US008407722B2

(12) **United States Patent**
Tuttle et al.

(10) **Patent No.: US 8,407,722 B2**
(45) **Date of Patent: Mar. 26, 2013**

(54) **ASYNCHRONOUS MESSAGING USING A NODE SPECIALIZATION ARCHITECTURE IN THE DYNAMIC ROUTING NETWORK**

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(*) Notice: Subject to any disclaimer, the term of this patent is extended or adjusted under 35 U.S.C. 154(b) by 1127 days.

(21) Appl. No.: **11/396,251**

(22) Filed: **Mar. 30, 2006**

(65) **Prior Publication Data**

US 2007/0239822 A1 Oct. 11, 2007

Related U.S. Application Data

(63) Continuation of application No. 10/105,018, filed on Mar. 21, 2002, now Pat. No. 7,051,070, which is a continuation-in-part of application No. 10/017,182, filed on Dec. 14, 2001, now Pat. No. 7,043,525.

(Continued)

(51) **Int. Cl.**
G06F 3/00 (2006.01)
G06F 15/16 (2006.01)

(52) **U.S. Cl.** **719/316**; 709/206

(58) **Field of Classification Search** 709/203, 709/219; 719/309

See application file for complete search history.

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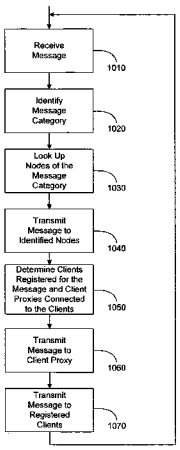
(57) **ABSTRACT**

A network routes update messages containing updates to properties of live objects from input sources to clients having the objects. When the clients receive live objects, the clients identify the object IDs associated with the objects and register the object IDs with the routing network. The routing network is adapted to selectively send update messages to nodes in the network and the nodes forward the messages to the clients. One implementation uses a hierarchy of registries to indicate which nodes and clients receive which update messages. Another implementation assigns update messages to one or more of N categories and nodes to one or more of M types, and the gateways maintain mapping between categories and types. To ensure that clients receive all of the update messages for which they register, the clients connect to client proxies that in turn connect to at least one node of each type.

37 Claims, 12 Drawing Sheets



Related U.S. Application Data

(60) Provisional application No. 60/256,613, filed on Dec. 18, 2000, provisional application No. 60/276,847, filed on Mar. 16, 2001, provisional application No. 60/278,303, filed on Mar. 21, 2001, provisional application No. 60/279,608, filed on Mar. 28, 2001, provisional application No. 60/280,627, filed on Mar. 29, 2001.

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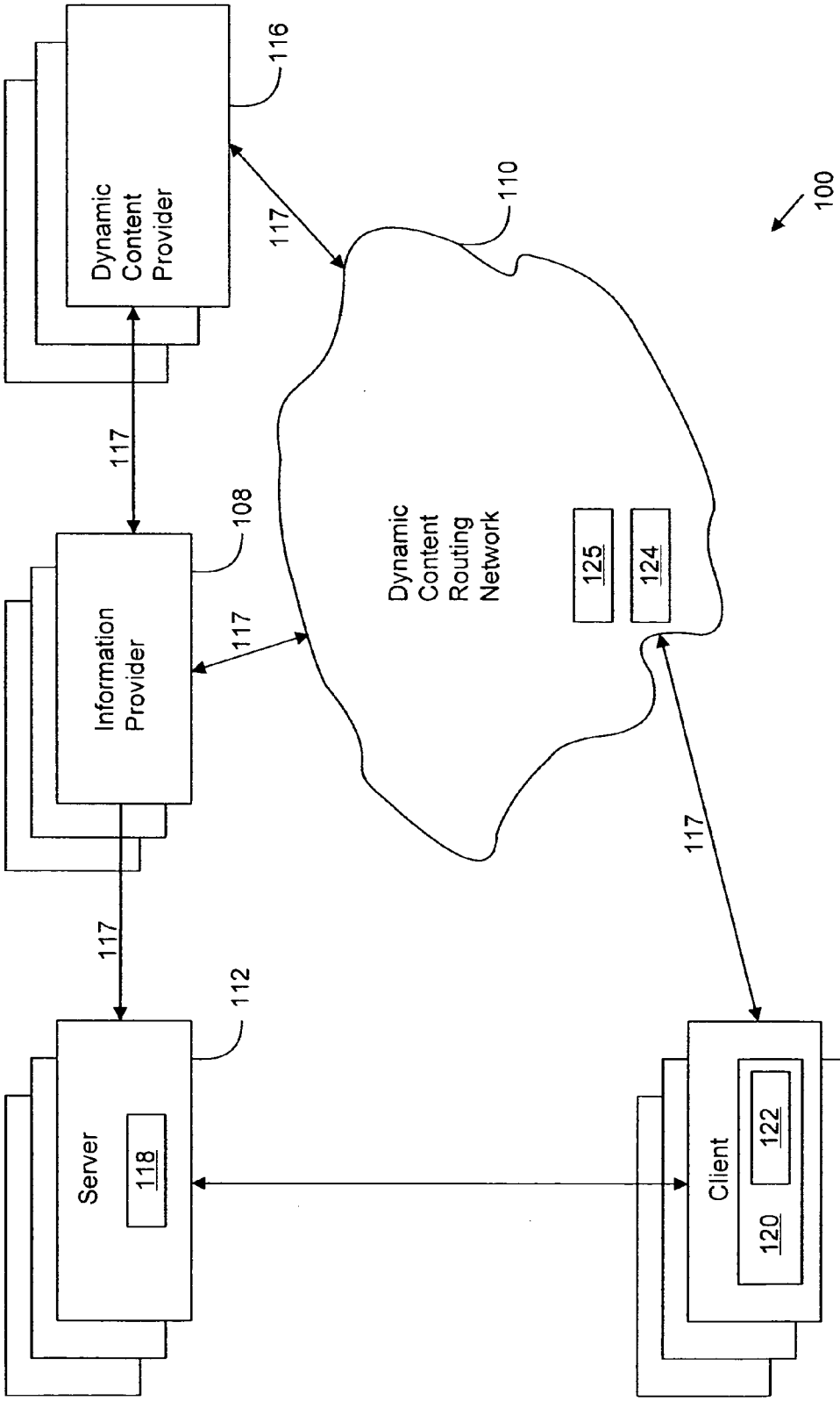
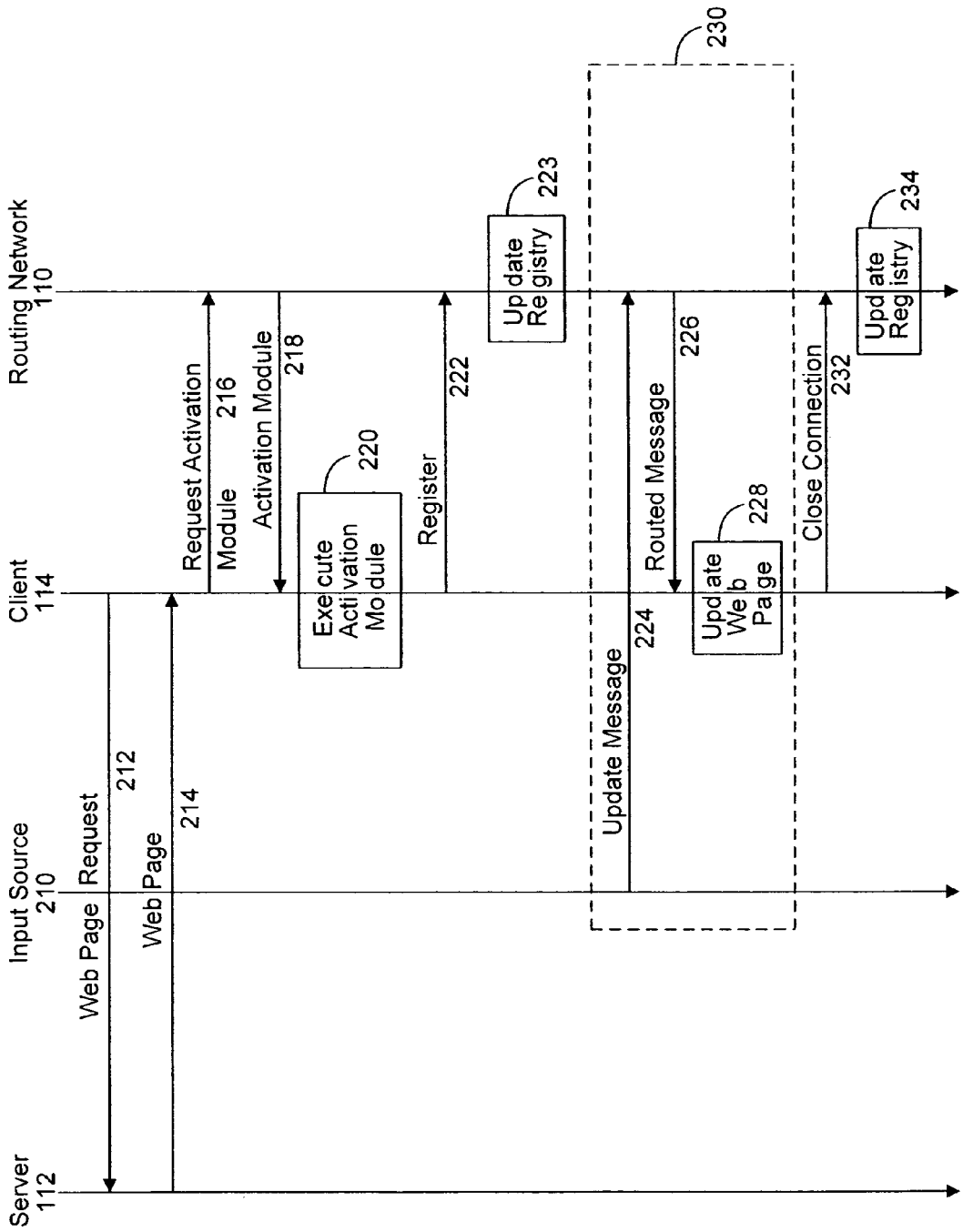

Dynamic Content Provider
116
117
117
Information Provider
108
117
117
Server
118
112
110
Dynamic Content Routing Network
125
124
100
117
Client
120
122
114

Fig. 1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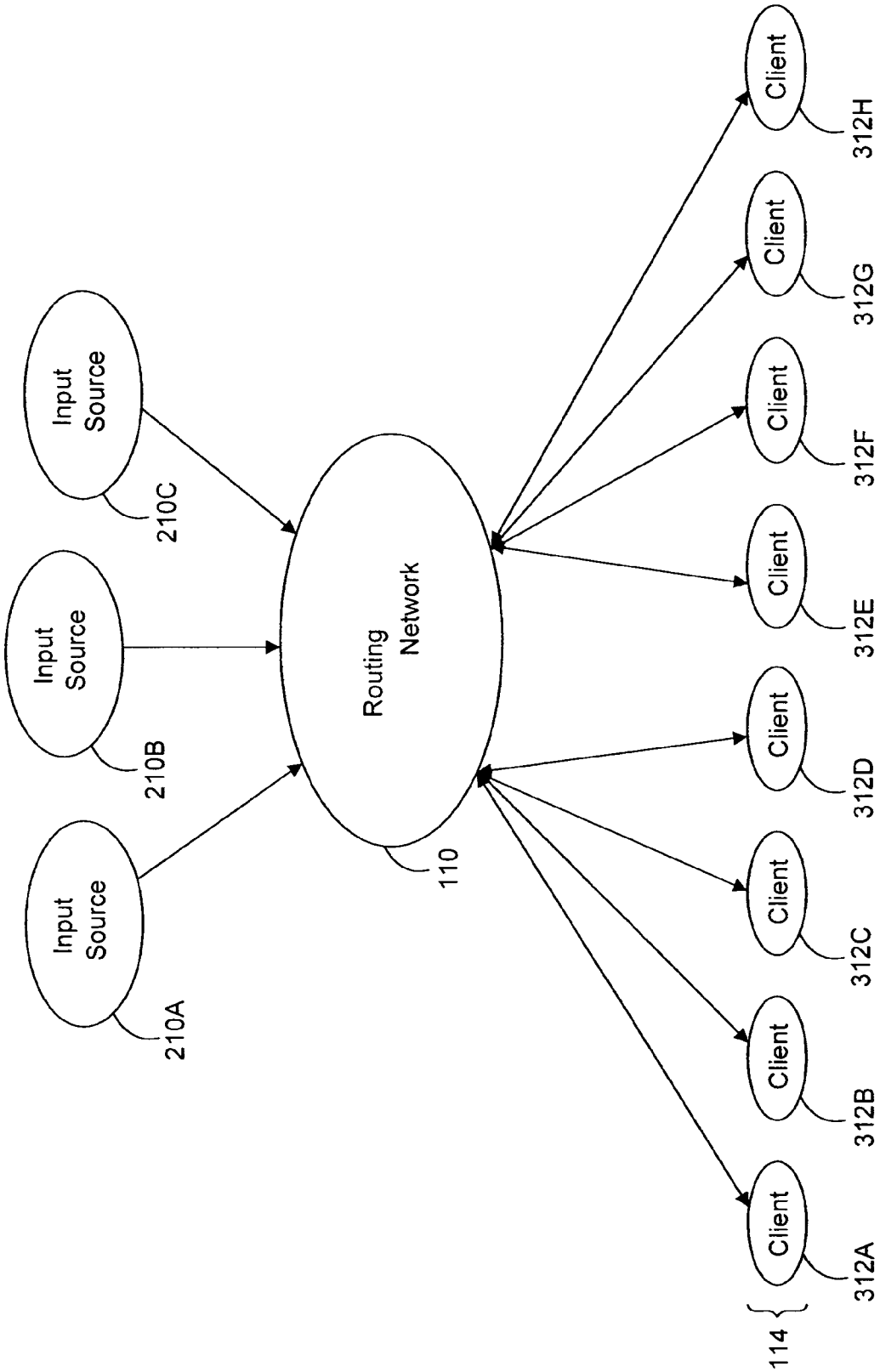
Input Source 210
Client 114
Routing Network 110
Web Page
Request
212
Web Page
214
Request Activation Module
216
Activation Module
218
Execute Activation Module
220
Register
222
Update Registry
223
Update Message
224
Routed Message
226
Update Web Page
228
230
Close Connection
232
Update Registry
234

