Address of Parity Bit Accumulators (Rate 9/10)	
4 5155 3858	
5 1517 1312	
6 2554 3158	
7 5280 2643	
8 4990 1353	
9 5648 1170	
10 1152 4366	
11 3561 5368	
12 3581 1411	
13 5647 4661	
14 1542 5401	
15 5078 2687	
16 316 1755	
17 3392 1991	

12. A method according to claim 11, wherein the row indices of 1's in other column indices m (m modulo $360 \neq 0$ and $m < k_{ldpc}$) of the parity check matrix are given by $\{x+m \bmod 360 \times q\} \bmod (n_{ldpc} - k_{ldpc})$, where $q=60$ for rate $\frac{2}{3}$ LDPC code, $q=30$ for rate $\frac{5}{6}$ LDPC code, $q=90$ for rate $\frac{1}{2}$ LDPC code, $q=45$ for rate $\frac{3}{4}$ LDPC code, $q=36$ for rate $\frac{5}{6}$ LDPC code, $q=72$ for rate $\frac{3}{5}$ LDPC code, $q=20$ for rate $\frac{8}{9}$ LDPC code, $q=18$ for rate $\frac{9}{10}$ LDPC code, wherein x denotes an entry at the j^{th} row of Tables 1-7, where $j = \text{int}\{m/360\}$, and $\text{int}\{\cdot\}$ denotes the integer function, the row indices of 1's in the column index $m = k_{ldpc} + j$ ($j=0, 1, 2, \dots, n_{ldpc} - k_{ldpc} - 2$) of the parity check matrix being given by j and $j+1$, the row index of 1 in the column index $n_{ldpc} - 1$ of the parity check matrix being given by $n_{ldpc} - k_{ldpc} - 1$.

13. A computer-readable medium bearing instructions for encoding, said instruction, being arranged, upon execution, to cause one or more processors to perform the method of claim 1.

14. An encoder for generating Low Density Parity Check (LDPC) codes, comprising:

memory storing information representing a structured parity check matrix of the LDPC codes, the information being organized in tabular form, wherein each row represents occurrences of one values within a first column of a group of columns of the parity check matrix, the rows correspond to groups of columns of the parity check matrix, wherein subsequent columns within each of the groups are derived according to a predetermined operation; and

means for retrieving the stored information representing the parity check matrix to output an LDPC coded signal.

15. An encoder according to claim 14, wherein the predetermined operation specifies one of a cyclic shift on the first column of each of the group, and addition of a constant to the first column of each of the group, the constant being dependent on code rate of the LDPC code.

16. An encoder according to claim 14, wherein an i^{th} parity bit is determined by adding the $(i-1)^{th}$ parity bit and the j^{th} information bit if the j^{th} entry in the i^{th} row of the parity check matrix is 1.

17. An encoder according to claim 14, wherein parity bit accumulators are initialized to zero, the first information bit in the j^{th} group of M information bits is accumulated in the i^{th} parity bit accumulator if the i^{th} entry in $(jM)^{th}$ column of the parity check matrix is 1, where $j=0, 1, 2, 3, \dots, k_{ldpc}/M - 1$,

the remaining $(M-1)$ information bits $m=jM+1, jM+2, jM+3, \dots, (j+1)M-1$ of the j^{th} group being accumulated in the parity bit accumulators according to $\{x+m \bmod M \times q\} \bmod (n_{ldpc} - k_{ldpc})$, wherein x denotes the address of the parity bit accumulator corresponding to the first bit, jM , in the group, and q is a code rate dependent constant, after all of the information bits are exhausted, operations, starting with $i=1$ are performed according to $p_i = p_i \oplus p_{i-1}$, $i=1, 2, \dots, n_{ldpc} - k_{ldpc} - 1$, wherein final content of p_i , $i=0, 1, \dots, n_{ldpc} - k_{ldpc} - 1$ is equal to the parity bit p_i .

18. An encoder according to claim 17, wherein $M=360$.

19. An encoder according to claim 14, wherein the code dependent constant q is 60, 30, 90, 45, 36, 72, 20, and 18 for code rates $\frac{2}{3}$, $\frac{5}{6}$, $\frac{1}{2}$, $\frac{3}{4}$, $\frac{5}{6}$, $\frac{3}{5}$, $\frac{8}{9}$, and $\frac{9}{10}$, respectively.

20. An encoder according to claim 11, wherein the LDPC coded signal is modulated according to a signal constellation that includes one of 8-PSK (Phase Shift Keying), 16-QAM (Quadrature Amplitude Modulation), QPSK (Quadrature Phase Shift Keying), 16-APSK (Amplitude Phase Shift Keying) and 32-APSK.