Fig. 2


Input Source
210A
Input Source
210B
Input Source
210C
Routing Network
110
Client
312A
Client
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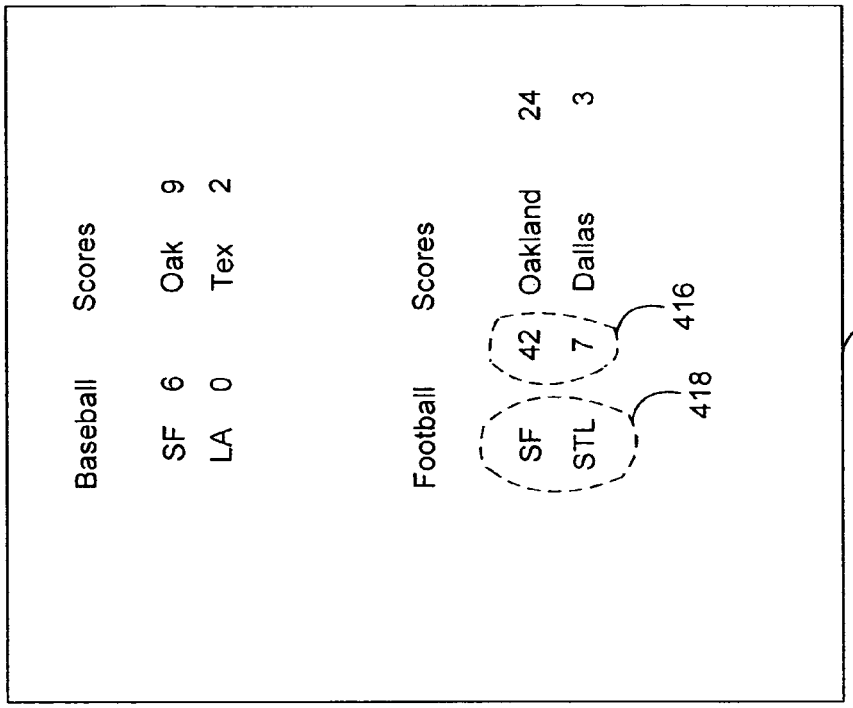
Client
312C
Client
312D
Client
312E
Client
312F
Client
312G
Client
312H
114

Fig. 3


Baseball Scores
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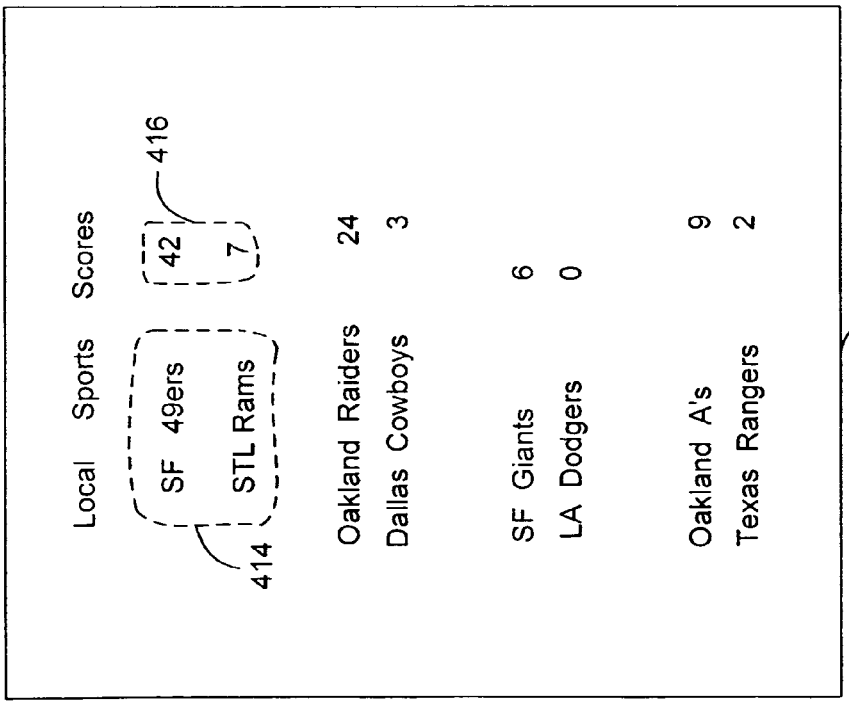
LA 0 Tex 2
Football Scores
SF 42 Oakland 24
STL 7 Dallas 3
416
418
412


Local Sports Scores
SF 49ers 42
STL Rams 7
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414
Oakland Raiders 24
Dallas Cowboys 3
SF Giants 6
LA Dodgers 0
Oakland A's 9
Texas Rangers 2
410

Fig. 4

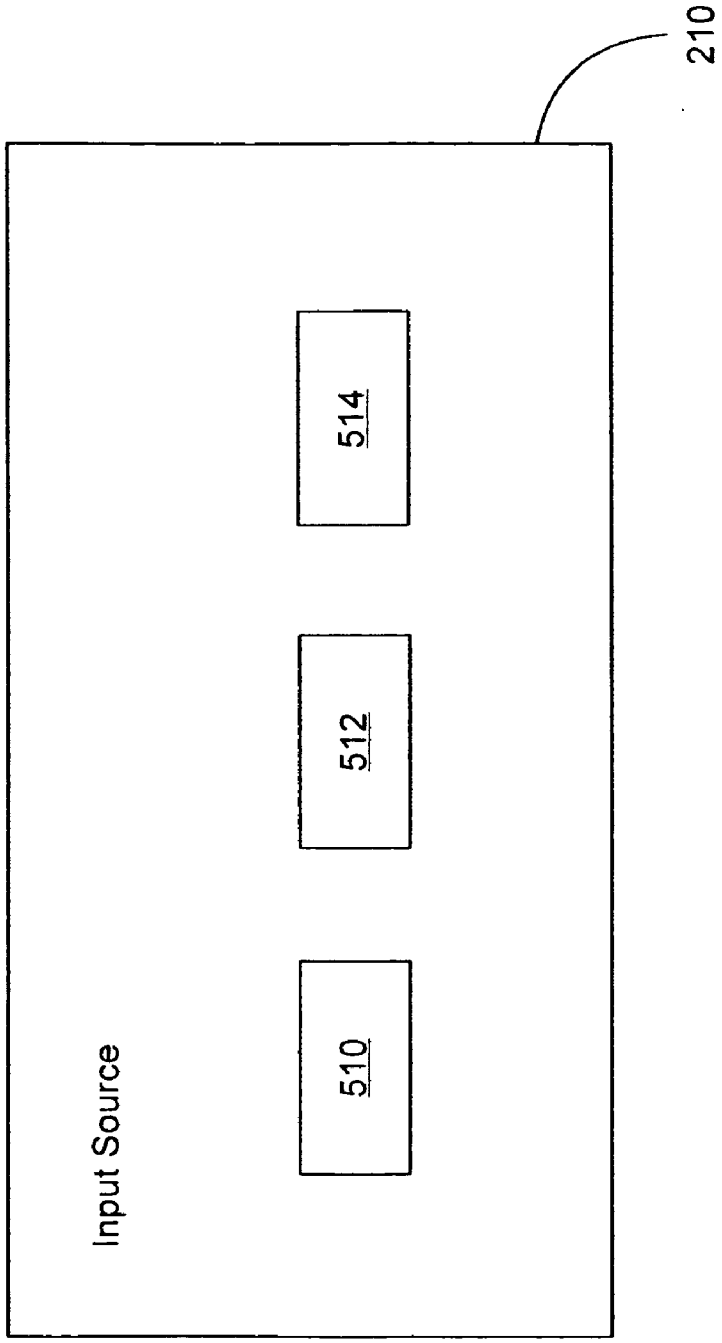
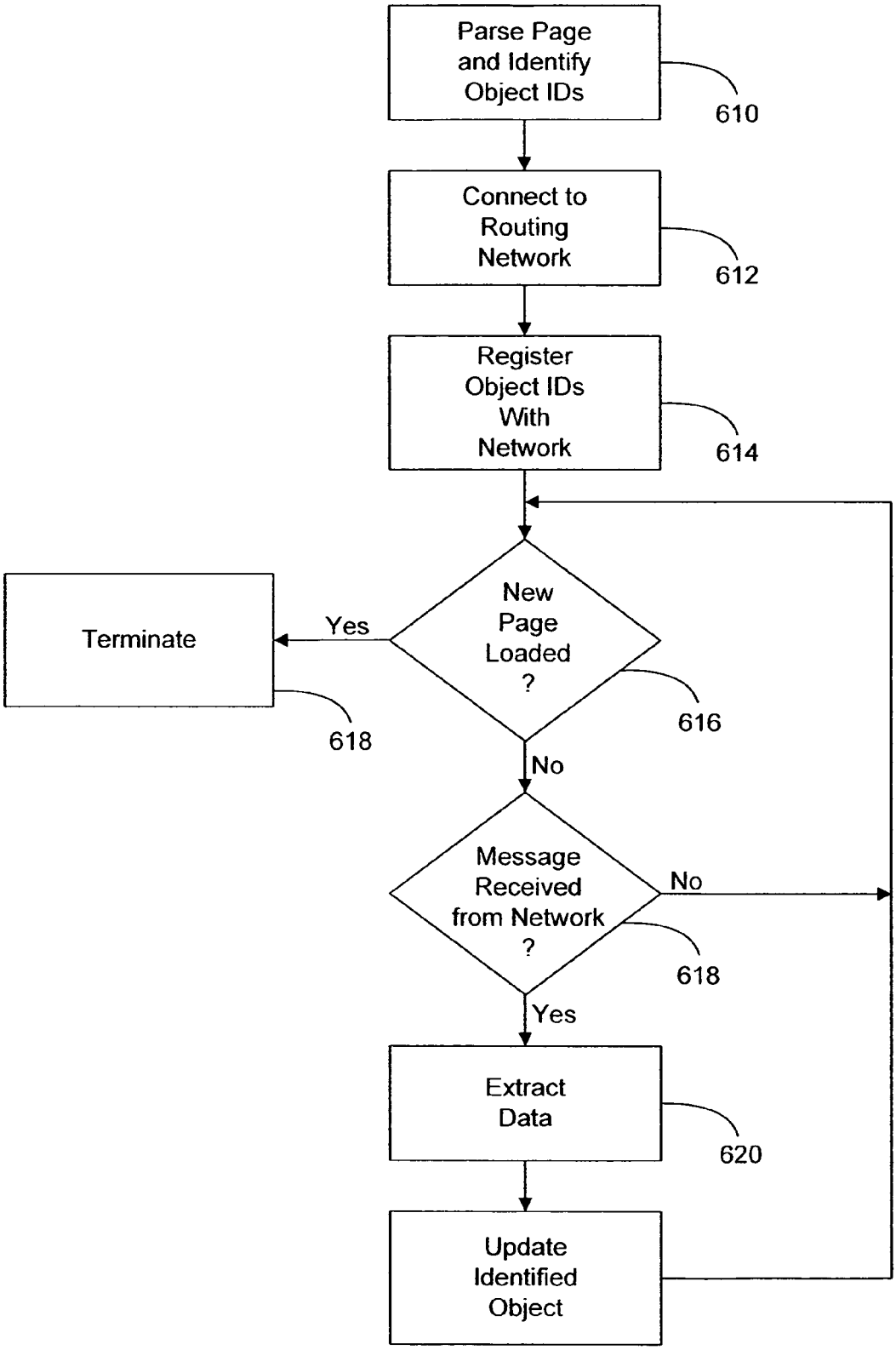


Fig. 5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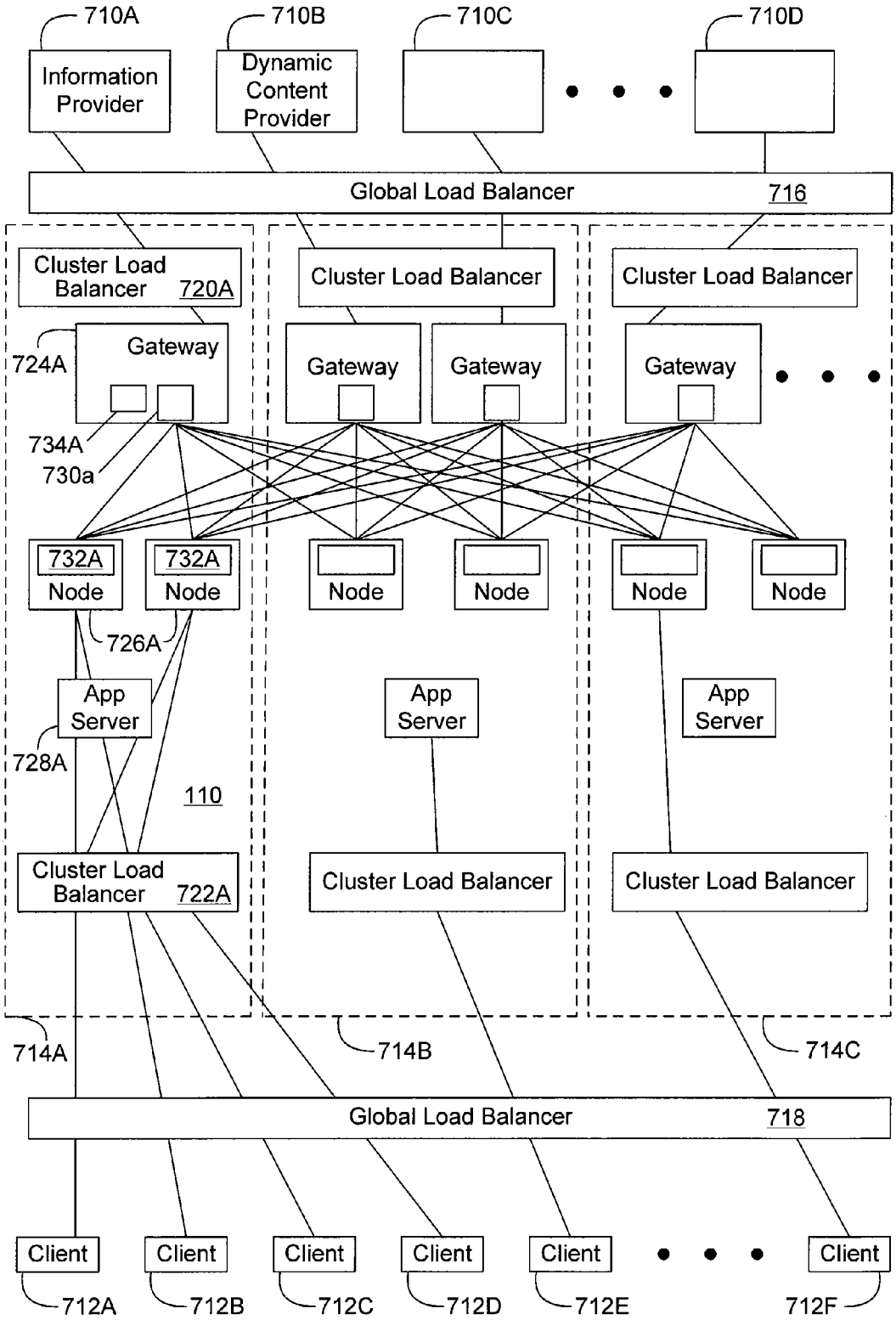
610
Connect to Routing Network
612
Register Object IDs With Network
614
New Page Loaded ?
616
Yes
Terminate
618
No
Message Received from Network ?
618
No
Yes
Extract Data
620
Update Identified Object

Fig. 6


710A
710B
710C
710D
Information Provider
Dynamic Content Provider
Global Load Balancer
716
Cluster Load Balancer
720A
Cluster Load Balancer
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Gateway
Gateway
Gateway
734A
730a
732A
Node
732A
Node
Node
Node
Node
Node
726A
App Server
App Server
App Server
728A
110
Cluster Load Balancer
722A
Cluster Load Balancer
Cluster Load Balancer
714A
714B
714C
Global Load Balancer
718
Client
Client
Client
Client
Client
Client
712A
712B
712C
712D
712E
712F

Fig. 7

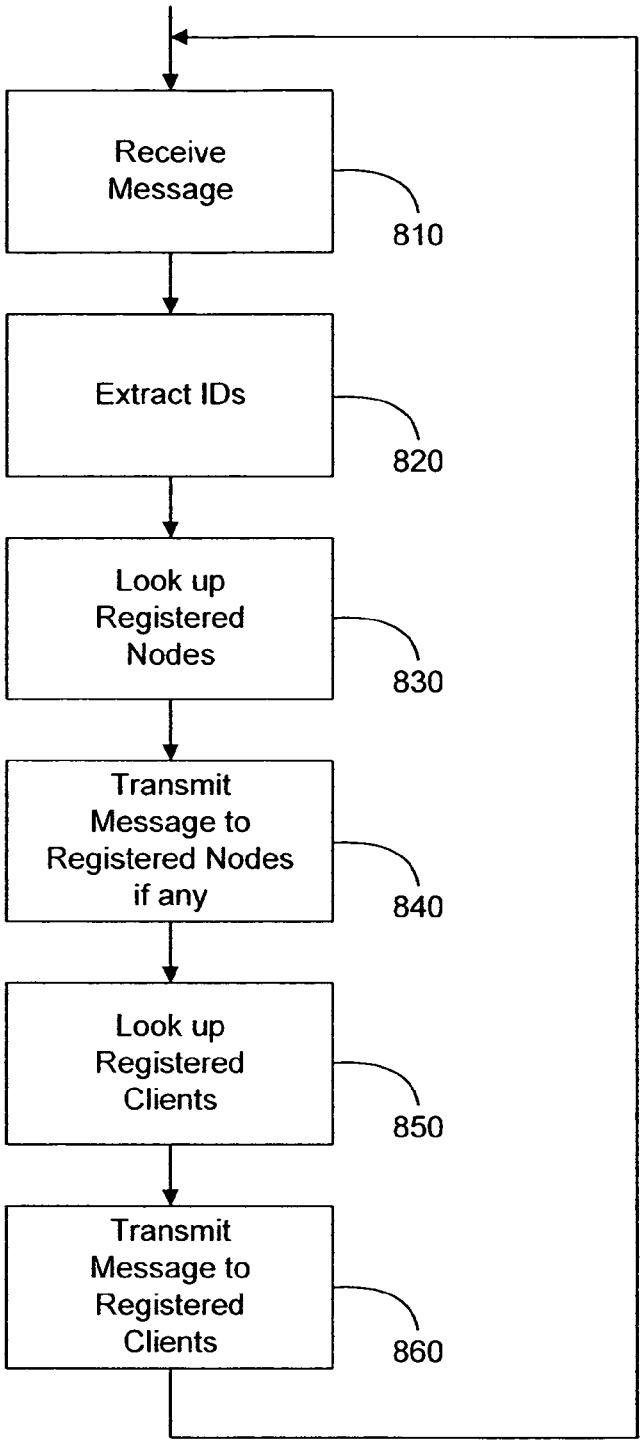
Receive Message
810
Extract IDs
820
Look up Registered Nodes
830
Transmit Message to Registered Nodes if any
840
Look up Registered Clients
850
Transmit Message to Registered Clients
860

Fig. 8

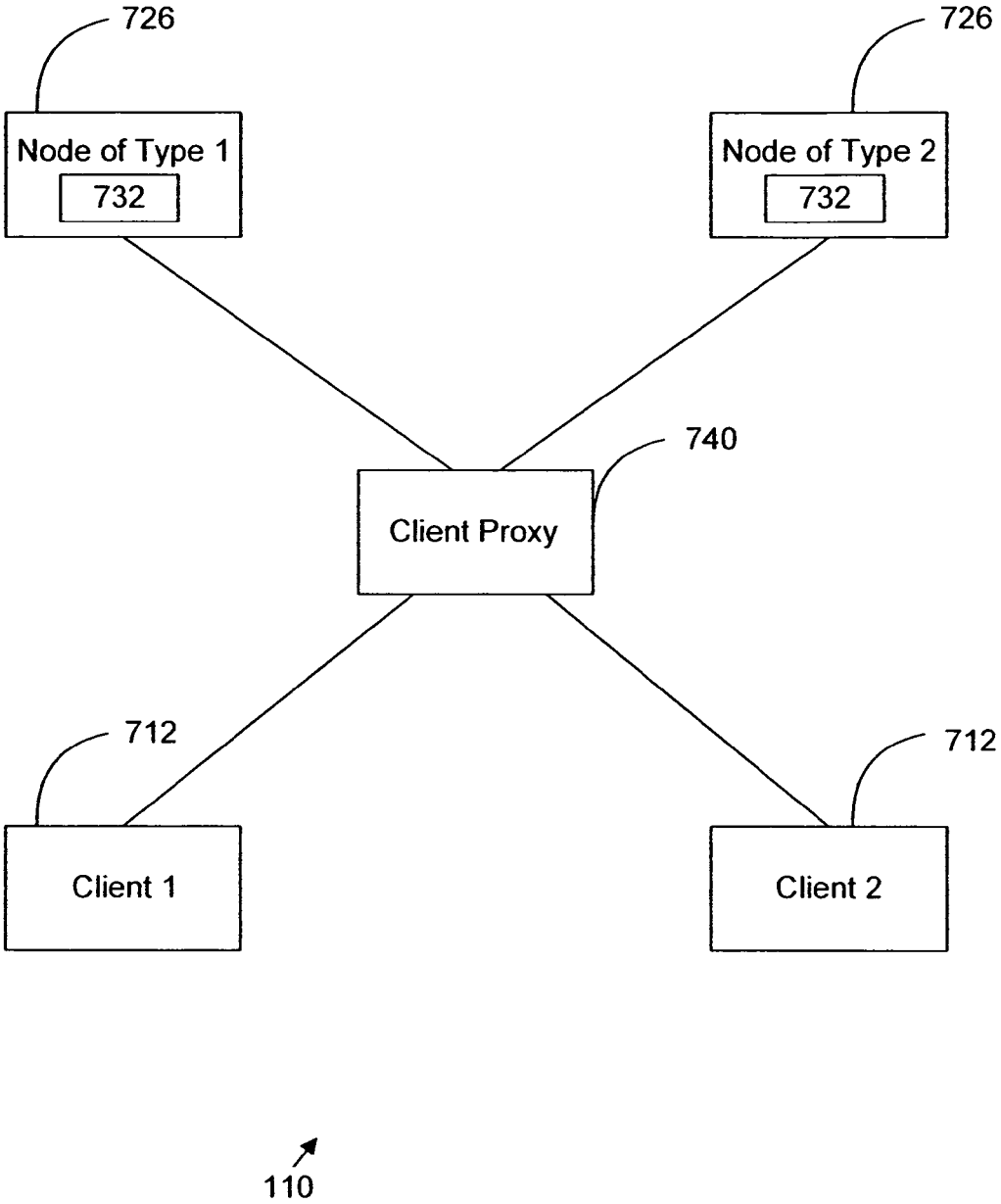
726
Node of Type 1
732
726
Node of Type 2
732
740
Client Proxy
712
Client 1
712
Client 2
110

Fig. 9

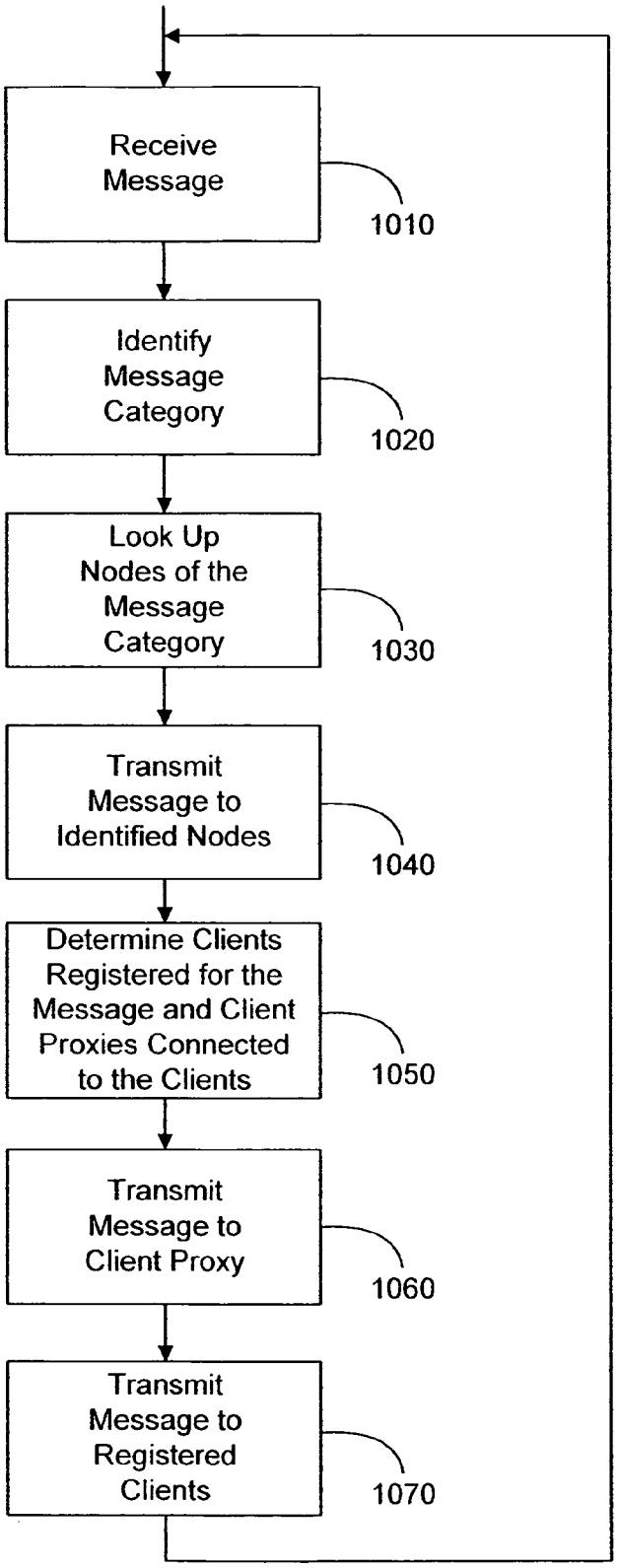
Receive Message
1010
Identify Message Category
1020
Look Up Nodes of the Message Category
1030
Transmit Message to Identified Nodes
1040
Determine Clients Registered for the Message and Client Proxies Connected to the Clients
1050
Transmit Message to Client Proxy
1060
Transmit Message to Registered Clients
1070

Fig. 10

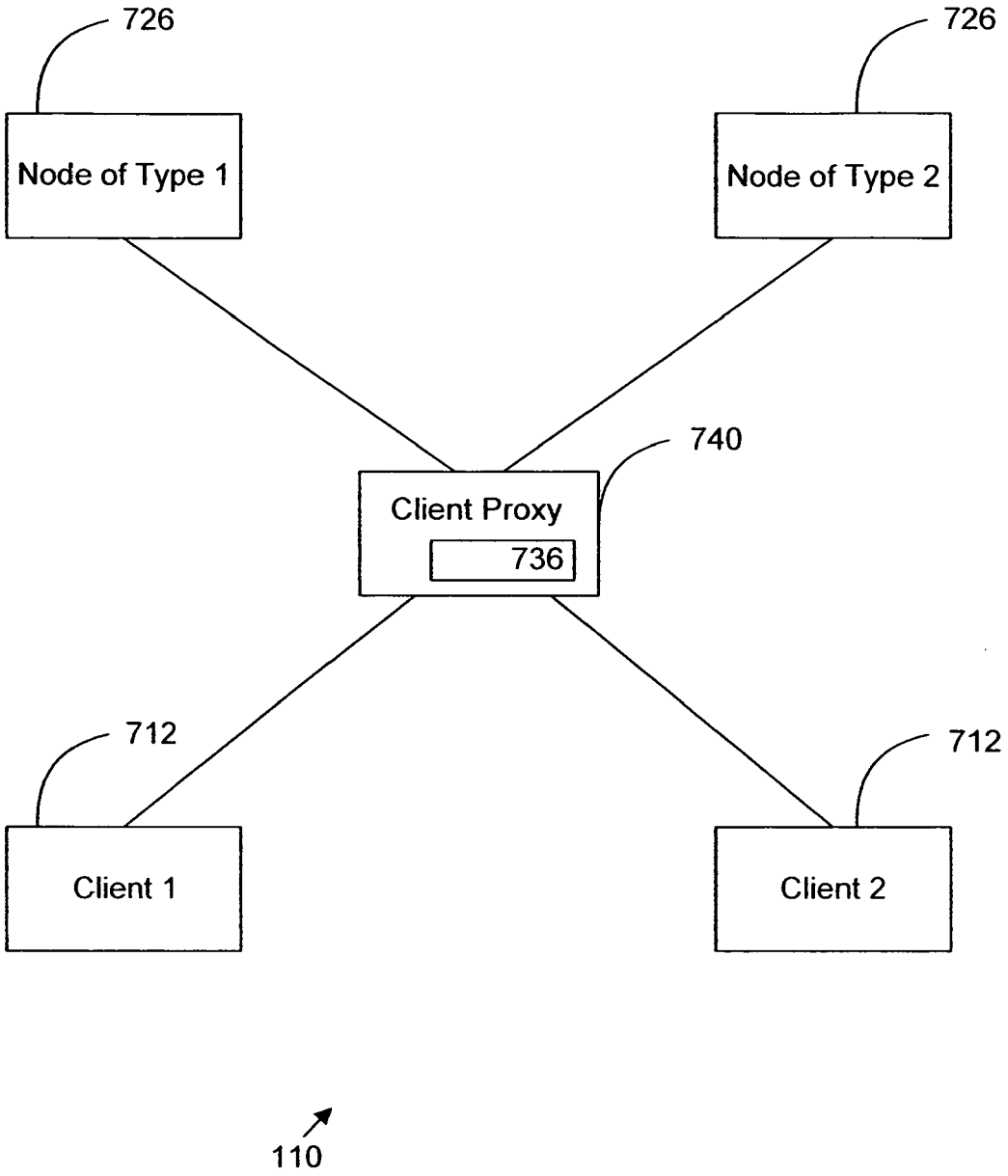
726
Node of Type 1
726
Node of Type 2
740
Client Proxy
736
712
Client 1
712
Client 2
110