21. An encoder according to claim 14, further comprising:

a Bose Chaudhuri Hocquenghem (BCH) encoder configured to encode an input signal using BCH codes, wherein the output LDPC coded signal corresponding to the input signal represents a code having an outer BCH code and an inner LDPC code.

22. An encoder according to claim 21, wherein the number of redundant BCH bits is $n_{BCH} - k_{BCH} = 16 \times t$, wherein t represents error correcting capability of the BCH code.

23. An encoder according to claim 21, wherein the error correction capability of the BCH code is 12 bits when used in concatenation with rate $\frac{1}{2}$, $\frac{3}{4}$, $\frac{5}{6}$ and $\frac{3}{5}$ LDPC codes, is 10 bits when used in concatenation with rate $\frac{2}{3}$ and $\frac{5}{6}$ LDPC codes, and is 8 bits when used in concatenation with rate $\frac{8}{9}$ and $\frac{9}{10}$ LDPC codes.

24. A transmitter utilizing Low Density Parity Check (LDPC) coding, comprising:

memory storing information representing a structured parity check matrix of the LDPC codes, the information being organized in tabular form, wherein each row represents occurrences of one values within a first column of a group of columns of the parity check matrix, the rows correspond to groups of columns of the parity check matrix, wherein subsequent columns within each of the groups are derived according to a predetermined operation; and

an LDPC encoder configured to access the stored information from the memory to output an LDPC coded signal.

25. A transmitter according to claim 24, wherein the predetermined operation specifies one of a cyclic shift on the first column of each of the group, and addition of a constant to the first column of each of the group, the constant being dependent on code rate of the LDPC code.

26. A transmitter according to claim 24, wherein an i^{th} parity bit is determined by adding the $(i-1)^{th}$ parity bit and the j^{th} information bit if the j^{th} entry in the i^{th} row of the parity check matrix is 1.

27. A transmitter according to claim 24, wherein parity bit accumulators are initialized to zero, the first information bit in the j^{th} group of M information bits is accumulated in the i^{th} parity bit accumulator if the i^{th} entry in $(jM)^{th}$ column of the parity check matrix is 1, where $j=0, 1, 2, 3, \dots, k_{ldpc}/M - 1$,

the remaining $(M-1)$ information bits $m=jM+1, jM+2, jM+3, \dots, (j+1)M-1$ of the j^{th} group being accumulated in the parity bit accumulators according to $\{x+m \bmod M \times q\} \bmod (n_{\text{ldpc}} - k_{\text{ldpc}})$, wherein x denotes the address of the parity bit accumulator corresponding to the first bit, jM , in the group, and q is a code rate dependent constant, after all of the information bits are exhausted, operations, starting with $i=1$ are performed according to $p_i = p_i \oplus p_{i-1}$, $i=1, 2, \dots, n_{\text{ldpc}} - k_{\text{ldpc}} - 1$, wherein final content of p_i , $i=0, 1, \dots, n_{\text{ldpc}} - k_{\text{ldpc}} - 1$ is equal to the parity bit p_i .

28. A transmitter according to claim 27, wherein $M=360$.

29. A transmitter according to claim 24, wherein the code dependent constant q is 60, 30, 90, 45, 36, 72, 20, and 18 for code rates $\frac{2}{3}$, $\frac{5}{6}$, $\frac{1}{2}$, $\frac{3}{4}$, $\frac{4}{5}$, $\frac{3}{5}$, $\frac{8}{9}$, and $\frac{9}{10}$, respectively.

30. A transmitter according to claim 24, wherein the LDPC coded signal is modulated according to a signal constellation that includes one of 8-PSK (Phase Shift Keying), 16-QAM (Quadrature Amplitude Modulation), QPSK (Quadrature Phase Shift Keying), 16-APSK (Amplitude Phase Shift Keying) and 32-APSK.

31. A transmitter according to claim 24, further comprising:

a Bose Chaudhuri Hocquenghem (BCH) transmitter configured to encode an input signal using BCH codes, wherein the output LDPC coded signal corresponding to the input signal represents a code having an outer BCH code and an inner LDPC code.

32. A transmitter according to claim 31, wherein the number of redundant BCH bits is $n_{\text{BCH}} - k_{\text{BCH}} = 16 \times t$, wherein t represents error correcting capability of the BCH code.

33. A transmitter according to claim 31, wherein the error correction capability of the BCH code is 12 bits when used in concatenation with rate $\frac{1}{2}$, $\frac{3}{4}$, $\frac{4}{5}$ and $\frac{3}{5}$ LDPC codes, is 10 bits when used in concatenation with rate $\frac{2}{3}$ and $\frac{5}{6}$ LDPC codes, and is 8 bits when used in concatenation with rate $\frac{8}{9}$ and $\frac{9}{10}$ LDPC codes.

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