Fig. 11

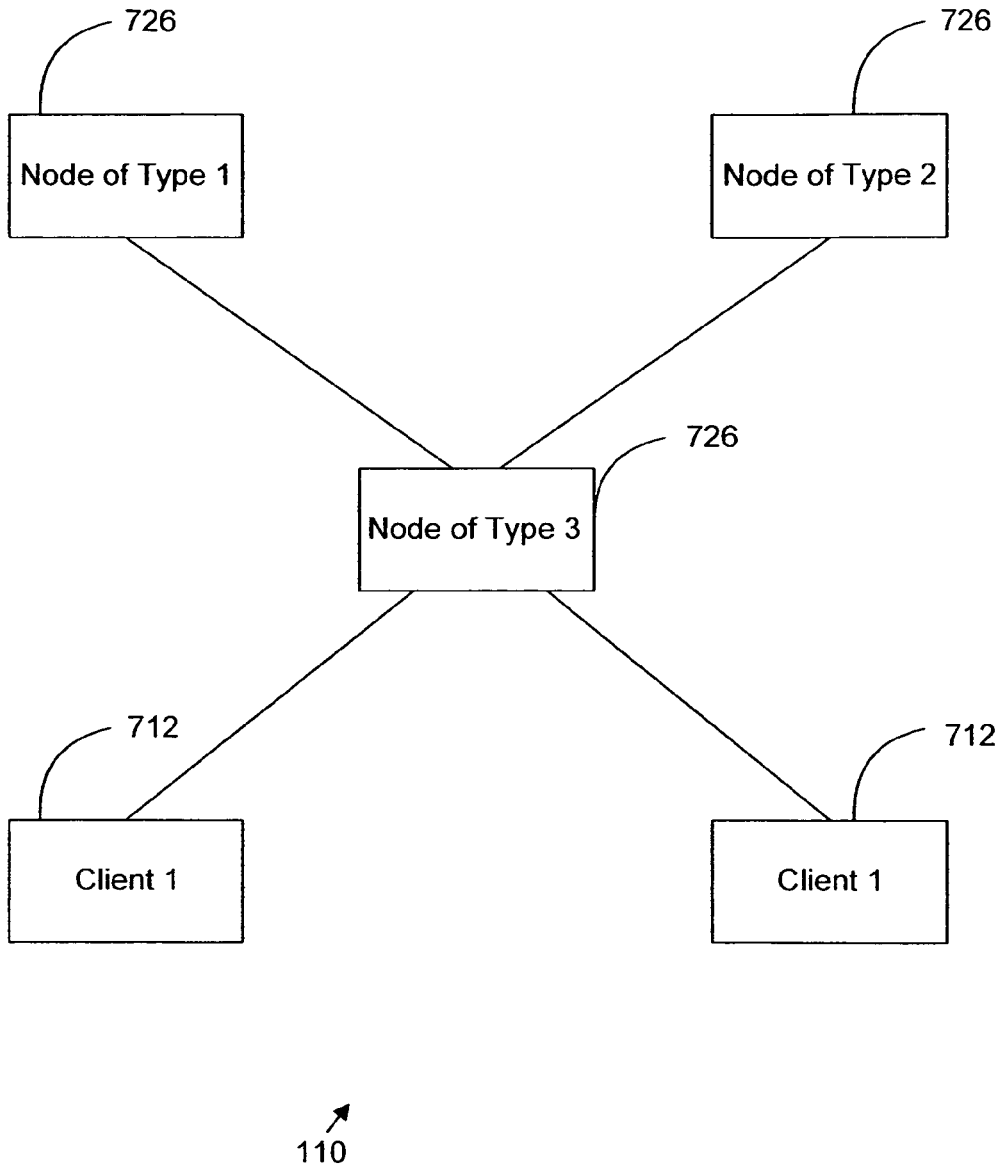
726
Node of Type 1
726
Node of Type 2
726
Node of Type 3
712
Client 1
712
Client 1
110

Fig. 12

ASYNCHRONOUS MESSAGING USING A NODE SPECIALIZATION ARCHITECTURE IN THE DYNAMIC ROUTING NETWORK

CROSS-REFERENCE TO RELATED APPLICATIONS

This application is a continuation of U.S. application Ser. No. 10/105,018 filed Mar. 21, 2002 (now U.S. Pat. No. 7,051,070), which is a continuation-in-part of U.S. application Ser. No. 10/017,182 filed Dec. 14, 2001 now U.S. Pat. No. 7,043,525, which claims the benefit of U.S. Provisional Application No. 60/256,613, filed Dec. 18, 2000, U.S. Provisional Application No. 60/276,847, filed Mar. 16, 2001, U.S. Provisional Application No. 60/278,303, filed Mar. 21, 2001, U.S. Provisional Application No. 60/279,608, filed Mar. 28, 2001, and U.S. Provisional Application No. 60/280,627, filed Mar. 29, 2001, all of which are hereby incorporated by reference herein.

BACKGROUND

1. Field of the Invention

This invention pertains in general to transferring information through digital networks and in particular to transferring information for remotely updating content at client devices through the digital networks.

2. Background Art

The Internet is a digital network of computers. An individual computer on the Internet is typically identified by an internet protocol (IP) address. A computer on the Internet sends a packet of information to another computer by routing the packet to a logical port at the destination computer's IP address. The destination computer interprets the packet according to one of several possible protocols determined by the port to which the packet was sent.

The World Wide Web (the "Web") is a collection of technology and content available on the Internet that allows the content to be routed from server computers to particular destination computers. The Web includes a large number of web pages residing on many different servers. Web pages contain one or more files, or references to one or more files, specifying instructions for presenting the web page and content, such as text, images, applets, video, and/or audio.

Web pages use a variety of definitional and programming languages to control how information is presented. The most fundamental of these is the Hypertext Markup Language (HTML). HTML uses a system of "tags" to specify how content should be displayed. Recent advances in HTML introduce "style sheets" which help separate content information from display information. HTML has also been modified and extended to provide new capabilities. For example, Extensible Markup Language (XML) adds semantic content to web pages. In addition, Dynamic HTML (DHTML) adds some dynamic content to web pages.

A web page may also include one or more programs for controlling how the web page is displayed. For example, JAVA® applets and JAVASCRIPT® scripts may be used to control the display of a web page. In addition, DHTML uses scripts to control the dynamic content. Thus, a web page designer can use applets and scripts to produce animation effects or modify the display based on user interaction. For example, the designer can write a script that changes the color of a piece of text when a user clicks on a button.

Devices that display/execute web pages are often called "client devices" or simply "clients." Client devices include personal computers, web-enabled set-top boxes and televi-

sions, cellular telephones, personal digital assistants and other handheld devices, and special-purpose web-browsing appliances. Client devices typically employ a program called a "web browser" for interpreting the HTML or other display instructions in the web page and displaying the content accordingly. Most web browsers include special functionality, such as a Java Virtual Machine, for executing JAVA® applets and/or other applets or scripts embedded in the web pages.

A client device specifies a web page or other document on the web using a Uniform Resource Locator (URL). A URL has the form "service://server/path/file." Here "service" refers to the protocol to be used, such as the file transfer protocol (FTP) or the hypertext transport protocol (HTTP). "Server" is the IP address of the server containing the page, and "path/file" specifies the particular web page on the server.

The Web suffers from a substantial limitation with respect to dynamically updating content in a web page at a client device. The Web's only mode of operation is for a client device to first request a page from a server and then for the server to send the requested page to the client device. Once the server delivers the page to the client, it typically terminates its connection to the client, and does not retain any information about the client or the page that was sent. For this reason, servers are typically "stateless." As a result, client devices drive and control the flow of information around the Web. While client-side control is appropriate in some situations, it does not permit efficient updating of data at the client devices. For example, if a web page contains information that may change, such as the score of a baseball game or a stock quote, the server has no way to inform the client devices that are viewing the page of the change. Instead, the client devices must ask the server for the updated information. However, the client devices do not know when the information on the web page has changed, and thus do not know to ask for the update.

There are some simple web programming techniques that attempt to update content on client device-side web pages. One approach that web designers use is to rely on the client devices to periodically re-request web pages. This updating can be performed as the result of user action (such as pressing the "refresh" button) or can be automated to occur on a particular schedule (such as by using the HTML Meta Refresh tag to cause the client device to request the page every 'X' seconds). Although this technique provides client devices with more up-to-date information, it is very wasteful of resources. In particular, the web server must resend the page even if nothing has changed, and, even when something has changed, it must resend the entire web page rather than just the updated information, which may be only a very small part of the page. Further, attempting to reduce unnecessary requests by decreasing the request rate results in decreasing the currency of the data. This is an unalterable trade off in a client-driven approach.

The performance of automatic refreshing can be improved somewhat by putting information that may change in a separate frame from information that is less likely to change, and only refreshing the separate frame. A few web designers even write custom JAVA applets to limit refreshing to individual components on a page, such as the score of a soccer game. A willingness to go to such effort illustrates the serious drain of resources caused by frequent refreshing. Nevertheless, even custom JAVA applets are not a meaningful attack on this problem. Custom applets require a large separate development effort for each item on each page that might need to be updated. More importantly, most custom applets still update content based upon client-driven requests, although it is possible to design an applet that accepts "pushed" messages. This

solution is not scalable to provide updated information for large numbers of client devices and for large numbers of web pages.

Therefore, there is a need in the art for an efficient way to provide dynamic content to a web page at a client device.

DISCLOSURE OF THE INVENTION

The above need is met by a dynamic content routing network that routes messages containing data for updating properties of live objects to clients displaying web pages or other representations of data containing the live objects. The web server that initially provides the web pages to the clients does not need to track which clients are currently displaying the live objects. Instead, the information provider or a dynamic content provider (generically referred to as an "input source") that provided the live object simply sends an update message to the routing network. This routing utilizes bandwidth efficiently because the update messages are provided to the clients only when the live objects change.

The routing network is adapted to selectively send messages to the nodes in the network. In one embodiment, a hierarchy of registrations is used. Each gateway in the routing network maintains the mappings between the live: objects and the nodes that have registered for the live objects. Each node in the routing network, in turn, maintains the mappings between the live objects and the clients that display them. An input source provides a message to a gateway in each cluster in the routing network. Each gateway forwards to each node only messages that reference the objects for which it has registered. Each node forwards to each client only messages that reference the objects for which it has registered. Adding node functionality to the gateway and client functionality to the node advantageously allows the routing network to decide which nodes should receive an update message. As a result, messages are sent to only nodes that have registered for the messages. Furthermore, each node receives all the messages that the clients connected to that node are interested in.

In another embodiment, all messages from an input source are assigned to one or more of N categories. Also, the nodes are assigned to one or more of M types, and mappings are created between message categories and node types. Each gateway keeps track of these mappings. When a gateway receives messages from input sources, the gateway identifies the categories of the messages and routes the messages to the nodes of the types to which the categories are mapped. To ensure that clients have access to the messages they need, clients are allowed to communicate with nodes of several types using client proxies connected between the clients and the nodes. There are at least two ways to implement the client proxy embodiment. The implementations differ primarily in where the client registration information is stored. In the first implementation, client registration information is stored at the nodes, and the client proxy merely passes messages among the clients and nodes. When a node directs a message to the client, the node passes the message to the client proxy along with a pointer to the client socket. When the client proxy receives the message, it simply pushes it to the client socket.

In the second implementation of the client proxy embodiment, the client proxy stores client registration information. In this embodiment, each node stores the object IDs registered by client proxies connected to the node. The client proxy registers with the nodes for all objects for which it needs to receive updates. The nodes receive the messages from the input sources, determine which client proxies have registered for the messages, and send the messages to the appropriate client proxies. The client proxies, in turn, transmit the messages to the clients that have registered for the messages.

The features and advantages described in this summary and the following detailed description are not all-inclusive, and particularly, many additional features and advantages will be apparent to one of ordinary skill in the art in view of the drawings, specification, and claims hereof.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. **1** is a high-level block diagram illustrating an environment containing a dynamic content routing network;

FIG. **2** is an interaction diagram illustrating interactions among a server, information provider, dynamic content provider, client, and routing network to update a property of a live object on a web page;

FIG. **3** is a high-level diagram graphically indicating the many-to-many mapping performed by the routing network;

FIG. **4** illustrates two different web pages containing sports scores;

FIG. **5** is a block diagram illustrating an input source and the tools available to it for generating the update messages;

FIG. **6** is a flow chart illustrating the steps performed by an embodiment of an activation module;

FIG. **7** is a block diagram illustrating a lower-level view of the routing network according to an embodiment of the present invention;

FIG. **8** is a flow chart illustrating steps performed by a gateway and a node in a cluster to perform object-based routing of a message received from an input source in an embodiment using a hierarchy of registries;

FIG. **9** is a block diagram illustrating a high-level view of the routing network in an embodiment adapted to use message categories, node types, and client proxies;

FIG. **10** is a flow chart illustrating steps performed by a gateway, a node that stores client registration information, and a pass-through client proxy to perform object-based routing of a message received from an input source;

FIG. **11** is a block diagram illustrating a high-level view of the routing network for an embodiment in which the client proxy stores the client registration information; and

FIG. **12** is a block diagram illustrating a high-level view of the routing network for an embodiment in which the nodes adopt client proxy functionality.

The figures depict an embodiment of the present invention for purposes of illustration only. One skilled in the art will readily recognize from the following description that alternative embodiments of the structures and methods illustrated herein may be employed without departing from the principles of the invention described herein.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

FIG. **1** is a high-level block diagram illustrating an environment **100** containing a dynamic content routing network **110** (hereafter referred to as the "routing network"). The environment **100** also contains a server **112** in communication with a client **114**, an information provider **108**, and a dynamic content provider **116**. Although a typical environment **100** will have hundreds of servers **112** and information providers **108**, thousands (or even millions) of clients **114**, and multiple dynamic content providers **116**, FIG. **1** illustrates only one of each of these entities in order to enhance the clarity of this description.

The server **112**, client **114**, information provider **108**, dynamic content provider **116**, and routing network **110** are

preferably in communication via conventional communications links **117** such as those comprising the Internet. The communications links **117** include known wired communications media, such as dedicated or shared data, cable television or telephone lines, and/or known wireless communications media, such as communications over the cellular telephone network using protocols such as the global system for mobile communications (GSM), code division multiple access (CDMA), time division multiple access (TDMA), etc.

In one embodiment, the entities may each be in communication with one or more Internet Service Providers (ISPs) (not shown) that provide each entity with access to other computers on the Internet. In addition, the server **112**, client **114**, information provider **108**, dynamic content provider **116**, and routing network **110** are preferably each identified by at least one Internet Protocol (IP) address such as "66.35.209.224." The IP address may also have one or more domain names associated with it, such as "bangnetworks.com." Alternative embodiments of the present invention may use alternative addressing schemes and/or naming conventions instead of, or in addition to, those described herein. For example, embodiments wherein one or more of the clients are cellular telephones or other portable devices may rely on different addressing schemes.

Preferably, the information provider **108** provides web pages or other representations of data to the server **112**. The web pages contain one or more "live objects," which are designated to be real-time dynamically-updateable objects. Each live object is identified by an object identifier, or object ID. Preferably, the server **112** provides the web pages **118** to multiple clients **114**. The clients **114** contact the routing network **110** and register for update messages for the object IDs on the web page. The routing network **110**, in turn, preferably maintains a registry indicating which clients have registered for which object IDs.

The information provider **108** and/or dynamic content provider **116** send update messages to the routing network **110**. These messages can be sent any time the information provider **108** or dynamic content provider **116** wants to update a property of a live object. Each update message preferably identifies a live object and contains data for updating a property of the identified live object. The routing network **110** accesses the registry and determines which clients have registered for the identified object. Then, the routing network **110** routes the update message to the appropriate clients. Upon receipt of an update message, the clients **114** update the specified property of the live object.

The routing network **110** provides an efficient one-to-many mapping of objects to clients (and by inference of information, a many-to-many mapping of information providers **108**/dynamic content providers **116** to clients) through object-based routing. Messages provided by the information provider **108** and/or dynamic content provider **116** to the routing network **110** are not routed to the clients **114** based entirely on a specified destination; more specifically, they are not routed based on the IP address of the client, as in conventional IP routing schemes. Instead, the messages are routed based on the live objects referenced by the message.

The mapping and object-based routing provided by the routing network **110** allow the information provider **108** and dynamic content provider **116** to update properties of live objects at a dynamically changing cross-section of clients in real-time, without requiring the information provider or dynamic content provider to track the clients or web pages being viewed by the clients. The clients **114**, in turn, do not need to have any a priori knowledge of object IDs—they "discover" which IDs they should register when they receives the web pages **118** from the server **112**.

Object-based routing also allows information providers **108** to dynamically update content on web pages without requiring the clients **114** to re-request the content, and without requiring the information providers **108** or servers **112** to maintain connections with the clients. In this manner, significantly more clients can receive updated content from a given information provider **108** than would be possible utilizing conventional client-side request-driven transmission control protocol/Internet Protocol (TCP/IP) connections between the clients and the server **112**.

Turning now to the individual entities illustrated in FIG. **1**, the server **112** is preferably a conventional computer system configured to act as a web server and serves web pages **118** and other data representations to clients **114**. The web pages **118** provided by the server **112** are associated with one or more information providers **108**.

An information provider **108** is an entity providing one or more web pages **118**, information contained in web pages, and/or other representations of data served by the server **112**. The information provider **108** preferably has a conventional computer system coupled to the Internet. In one embodiment, the server **112** is directly controlled by the information provider **108** (e.g., the server is physically located at the information provider and/or is dedicated to serving only the information provider's web pages). In this embodiment, the server **112** and information provider **108** can be treated as the same entity. In an alternative embodiment, the server **112** serves web pages from multiple information providers.

As is known in the art, the web pages **118** and other content on the server **112** are specified by uniform resource locators (URLs) having the form "service://server/path/web page." Typically, web pages **118** are obtained via the hypertext transport protocol (HTTP) and thus an exemplary URL for retrieving the web page "b1.html" from the web server having the domain name "www.bangnetworks.com" is "http://www.bangnetworks.com/news/b1.html."

As used herein, a "web page" is a block of data available from the server **112**. In the simplest case, a web page is a file written in the hypertext markup language (HTML). The web page may also contain or refer to one or more other blocks of data, such as other files, text, images, applets, video, and/or audio. In addition, the web page may contain instructions for presenting the web page and its content, such as HTML tags and style sheets. The instructions may also be in the Extensible Markup Language (XML), which is related to HTML and adds semantic content to web pages or the Dynamic HTML (DHTML), which adds some dynamic content to web pages. Additionally, the instructions may take the form of one or more programs such as JAVA® applets and JAVASCRIPT® and/or DHTML scripts.

As used herein, the phrase "web page" also refers to other representations of data served by the server **112** regardless of whether these data representations include characteristics of conventional web pages. These data representations include, for example, application programs and data intended for the web browser **120** or other application programs residing at the clients **114** or elsewhere, such as spreadsheet or textual (e.g., word processing) data, etc.

In a preferred embodiment, objects at the client, such as web pages and elements of web pages, can be designated as "live" by the information provider **108**. Properties of a live object can be dynamically updated in real-time at the client **114** by the information provider **108** or another entity acting on behalf of the information provider. As used herein, an "object" is any datum or data at the client **114** that can be

individually identified or accessed. Examples of objects include elements of web pages such as text characters and strings, images, frames, tables, audio, video, applets, scripts, HTML, XML, and other code forming the web page, variables and other information used by applets, scripts and/or code, URLs embedded in the web page, etc. Application and operating system constructs are also objects. For example, cells of spreadsheets, text in word processor documents, and title bars and messages displayed by the operating system or applications are objects. Preferably, multiple objects can be grouped together into a single, logical object. Thus, an object can be defined at any desired or useful level of granularity.

Since content on a web page is conceptualized and organized by "object," the present invention essentially abstracts web pages and web page content, and other modules and/or functionality at the client **114**, away from the HTML code or other conventional representation. This abstraction allows the information provider **108** to update a property of an object without concern for the location, display format, or other specifics of how the data is being represented at the client **114**.

Live objects have associated "properties" which include any modifiable data related to the object or referenced with respect to the object. The information provider **108** typically, but not necessarily, provides initial settings for the properties of live objects provided to the client **114**. The properties may or may not affect the visual representation of the object in the web page or other data representation. A property may affect an internal aspect of the object and, thus, a change to the property may not have any direct effect on a web page containing the object. For example, the property may affect whether particular aspects of the object are modifiable, how the object responds to user input or other stimuli, etc. Additionally, a property may also have a direct effect on how the object is displayed at the client **114**. For example, the property may affect the content, color, typeface, size, formatting, or other attribute of text, images, or other data displayed by the object. Other properties may occupy parts of the spectrum between having no effect on the visible representation of the object and having a direct effect on the visible representation of the object. For example, a web page showing scores of football games may include a list of games and the current scores of the games as of the time the server **112** serves the web page. The list of games, subset of games to be displayed, and the scores of the games can be designated as live objects (or properties of a single live object) and updated as necessary or desired.

A property can also preferably include instantiating an instance of the object or invoking functionality of the object. For example, a property of a browser window object may include functionality for instantiating another browser window. This function can be invoked as a logical change to a property of the object. The second browser window can be referenced through the original browser window (i.e., object) or designated as a new live object.

An information provider **108** or other entity preferably updates a live object at a client **114** via an update message. In general, an update message identifies the live object and, if necessary, the property of the live object, and contains data for updating the property. In one embodiment, the data may be the actual value for the property or executable code for causing the object's property to be updated. For example, the data may be a simple numerical or textual value, e.g., "4," to which the property should be set, and/or the data may be JAVASCRIPT® code or a call to a JAVASCRIPT® function at the client that effects the desired change to the property of the object.

The update message preferably implicitly or explicitly identifies a handler at the client **114** for use in updating the live object's property. In one embodiment, the client **114** utilizes a default handler when the message implicitly specifies the handler (e.g. when the message does not identify a specific handler). In one embodiment, if the update message specifies the actual value for the property, a default handler generates JAVASCRIPT® code for changing the property to the specified value. If the data in the update message are JAVASCRIPT® code, the default handler does not perform any processing of the code. In either case, the default handlers preferably use LiveConnect to execute the JAVASCRIPT® code in a Java Virtual Machine (JVM) **122** at the client **114** and thereby update the property of the live object.

For certain objects and/or data types, the default handlers are not appropriate. In these cases, the message preferably explicitly identifies a handler for performing the update. For example, the message may explicitly specify a function to call on the data or the message may explicitly identify the environment in which the data should be executed. For example, the data in the update message may include code for execution by a software "plug-in" such as MACROMEDIA FLASH® and the message may explicitly identify FLASH as the handler.

The information provider **108** preferably designates an object as "live" by including a unique identifier for the object, the object ID, in the web page or other data representation provided to the client **114**. In one embodiment, the information provider **108** encodes the object ID in an object's corresponding HTML "ID" attribute using the following HTML expression:

ID="elementIdentifier,"

where "elementIdentifier" is the object ID and is preferably a string. The string can encode any information desired by the information provider **108** or other entity establishing the object ID and in one embodiment is a simple textual and/or numeric identifier. In one embodiment, the information provider **108** begins the object ID with a predefined token, such as "Bang$," in order to distinguish live objects from other objects that happen to have defined ID attributes. For example, an object can have the object ID "Bang$elementIdentifier."

In the preferred embodiment, each information provider **108** optionally encodes a unique information provider ID in its object IDs in order to prevent naming collisions between the object IDs of different information providers. In one embodiment, the information provider ID is a textual and/or numeric identifier. The information provider **108** may specify the information provider ID and the object ID as part of a hierarchical namespace. For example, in one embodiment objects are named as follows: "$namespace1$ [namespace2$ . . . $namespaceN$]objectId," where "$namespace1" is the information provider ID and the "$" operates as the name separator and defines additional optional levels of a namespace hierarchy. One embodiment of the system **100** supports typical directory services functionality. For example, two dollar sign characters appearing together, "$$," refers to the top level of the namespace hierarchy.

Thus, the object ID for a live object is preferably formed from a combination of the predefined token, the information provider ID namespace, and a value assigned by the information provider **108**. For example, the object ID for a live object representing the real time price of a stock having the symbol "BANG" might be: "Bang$$informationProviderID-$equities$realtime$bang." In this example, "Bang$" is the predefined token that signifies a live object, "$informationProviderID" is the ID identifying the information provider,

"$equities$realtime$" defines levels of a namespace hierarchy, and "bang" identifies the specific object.

In some embodiments and situations, the object ID utilizes relative names. For example, an information provider **108** referring to its own object IDs is implicitly in its own namespace. Accordingly, the information provider **108** does not need to include the information Provider ID in the object IDs it utilizes internally. In one embodiment, the information provider ID is not explicitly encoded into the object ID. Instead, the information provider ID is encoded elsewhere in the web page in order to provide scope to the page's object IDs.

In one embodiment, the object ID identifies a point (i.e., a node in a tree) in a Document Object Model (DOM) representation of a web page or other document at the client **114**. The DOM is a platform- and language-neutral interface that represents a document as a hierarchy of objects. The DOM also provides an interface that allows programs and scripts to dynamically access and update properties of the objects. Object properties can be inherited by descendent objects.

In this embodiment, the client **114** preferably executes an update message in the context of the specified point in the DOM representation. The update may specify a change to a property of the object at the identified point. The update also may specify a change to a parent or descendent of the object at the identified point. In each case, the update is executed relative to the specified point in the DOM representation. In one embodiment, points in the DOM representation specify how to update properties of live objects located at those points. Thus, the same update may be interpreted differently depending upon the identified live object's location in the DOM representation.

For example, assume there is an object in the DOM representation identified as "window.document.frame[3].ObjectID." Also assume that the object has an "innerText" property located at "window.document.frame[3].ObjectID.innerText" that specifies the text displayed by the object. An update message can change the text displayed by the object by specifying "ObjectID" and the new value for the innerText property.

An advantage of utilizing object IDs to specify objects is that the information provider **108** or other entity providing the update message can access and change properties of objects without knowing the object's actual location in the DOM representation. Indeed, the object may be in different locations in different DOM representations and/or in multiple locations in the same DOM representation. In any of these cases, the update message will change the specified properties of all of the objects having the given object ID.

Depending upon the particular embodiment of the environment **100**, the information provider **108** and/or the dynamic content provider **116** provides update messages to the routing network **110**. The dynamic content provider **116** is preferably a conventional computer system operated by an entity that provides real-time information, such as stock prices and/or sports scores. In one embodiment, the information provider **108** receives updated properties for the live objects from the dynamic content provider **116** or another source (or generates the updated properties internally). Then, the information provider **108** sends an update message specifying the object ID and the change to the object property to the routing network **110**. In this embodiment, the dynamic content provider **116** may be absent from the environment **100**.

In another embodiment, the dynamic content provider **116** provides the object IDs for live objects to one or more information providers **108** and the information providers **108** distribute the live objects to the clients **114**. Then, the dynamic content provider **116** sends messages specifying the changes to the properties of the live objects to the routing network **110**. For example, the dynamic content provider **116** distributes an object ID associated with the score of a particular baseball game to the information providers **108**. Then, the dynamic content provider **116** sends a message specifying the object ID and an update to a property of the object that controls the displayed score of the particular baseball game to the routing network **110**. These two embodiments are not mutually exclusive and, therefore, some updates may be provided to the routing network **110** by the information provider **108** while others are provided by the dynamic content provider **116**.

The client **114** is a device that retrieves web pages **118** and/or other information from the server **112**. In one embodiment, the client **114** is a conventional personal computer used by a person to access information on the Internet. In alternative embodiments, the client **114** is a different consumer electronic device having Internet connectivity, such as an Internet-enabled television, a cellular telephone, a personal digital assistant (PDA), a web browsing appliance, etc. The client **114** preferably, but not necessarily, has an associated display device.

The client **114** preferably executes a web browser **120**, such as MICROSOFT INTERNET EXPLORER®, for retrieving web pages and displaying them on the display device. In embodiments where the client receives data representations from the server **112** other than conventional web pages, the web browser **120** does not necessarily share similarities with conventional web browsers. Preferably, the web browser **120** contains a JVM **122** for executing JAVA® applets and/or scripts. The web browser **120** also preferably contains Dynamic HTML capabilities, such as support for JAVASCRIPT® (or another scripting language, such as VBScript) and the Document Object Model (DOM), and enables communications between JAVA® and the scripting languages. In one embodiment, the web browser **120** supports the LiveConnect standard for enabling communication between JAVA® applets and scripts written in the supported scripting languages. The web browser **120** can also be extended through software plug-ins such as MACROMEDIA FLASH®, REAL NETWORKS REALPLAYER®, and/or APPLE QUICKTIME®. In alternative embodiments, the functionality of the JVM **122** and/or other aspects of the web browser **120** are provided by one or more other functional units within the client **114**. The term "module" is used herein to refer to software computer program code and/or any hardware or circuitry utilized to provide the functionality attributed to the module. The web browser **120** and JVM **122** are examples of modules in the client **114**.

In some embodiments, the client **114** does not necessarily have a display device, web browser **120** and/or other components associated with a typical consumer device. The client **114**, for example, may be a dedicated purpose device having certain aspects of web connectivity such as an embedded HTTP client in a web-enabled appliance or in a controller for an automobile, audio-visual equipment, or some other device.

A web page **118** provided from the server **112** to the client **114** preferably includes instructions for enabling the live objects on the web page. The instructions cause the client **114** to automatically and transparently (i.e., without user interaction) contact the routing network **110** and download an activation module **124** for activating the live objects. In one embodiment, the instructions comprise a URL specifying the location of the activation module **124** at the routing network **110**. In an alternative embodiment, the client **114** obtains the activation module **124** from the server **112** or another source.

The activation module **124** preferably contains JAVA® instructions for execution by the JVM **122**. However, alternative embodiments of the module **124** may encode the instructions in the web page **118** and/or the activation module **124** using different languages and/or techniques. For example, the instructions and/or activation module **124** can be embedded in the web browser **120** or operating system, either as native code or as plug-ins. In these alternative embodiments, the web browser **120** does not have to download the activation module **124** from an external source.

The activation module **124** preferably registers object IDs from the web page **118** downloaded by the client **114** with the routing network **110** and updates the live objects in response to update messages received from the network. The routing network **110** records the registrations in the registry **125**. The client's registrations preferably remain in effect as long as the client is displaying the associated web page **118**, although other embodiments of the system **100** may use different criteria for determining when to terminate the client's registrations.

FIG. **2** is an interaction diagram illustrating interactions among the server **112**, information provider **108**/dynamic content provider **116** (generically referred to as an "input source **210**"), client **114**, and the routing network **110** to update a property of a live object. Initially, the client **114** sends **212** a web page request to the server **112**. In response, the server **112** provides **214** to the client **114** the web page containing or otherwise identifying the one or more live objects. Instructions encoded in the web page preferably cause the client **114** to transparently request **216** the activation module **124** from the routing network **110**. In response, the routing network **110** sends **218** the activation module **124**. The client **114** executes **220** the activation module **124**, which identifies the object IDs of the live objects at the client and registers **222** the object IDs with the routing network **110**. The routing network **110** updates **223** its registry to identify the object IDs for which the client **114** has registered.

At some point, the input source **210** sends **224** an update message to the routing network **110** in order to change a property of a live object at the client **114**. In one embodiment, the message from the input source **210** to the routing network **110** contains only a single object ID and an update to a property of the identified object. In another embodiment, the message contains multiple object IDs and the corresponding property updates. In this latter embodiment, the message may have an associated "Batch ID" that identifies the message as having multiple object IDs and updates. Preferably, the information provider **108** can include a batch ID in a web page **118** in the same manner as including an object ID. Likewise, the client **114** can preferably register for a batch ID with the routing network **110** in the same manner as an object ID. In fact, the batch ID can be the same as the object ID so that the client **114** registers for both batch and non-batch messages by registering one ID. Alternatively, separate procedures can be established for registering batch messages. The client **114** preferably processes the component messages of a batch as if each message were delivered separately.

The routing network **110**, in turn, routes **226** the message to each client **114** that has registered for the specified object ID, preferably by utilizing standard Internet communications protocols, such as IP addresses, etc. The activation module **124** at the client **114** processes the message and updates **228** the property of the identified live object. If live objects having the same object ID appear in multiple locations at the client **114** (e.g., at multiple locations on a web page being displayed at the client), the activation module **124** preferably updates each of the live objects having the specified ID. As a result, the routing network **110** allows live objects at the client **114** to be dynamically updated. Preferably, this routing and updating happens quickly enough to be considered "real-time" for the purposes of the input source **210**.

This update process, indicated within the dashed box **230** in FIG. **2**, can repeat an indefinite number of times and is fully asynchronous as to the information provider **210** and client **114**. For example, the input source **210** may send regular update messages to the routing network **110** as the score of a sporting event changes or a stock price fluctuates, but may stop sending update messages once the sporting event ends or stock market closes. When the client **114** ends the display of a web page containing the live object, or otherwise no longer desires to receive update messages, the client preferably closes **232** the connection with the routing network **110**. The routing network **110**, in turn, updates **234** the registry **125** to remove the client's object registrations. In another embodiment, the client **114** sends messages to the routing network **110** that selectively register and/or de-register the client from one or more objects yet leaves the connection open in order to receive update messages pertaining to other objects.

FIG. **3** is a high-level diagram graphically indicating the many-to-many mapping performed by the routing network **110**. Multiple input sources (labeled **210**A-C) send update messages to the routing network **110**. Each update message preferably specifies at least one object ID and an update to a property of the identified object. The routing network **110**, in turn, selectively routes the update messages to the clients **114** that have registered for the given object ID from the given input source **210**. In FIG. **3**, assume for example that clients **312**A and **312**B have registered for a given object ID while the other clients have not registered for the object ID. Accordingly, the routing network **110** routes the update message to clients **3 12**A and **312**B, but does not route the message to clients **312**C-**312**H.

FIG. **4** illustrates an example of the capabilities of the dynamic content routing network **110**. FIG. **4** illustrates two different web pages **410**, **412** containing sports scores. Although the web pages are formatted differently, each page contains the same scores for two professional football games and two professional baseball games. Web page **410** contains all four games under the heading "Local Sports Scores" while web page **412** contains the baseball games under the heading "Baseball Scores" and the football games under the heading "Football Scores."

There are various ways to internally represent the games and scores in the web pages using live objects. In one embodiment, a "game" object is defined having properties for the two teams involved in the game and the score associated with each team. The game object is placed at a selected position in the web page and the properties of the object cause the information about the game to be displayed on the page. In another embodiment, "team" and "score" objects are defined, with the team object having a property defining the name of a team and the score object having a property defining a score. In this second embodiment, the team and score objects are placed at selected locations on the page so that the proper teams and scores are aligned when the page is rendered. In yet another embodiment, an object is defined having properties for the name of one team and a score associated with that team. Then, pairs of the objects are placed in the page in the proper alignment to indicate the games and scores. In another embodiment, an object is defined having properties specifying names of two teams and a separate object is defined having properties specifying two scores. In this last embodiment, the two objects are placed in the page so that the names

of the teams align with the associated scores. Obviously, additional variations of these representations are possible.

Assume for the example of FIG. **4** that the names of teams in a game are specified by a "names" object having properties for the two team names and the scores in the game are specified by a "scores" object having properties for two scores. In web page **410**, a names object **414** having properties set to identify the "SF 49ers" and the "STL Rams" is located directly under the "Local Sports Scores" heading. A scores object **416** having a property set to identify the score of the game as "42" to "7" is directly to the right of the names object **414**. In web page **412**, the properties of the second names object **418** identify the same game using slightly different terminology: "SF" and "STL." However, this names object **418** is aligned with the same scores object **416** as is utilized in web page **410**.

Thus, the same scores object **416** is utilized in different positions in each web page **410**, **412**. In order to update the score of the San Francisco 49ers vs. St. Louis Rams football game on both web pages, the input source **210** simply sends an update message to the routing network **110** specifying the object ID for the scores object **416** and the update to the score property. The routing network **110** routes the update message to the appropriate clients **114**, and the clients update the appropriate score regardless of the particular page layout.

The input source **210**, i.e., the information provider **108** and/or dynamic content provider **116** can use a variety of tools to generate the update messages. FIG. **5** is a block diagram illustrating an input source **210** and the tools available to it for generating the update messages. Other tools can be utilized in addition to or instead of the ones described herein.

Preferably, the tools allow the input source **210** to access an application programming interface (API) provided by the routing network **110** for accepting messages. In one embodiment, the messages sent by the input source **210** are in the same format as utilized by the activation module **124** at the client **114**. In an alternative embodiment, the messages provided to the routing network **110** are in a different format and the routing network translates the messages into the format utilized by the activation module **124**.

In one embodiment, the input source **210** utilizes a data pump module **510** to access the API. The data pump module **510** reads an extensible markup language (XML) file containing one or more object IDs and the new values for the identified objects at regular intervals and automatically generates API calls that send messages representing changes to object properties to the routing network **110**. In another embodiment, the data pump module **510** is event-driven and reads the XML file in response to a change in the file or some other occurrence.

In another embodiment, the input source **210** utilizes a director console module **512** to access the API. Preferably, the director console module **512** presents an administrator with a graphical interface displaying the contents of the web page **118**. For example, the administrator may use the director console **512** to edit textual data, images, and/or any objects or properties of objects on the web page. After editing, the administrator uses a "send update" button or similar technique to cause the director console module **512** to send messages for the changed objects and properties to the routing network **110** via the API.

In another embodiment, the information provider **108** and dynamic content provider **116** work together as the input source **210** by using a content management system module **514** to access the API. Preferably, the content management system module **514** resides at the information provider **108** and receives object property updates from the dynamic content provider **116**. The content management system module **514** preferably updates the properties of the live objects in the web page **118** stored at the server **112** and also sends messages for the changed properties to the routing network **110**. In this manner, the web page **118** at the server **112** and the web page displayed at the client **114** are updated almost simultaneously. In one embodiment, the dynamic content provider **116** sends the update messages to the routing network **110** instead of to the information provider **108**. Embodiments of the system **100** can also utilize any combination of the content management techniques described herein.

For example, the tools described above can generate a message having the following code for updating the text displayed by a score object to "2":

```
LiveObject score=new LiveObject("Bang$homeScoreID");
score.setProperty("innerText", "2").
```

This code sets the innerText property of the object having object ID "Bang$homeScoreID" to "2." The tools use the API to pass this message to the routing network **110**.

Turning now to the actions performed at the client **114**, FIG. **6** is a flow chart illustrating the steps performed by an embodiment of the activation module **124**. Those of skill in the art will recognize that different embodiments may perform the steps of FIG. **6** in different orders. The activation module **124** generally performs three functions: register object IDs with the routing network **110**, handle messages received by the client **114** from the network in order to update the properties of live objects, and control communications between the client and the network.

In order to register object IDs, the activation module **124** preferably parses **610** the web page **118** received from the server **112** and identifies the object IDs of the live objects. In an alternative embodiment, the activation module **124** identifies only a subset of the object IDs, such as the IDs of only live objects that are currently being displayed by the web browser **120**. Alternatively, a list of object IDs may be pre-encoded in the web page in addition to the objects themselves, thereby enabling easy identification by the activation module **124**. In yet another embodiment, a user of the client **114** selects the object IDs to register.

The activation module **124** preferably opens **612** a connection between the client **114** and the routing network **110**. The activation module **124** can open **612** this connection before or after the activation module receives and/or parses the web page **118**. In some cases, the client **114** is located behind a firewall that puts a restriction on the types of connection requests the client can make. A firewall might, for example, block all non-HTTP traffic. For this reason, the activation module **124** preferably wraps the connection request in an HTTP header in order to get the request to the routing network **110** through the firewall.

The activation module **124** uses the connection between the client **114** and routing network **110** to register **614** the object IDs by communicating to the routing network **116** a vector (e.g., a list or array) containing the identified object IDs. In order to accomplish this task through the firewall, the activation module **124** preferably puts the vector into a string, referred to as "object data," and then preferably creates an HTTP message to communicate the object data to the routing network **110**. A schematic example is as follows:

```
POST/HTTP/1.\r\n
Content-Length: <length of object data>\r\n
\r\n
<object data>
```

where <object data> is the object ID list. When the routing network **110** receives such an HTTP request, it extracts the object data and updates the registry **125** to indicate that the client **114** has registered for the identified objects.

If the web browser **120** loads **616** a new page, or otherwise terminates display of the objects on the initial page, the activation module **124** associated with the initial web page preferably terminates **618** the client's connection with the routing network **110**. Those of skill in the art will recognize that this termination **618** can occur asynchronously with the other steps illustrated in FIG. **6**. Thus, the location of steps **616** and **618** represents only one possible place in the sequence of steps where the termination may occur.

If the connection is not terminated, the activation module **124** preferably waits until it receives **618** a message from the routing network **110** specifying an object ID and an update to a property of the identified object. In one embodiment, this message is received as HTTP data. Upon receipt of the message, the activation module **124** preferably extracts **620** the object ID and update from the HTTP data. Then, the activation module **124** updates **622** the property of the identified object, or causes the object to be updated, as specified by the message.

The sequence of receiving messages **618**, extracting data **620**, and updating objects **622** is preferably repeated until a new page is loaded **616** or the connection with the routing network **110** is otherwise terminated. Although not shown in FIG. **6**, in certain circumstances, such as when a user action with respect to the web page **118** activates a new live object, the activation module **124** may register new object IDs with the routing network **110** without first downloading and parsing a new page. In one embodiment, if the newly-loaded page contains live objects, then the process of downloading the activation module **124** and updating the objects as described by FIG. **6** is repeated. In an alternative embodiment, the activation module **124** remains active at the client **114** and, therefore, the client does not re-download the activation module from the routing network **110**. Instead, the already-present activation module **124** performs the live-enabling process on the new page.

FIG. **7** is a block diagram illustrating a lower-level view of the routing network **110** according to one embodiment of the present invention. FIG. **7** illustrates multiple input sources (labeled **710**A-D) representative of sources providing messages to the routing network **110**, such as an information provider **710**A and a dynamic content provider **710**B. FIG. **7** also illustrates multiple clients (labeled **712**A-F) representative of the many clients in communication with the routing network **110** at any given instant.

Internally, the routing network **110** is preferably divided into one or more clusters **714**. In FIG. **7**, the routing network **110** has three clusters **714**A, **714**B, **714**C, although the number of clusters can vary depending upon the processing needs of the network. An input-side global load balancer **716** preferably routes messages from the input sources **710** to the clusters **714**. Similarly, a client-side global load balancer **718** preferably routes connection requests from the clients **712** to the clusters **714**. The load balancers **716**, **718** are designed to ensure that load is distributed among the clusters **714** according to a predetermined heuristic. For example, the load may be distributed evenly among the clusters **714** or a more powerful cluster may be distributed a majority of the load. In one embodiment, one load balancer performs the functions of the input-side **716** and client-side **718** load balancers and utilizes conventional Domain Name System-(DNS-) based load balancing.

Each cluster **714**, of which cluster **714**A is representative, preferably contains an input-side cluster load balancer **720**A and a client-side cluster load balancer **722**A. The cluster load balancers **720**A, **722**A function similarly to the corresponding global load balancers **716**, **718** in that the input-side cluster load balancer **720**A balances and routes incoming messages among one or more gateways **724**A and the client-side cluster load balancer **722**A balances and routes incoming connection requests among one or more nodes **726**A and application servers **728**A. The gateways **724**A are connected to the nodes **726**A. In one embodiment every gateway **724**A is connected to every node **726**A and in another embodiment certain gateways are connected to only certain nodes.

Preferably, the routing network **110** utilizes conventional single-processor computer systems executing the Linux operating system (OS). Preferably, each component of the routing network **110** is implemented by a separate, dedicated computer system in order to enable the separate optimization of the components. The input/output (I/O) functionality of the OS is preferably enhanced through the use of a non-blocking OS package such as NBIO available from the University of California, Berkeley, Calif. Based on the assumption that connections with the nodes **728** are long-lived, the OS is preferably configured to not allocate resources toward monitoring idle connections. Instead, the well-known/dev/poll patch is preferably applied to the OS in order to provide advanced socket polling capabilities.

Those skilled in the art will recognize that there are many ways to use the functionality of the routing network **110** to route update messages to clients **710**. For example, in one embodiment, every message is distributed to every node **726**. In another embodiment, the routing network **110** selectively sends messages to the nodes **726** in the routing network **110**. Selectively sending messages to the nodes in the routing network presents at least the two difficulties. First, for a given message, a decision needs to be made as to which nodes should receive it. Second, a client must receive all messages in which it is interested. There are at least two approaches meeting these difficulties. The first approach uses a hierarchy of registries at the gateways and nodes to respectively keep track which messages to send to the nodes and clients. The second approach assigns messages to one or more categories, assigns nodes to one or more types, and maintains mappings between categories and types. This latter approach also uses client proxies to allow clients **712** to communicate with multiple nodes of different types.

FIG. **7** illustrates the embodiment using the hierarchy of registries. A node **726** registers with each gateway **724** in all of the clusters and indicates which messages it needs and the clients do the same with the nodes. Each gateway **724** preferably maintains a registry **734** containing the object IDs registered by the nodes **726** connected to the gateway **724**. In one embodiment, the gateway registry **734** associates each object ID with a linked list containing one entry for each node **726** that has registered for that object ID. In another embodiment, the gateway registry **734** is a hash table containing the object ID registered by the nodes **726** connected to the gateway **724**. A node **726** preferably maintains a node registry **732** containing the object IDs registered by clients **712** connected to the node. The gateways **724** in each cluster **714** receive the messages from the input sources **710** and direct the messages to the appropriate node or nodes **726**. The nodes **726** preferably transmit messages received from the gateways **724** to the clients **712** that have registered in the object IDs identified by the messages.

In one embodiment, the node registry **732** associates each object ID with a linked list containing one entry for each

client **712** that has registered for that object ID. In another embodiment, the node registry **732** is a hash table containing the object ID registered by the clients **712** connected to the nodes **726**. Each entry in the linked list or hash table preferably contains a pointer to a socket representing the connection to the corresponding client **712**. As is known in the art, the pointer to the socket, typically called a "file descriptor," represents an address to which the node can write in order to send the message to the corresponding client. Gateways **724** can also use file descriptors in this manner to store node addresses. Alternative embodiments of the present invention utilize other data structures in addition to, or instead of, the hash table and linked list, and/or may utilize different data within the data structures.

Preferably, the node **726** adds an entry to its registry **732** every time a client **712** registers an interest in an object and deletes the corresponding entry from the registry when the client **712** disconnects from the node or otherwise indicates that it is no longer interested in a particular object. If the node **726** determines that the client **712** registered for an object ID that was not previously registered on that node, the node preferably registers that object ID with the gateways **724** to which it is connected. Similarly, if the node **726** determines that the client **712** deregistered an object ID for which it was the last interested client, the node **726** deregisters that object ID with the gateways **724** to which it is connected. The gateways **724** update their registries **734** in response to the communication from the node **726**.

In alternative embodiments of the present invention, when the client **712** disconnects from the node **726** or otherwise indicates that it is no longer interested in a particular object, the node **726** waits for a period of time or until some event occurs before deregistering the object ID with the gateway **724**. For example, the node **726** could wait until it receives a message associated with that object ID to do the deregistration. Alternatively, the wait time can be fixed, random, or based on the frequency of registrations for that object ID, or registrations for the same input source **710** as the object being deregistered. This latter approach keeps the object IDs from more frequently used input sources **710** registered for a longer period of time. This waiting advantageously reduces the number of registration changes required between nodes and gateways in cases where it is likely that another client will soon register for the same object ID.

Since a gateway **724** does not control the rate at which it receives messages from input sources **710**, it is possible for the gateway to receive messages faster than it can process them (i.e., send the messages to the nodes). Therefore, each gateway **724** preferably maintains a queue **730** of messages that have been received but not yet processed in order to avoid losing messages. In one embodiment, the gateway **724** drops messages if the queue **730** becomes too long. In another embodiment, the gateway **724** utilizes priorities assigned to certain messages or input sources to determine which messages to drop.

The application server **728** within each node **714** preferably serves the activation module **124** to the clients **712** in response to client requests. In addition, the application server **728** serves any other modules that may be required or desired to support the environment **100**. In an alternative embodiment of the routing network, a single application server **728** fulfills all of the client requests. This application server **728** may be within a certain cluster **714** or independent of the clusters. However, this single-application-server embodiment is less desirable because it lacks redundancy.

FIG. **8** is a flow chart illustrating steps performed by a gateway **724** and a node **726** in a cluster **714** to perform object-based routing of a message received from an input source **710** in the embodiment using a hierarchy of registries. Initially, the gateway **724** receives **810** the message from the input source **710**. The gateway **724** extracts **820** the object ID from the message. The gateway **724** examines its registry **734** to determine the nodes **726** that have registered in the object ID. The gateway **724** transmits **840** the message to each of the registered nodes **726**. Each node **726** that receives the message uses its registry **732** to determine **850** which clients **712** have registered for the message. Each node **726** then forwards **860** the message to the registered clients **712**.

Adding node functionality to the gateway and client functionality to the node advantageously allows the routing network **110** to solve the difficulties identified above. For example, it allows the routing network **110** to decide which nodes should receive an updated message. As a result, messages are sent to only nodes that have registered for the message. Further, no matter which node a client connects to, that node will receive all messages that the client wants. One skilled in the art would understand that while the present invention allows registration at both a gateway level and node level, registration could be extended to any number of levels.

In the second approach for meeting the difficulties described above, all messages in the routing network **110** are assigned to one or more of N categories, and all of the nodes are assigned to one or more of M types. Mappings are created that specify which categories of messages are forwarded to which types of nodes. The mappings allow control over the amount of traffic processed by the nodes.

FIG. **9** is a block diagram illustrating a high-level view of the routing network **110** of FIG. **7** in an embodiment adapted to use the approach having message categories and node types. Although FIG. **9** illustrates only two node types—nodes of type **1** **726** and nodes of type **2** **726**—for purposes of simplicity, this embodiment of the network **110** can have any number of node types. FIG. **9** also illustrates a client proxy **740** for reasons described below. However, an embodiment of the present invention using message categories and node types does not necessarily utilize a client proxy **740**.

There are many different possible mappings between message categories and node types. In the simple case, there is one-to-one mapping between message categories and node types. For example, if the message is of category **1**, it is forwarded to the nodes of type **1**. In more complicated mappings, messages of one category are mapped to nodes of multiple types. For example, messages of category **1** are mapped to nodes of types **1**, **2**, and **3**, whereas messages of category **2** are mapped to nodes of types **2**, **3** and **4**. In short, any possible mapping of message categories to node types is possible and the number of message categories, N, does not have to be the same as the number of node types, M.

There are multiple ways to assign the messages into categories. One way is to assign all messages from a given input source **710** into a certain category. Another way is to explicitly specify the category in the object ID for the message. Yet another way is to utilize a hashing function or lookup table to partition messages into categories based on object IDs or other values. For example, in one embodiment a hash function is applied to the object ID to generate an integer between **1** and N, and this integer is the message category.

Nodes are preferably assigned to types based on information stored in the gateways **724**. In one embodiment, each gateway **724** holds a lookup table or other data structure that specifies the types to which each node is assigned. The lookup table also preferably stores the mappings between message categories and node types. When a gateway **724** receives a message from an input source, the gateway preferably deter-

mines the category of the message using one of the techniques described above. Then, the gateway **724** determines the node type (or types) to which the message category maps, and determines which nodes are of the given type. The gateway **724** routes the message to the appropriate nodes.

In one embodiment, the gateway **724** uses a combination of multiple techniques to determine the message categories, node types, and/or mappings. For example, a lookup table can be used to encode a priori knowledge about categories, types, and/or mappings and a hash table can be used to route messages for which there is no a priori knowledge. Continuing this example, assume that certain messages are assigned to a given category based on a table lookup, while other messages are assigned to categories based on a hashing function. In this example, the gateway **724** looks up the object ID (or other information, such as an input source ID) of an arriving message in a lookup table to determine if it has a specified category. If the object ID is stored in the lookup table, the gateway **724** determines the mappings for the category and routes the messages to the nodes of the appropriate types. If the object ID is not stored in the lookup table, the gateway **724** utilizes a hash function on the object ID (or other information) to determine the message category.

Message categorization advantageously allows the routing network **110** to decide which nodes get which messages. However, message categorization does not ensure that clients **712** have access to the messages they need. Assume a client **712** connects to a node of type **1** and the client **712** wants to receive messages of category **3**. If there is a simple one-to-one mapping of message categories to node types, nodes of type **1** will never receive messages of category **3**, and neither will the client **712**. To ensure that clients have access to the messages they need, the embodiment shown in FIG. **9** uses client proxies to allow clients **712** to communicate with multiple nodes of different types. As shown in FIG. **9**, the clients **712** connect to a client proxy **740** instead of the nodes **726**. Each client proxy **740** is connected to at least one node of each type. For example, as shown in FIG. **9**, the client proxy **740** connects to a node of type **1 726** and a node of type **2 726**.

There are at least two ways to implement the embodiment that uses client proxies to ensure that clients connect to multiple nodes of different types. The implementations vary primarily in where the client registration information is stored. In the first variation, client registration information is stored at the nodes, and the client proxy is adapted to simply pass update messages and registration information among the clients and nodes. In the second variation, the client proxy **740** is responsible for keeping track of client registrations.

FIG. **9** illustrates the variation where the nodes store the client registration information, and the client proxy **740** is adapted to pass messages and registration information among the clients and nodes. When a client sends registration information to the client proxy **740**, the client proxy passes it to the appropriate node or nodes. In one embodiment, the client proxy **740** sends client registration information to at least one node of every type, and the nodes ignore irrelevant registration information (e.g., registrations for messages of categories not handled by the node). In another embodiment, the client proxy **740** itself analyzes the registrations, and passes only registrations relevant for a given node type to a node of that type. In a third embodiment, the client contains functionality for determining which node types handle which registration requests, and the client tells the proxy **740** to which node types to pass the registration information.

The client proxy **740** preferably stores an identifier of its connection to each client **712**, such as a pointer to the socket for the client connection, and sends this identifier to the nodes along with the registration information. Each node **726**, in turn, maintains a registry **732** storing the relevant registration information it receives from the client proxy **740**. Preferably, at most one node of each type will contain registration information for a given client for messages of a category mapped to the node type. In addition, a node preferably does not store registration information for messages of categories not mapped to its node type.

The registry **732** for each node **726** preferably indicates the object IDs registered by the clients, the client proxy **740** to which the client is connected (e.g., a pointer to the socket at the node to which the client proxy is connected), and the identifier indicating the client's connection to the client proxy (e.g., the pointer to the socket at the client proxy to which the client is connected). When the node **726** receives an update message, it uses the registry to identify the clients to which it should forward the message, the client proxies to which the clients are connected, and the specific connections between the client proxies and the clients. The node routes the update message to the identified client proxies and includes the identifier (e.g., the pointer to the client socket) telling the client proxy where to send the update message. When the client proxy **740** receives the update message, it uses the identifier to send the message to the client or clients.

In one embodiment, the node registry **732** contains, for each object ID, a list with one entry for each client registered for that ID. Each entry contains a (name of client proxy, name of client) pair where each "name" is a socket identifier or some other information for identifying how to route the message to the named entity. In another embodiment, the registry **732** is adapted to more efficiently handle the case where multiple clients at a given client proxy are registered for the same object ID by placing a list of clients registered for the object ID in each entry. Thus, each entry in the list for a given object ID contains a (name of client proxy, list of names of clients on that proxy) pair. When a message having a particular object ID arrives at the node, the node walks down the corresponding entries in the registry for that ID, pushing one copy of the message to each listed client proxy. Along with the message, the node includes the list of clients on that proxy (e.g., a list of pointers to sockets for the clients) who have registered for the object ID. When the client proxy receives the message and the list, it simply pushes a copy of the message to each client.

The approach using the pass-through client proxy **740** described above is advantageous because it reduces memory loads because the client proxy stores only minimal state, and each node has a restricted list of object IDs for which it is responsible. It also expedites message routing because the client proxy does very little processing to forward a message. Plus, this approach separates different potential stress points in the network into different components, allowing the network to be tuned to provide good performance. In particular, the nodes handle a large amount of data and store a large amount of state, but do not have to hold open a large number of connections. The client proxies do not store a large amount of data or state, but hold open a large number of connections to the clients. In different embodiments, the ratios of nodes to client proxies can be varied to match the network to its requirements.

In the second variation of client proxies, the client proxy **740** stores client registration information. FIG. **11** illustrates a high-level diagram of the routing network **110** in which the client proxy **740** stores the client registration information. In this embodiment, the node registry **732** stores the object IDs registered by client proxies **740** connected to the node. Each client proxy **740** preferably maintains a client proxy registry

736 containing the object IDs registered by clients 712 connected to the client proxy 740. The client proxy 740 registers with the nodes for all objects for which it needs to receive updates. Thus, this variation is similar to the hierarchical registry embodiment described above.

An alternative to maintaining a separate client proxy as described above is to have the nodes themselves provide the client proxy functionality. In this alternative, clients connect to a node of a first type. When a client registers for messages of a category not handled by the node, the node passes the registration to a node of the appropriate type along with an identifier of the client connection. In this case, the node connected to the client adopts the functionality of the pass-through client proxy. The node can also adopt the functionality of the client proxy that stores client registration information, thereby creating a hierarchy of registrations among the nodes themselves.

FIG. 12 is a high-level block diagram of the routing network 110 of the embodiment in which the nodes themselves adopt client proxy functionality. There are at least two ways to implement this embodiment. In one implementation there is a preferred node type, such as node of type 3, that always serves as a client proxy. This node type may be designated to handle the most common categories of messages. As a result, the client proxy functionality would be utilized very rarely. As shown in FIG. 12, client 1 712 and client 2 712 are connected to the node of type 3 726, which in turn is connected to a node of type 1 726 and a node of type 2 726.

In another implementation, instead of having one preferred node type that always serves as a client proxy, each node type can serve as a node and as a client proxy. For example, when a client 712 connects to a node 726, it passes the node all of the object IDs the client wishes to register. For object IDs of message categories handled by that node, the node stores the registrations in its registry. For object IDs of messages in other categories, the node acts as a client proxy and passes the registration requests to nodes of the appropriate types. This implementation can be made more efficient by providing functionality in the activation module for identifying the most common categories of messages sought by the client, and causing the client to connect to a node of a type that receives those categories of messages.

An alternative to connecting clients 712 to nodes 726 using a client proxy 740 is having a client 712 maintain multiple connections to the routing network 110. This can be accomplished, for example, by attaching message-categorization functionality to the activation module 124. The activation module 124 preferably determines which categories of messages it needs and to which nodes types it should connect. The activation module 124 then makes a connection to one node of each relevant type and registers the appropriate object IDs with the appropriate nodes.

An alternative to using client proxies is to allow clients to register for only categories of messages handled by one type of node. For example, if each node type handles messages from only one input source, each client 712 can be required to register for messages from only one input source. This can be done, for example, by configuring the load balancer 718 to ensure that clients registering for messages from one input source connect to a node of the type that receives messages from that input source.

FIG. 10 is a flow chart illustrating steps performed by a gateway 724, a node 726 that stores client registration information, and a pass-through client proxy 740 to perform object-based routing of a message received from an input source 710. Initially, the gateway 724 receives 1010 the message from an input source 710. The gateway 724 determines 1020 to which category the message belongs using one or more of the techniques described above. After the gateway 724 has determined to which category the message belongs, the gateway 724 identifies 1030 node types to which the message category is mapped and routes 1040 the message to the nodes of those types. Each node 726 that receives the message extracts 1050 the object ID from the message. Each node 726 examines its registry 732 to determine which clients have registered for the message, the client proxies to which the clients are connected, and the connections between the client proxies and the clients. Each node 726 then forwards 1060 the message to the identified client proxies 740 along with the identifiers for the client connections (e.g., the client socket or a list of client sockets). When a client proxy 740 receives the message from a node 726, it pushes 1070 the message to the identified client socket or sockets.

The above description is included to illustrate the operation of the preferred embodiments and is not meant to limit the scope of the invention. The scope of the invention is to be limited only by the following claims. From the above discussion, many variations will be apparent to one skilled in the relevant art that would yet be encompassed by the spirit and scope of the invention. For example, the invention may be implemented using a tangible computer-readable medium having stored thereon computer-executable instructions.

The invention claimed is:

1. A method comprising:

receiving, using a processing device, an update message from an input source, the update message identifying a live object and containing data for updating a property of the live object;

identifying a category of the update message based on the input source;

determining a node having a node type to which the update message is to be routed based on a mapping of categories of update messages to node types, the mapping controlling an amount of update message traffic through nodes of a routing network;

routing, using the processing device, the update message to the node having the determined node type;

causing the node, through the update message, to determine a client, different from the input source, that has registered for updates of the live object;

causing the node to route the update message from the node to the client; and

causing the client to process the update message and to update the property of the lire object.

2. The method of claim 1, further comprising:

causing the node to extract an object identifier (ID) from the update message, to establish a connection with the client, and to determine if the client has registered for updates of the live object based on the object ID.

3. The method of claim 2, further comprising:

causing the node to determine at least one client proxy with which the client communicates;

causing the node to route the update message to the client proxy; and

causing the client proxy to route the update message to the client to determine the connection.

4. The method of claim 3, further comprising causing the node to maintain client registration information concerning the client connection.

5. The method of claim 3, further comprising causing the node to maintain client registration information concerning the client connection at the client proxy.

6. The method of claim 1, wherein the property of the live object has a direct effect on a visual representation of the live

object in a data representation, has an effect on an internal aspect of the live object and has no effect on the visual representation of the live object in the data representation, or has a direct effect on one aspect of the visual representation of the live object in the data representation and has no effect on other aspect of the visual representation of the live object in the data representation.

7. A routing network comprising:

a gateway device configured to:

receive an update message from an input source, the update message identifying a live object and containing data for updating a property of the live object,

identify a category of the update message based on the input source;

determine a node having a node type to which the update message is to be routed based on a mapping of a categories of update messages to node types, the mapping controlling an amount of update message traffic through nodes of a routing network; and route the update message; and

the node configured to:

receive the update message from the gateway device, wherein the node is configured to be mapped to the node type,

determine a client, different from the input source, that has registered for updates of the live object, and

route the update message from the node to the client, wherein the client is adapted to process the update message and to update the property of the live object.

8. The routing network of claim **7**, wherein the node is configured to extract an object identifier (ID) from the update message and to determine a connection to the client to route the data to the client.

9. The routing network of claim **8**, wherein the node is configured to determine the connection by determining at least one client proxy with which the client communicates and to route the update message to the client proxy.

10. The routing network of claim **9**, wherein client registration information concerning the client connection is configured to be maintained at the node.

11. The routing network of claim **9**, wherein client registration information concerning the client connection is configured to be maintained at the client proxy.

12. The routing network of claim **7**, wherein:

the gateway device is configured to route the update message to the node adapted to receive messages of more than one message category, and

the node is configured to determine a connection with at least one client proxy with which the client communicates and to route the data to the client proxy.

13. The routing network of claim **7**, wherein the node type is configured to identify the node that receives the update message from the gateway device.

14. A method comprising:

providing, using a processing device of an input source, a data representation to a client device, different from the input source, coupled to a routing network, wherein the data representation includes at least one live object recognizable by the client device, and causing the client device to respond to the live object of the data representation by determining an object identifier (ID) of the live object and to register for updates of the live object with the routing network, such that registering the client device with the routing network provides client connection information to a node in the routing network; and

sending, using the processing device of the input source, an update message to the routing network, wherein the update message identifies the live object and contains update data that updates a property of the live object,

wherein a gateway device at the routing network is configured to identify a category of the update message based on the input source, to determine a node type to which the identified category maps, and to route the update message to the node, having the node type, at the routing network,

wherein the node is configured to identify the client device as a registered device and to route the update message to the client device, and

wherein the client device processes the update message upon receipt to update the property of the live object at the client device.

15. The method of claim **14**, wherein providing the data representation to the client device includes providing the live object that causes the client device to register with a client proxy of the routing network.

16. The method of claim **14**, wherein providing the data representation to the client device includes providing the live object that causes the client device to register with the node of the routing network.

17. The method of claim **14**, wherein providing the data representation to the client device includes providing an activation module that is executed by the client device and that registers the live object with the routing network.

18. The method of claim **17**, wherein providing the activation module includes providing the activation module that is configured to determine a node type that handles registration and that causes the client device to register with the node corresponding to the node type.

19. The method of claim **17**, wherein providing the activation module includes providing the activation module that is configured to determine a message category of the data representation and that causes the client device to register with a node having a node type corresponding to the message category.

20. An apparatus comprising:

an input source device configured to provide a data representation to a client device, different from the input source, coupled to a routing network, wherein the data representation includes at least one live object that is recognizable by the client device, and that causes the client device to determine an object identifier (ID) of the live object to register for updates of the bye object with the routing network, such that registering the client device with the routing network provides client connection information to the routing network,

wherein the input source device is configured to route an update message to the routing network, wherein the update message identifies the live object and contains update data for updating a property of the live object,

wherein a gateway device at the network is configured to identify a category of the update message based on the input source, to determine a node type to which the identified category maps, and to route the update message to a node of the node type at the routing network,

wherein the node is configured to identify the client device as a registered device and route the update message containing the update data to the client device, and

wherein the client device is configured to process the update message upon receipt to update the property of the bye object at the client device.

21. The apparatus of claim **20**, wherein the live object of the data representation is configured to cause the client device to register with a client proxy of the routing network.

22. The apparatus of claim **20**, wherein the live object of the data representation is configured to cause the client device to register with the node of the routing network.

23. The apparatus of claim **20**, wherein the received data representation includes an activation module that is configured to be executed by the client device and adapted to register the live object with the routing network.

24. The apparatus of claim **23**, wherein the activation module is configured to determine a node type for handling registration and to cause the client device to register with the node of the node type.

25. The apparatus of claim **23**, wherein the activation module determines a message category of the data representation and causes the client device to register with the node having a node type corresponding to the message category.

26. An article of manufacture including a computer-readable storage medium having instructions stored thereon, execution of which by a computing device causes the computing device to perform operations comprising:

providing, using a processing device of an input source, a data representation to a client device, different from the input source, coupled to a routing network, wherein the data representation includes at least one live object that is recognizable by the client device, and that causes the client device to respond to the live object by determining an object identifier (ID) of the live object to register for updates of the live object with the routing network, such that registering the client device with the routing network provides client connection information to the routing network; and

sending, using the processing device of the input source, an update message to the routing network, wherein the update message identifies the live object and contains update data for updating a property of the live object,

wherein a gateway device at the routing network is configured to identify a category of the update message based on the input source, to determine a node type to which the identified category maps, and to route the update message to a node of the node type at the routing network,

wherein the node is configured to identify the client device as a registered device and to route the update message to the client device, and

wherein the client device is configured to process the update message upon receipt to update the property of the live object at the client device.

27. The article of manufacture of claim **26**, wherein the live object of the data representation causes the client device to register with a client proxy of the routing network.

28. The article of manufacture of claim **26**, wherein the live object of the data representation causes the client device to register with the node of the routing network.

29. The article of manufacture of claim **26**, wherein the received data representation includes an activation module executed by the client device and adapted to register the live object with the routing network.

30. The article of manufacture of claim **29**, wherein the activation module determines a node type that handles registration and causes the client device to register with the node of the node type.

31. The article of manufacture of claim **29**, wherein the activation module determines a message category of the data representation and causes the client device to register with the node having a node type corresponding to the message category.

32. A non-transitory computer readable storage medium having instructions stored thereon, the instructions comprising:

instructions for providing a data representation to a client device coupled to a routing network, wherein the data representation includes at least one live object recognizable by the client device, and wherein the client device is configured to respond to the live object of the data representation by determining an object identifier (ID) of the live object to register for updates of the live object with the routing network, such that registering the client device with the routing, network provides client connection information to the routing network; and

instructions for providing, using a processing device of an input source, different from the client device, an update message to the routing network, wherein the update message identifies the live object and contains update data for updating a property of the live object,

wherein a gateway device at the routing network is configured to identify a category of the update message based on the input source, to determine a node type to which the identified category maps, and to route the update message to a node of the node type at the routing network,

wherein the node is configured to identify the client device as a registered device and to route the update message containing the update data to the client device, and

wherein the client device is configured to process the update message upon receipt to update the property of the live object at the client device.

33. The non-transitory computer readable storage medium of claim **32**, wherein the live object of the data representation causes the client device to register with a client proxy of the routing network.

34. The non-transitory computer readable storage medium of claim **32**, wherein the live object of the data representation causes the client device to register with the node of the routing network.

35. The non-transitory computer readable storage medium of claim **32**, wherein the data representation includes an activation module that is executed by the client device and that is adapted to register the live object with the routing network.

36. The non-transitory computer readable storage medium of claim **35**, wherein the activation module determines a node type for handling registration and causes the client device to register with the node of the node type.

37. The non-transitory computer readable storage medium of claim **35**, wherein the activation module determines a message category of the data representation and causes the client device to register with the node having the node type corresponding to the message category.

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UNITED STATES PATENT AND TRADEMARK OFFICE

CERTIFICATE OF CORRECTION

PATENT NO. : 8,407,722 B2
APPLICATION NO. : 11/396251
DATED : March 26, 2013
INVENTOR(S) : Tuttle et al.

It is certified that error appears in the above-identified patent and that said Letters Patent is hereby corrected as shown below:

On Title Page 2, Item (56), under "OTHER PUBLICATIONS", in Column 2, Line 3, delete "Retrived" and insert -- Retrieved --, therefor.

On Title Page 3, Item (56), under "OTHER PUBLICATIONS", in Column 1, Line 12, delete "Slavak Republic," and insert -- Slovak Republic, --, therefor.

On Title Page 3, Item (56), under "OTHER PUBLICATIONS", in Column 1, Line 48, delete "library/default.asp?url=/library/en-us/dndotnet/hmtl/progthrepool." and insert -- library/default.asp?url=/library/en-us/dndotnet/html/progthrepool. --, therefor.

In the Specifications

In Column 3, Line 24, delete "live: objects" and insert -- live objects --, therefor.

In Column 12, Line 7, delete "provider 210" and insert -- provider 108 --, therefor.

In Column 14, Line 57, delete "network 116" and insert -- network 110--, therefor.

In the Claims

In Column 22, Line 47, in Claim 1, delete "of the lire" and insert -- of the live --, therefor.

In Column 24, Line 45, in Claim 20, delete "of the bye" and insert -- of the live --, therefor.

In Column 24, Line 64, in Claim 20, delete "bye object" and insert -- live object --, therefor.

In Column 26, Line 17, in Claim 32, delete "routing, network" and insert -- routing network --, therefor.

Signed and Sealed this
Thirteenth Day of May, 2014

Michelle K. Lee

Michelle K. Lee
Deputy Director of the United States Patent and Trademark Office

US008407722C1

(12) EX PARTE REEXAMINATION CERTIFICATE (12886th)

United States Patent

Tuttle et al.

(10) Number: US 8,407,722 C1

(45) Certificate Issued: Apr. 4, 2025

(54) **ASYNCHRONOUS MESSAGING USING A NODE SPECIALIZATION ARCHITECTURE IN THE DYNAMIC ROUTING NETWORK**

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(73) Assignee: **INTELLECTUAL VENTURES I LLC**

Reexamination Request:
No. 90/019,662, Sep. 16, 2024

Reexamination Certificate for:

Patent No.:	**8,407,722**
Issued:	**Mar. 26, 2013**
Appl. No.:	**11/396,251**
Filed:	**Mar. 30, 2006**

Certificate of Correction issued May 13, 2014

Related U.S. Application Data

(63) Continuation of application No. 10/105,018, filed on Mar. 21, 2002, now Pat. No. 7,051,070, which is a continuation-in-part of application No. 10/017,182, filed on Dec. 14, 2001, now Pat. No. 7,043,525.

(60) Provisional application No. 60/278,303, filed on Mar. 21, 2001, provisional application No. 60/280,627, filed on Mar. 29, 2001, provisional application No. 60/279,608, filed on Mar. 28, 2001, provisional application No. 60/276,847, filed on Mar. 16, 2001, provisional application No. 60/256,613, filed on Dec. 18, 2000.

(51) **Int. Cl.**

G06F 3/00	(2006.01)
G06F 15/16	(2006.01)
H04L 45/302	(2022.01)
H04L 67/55	(2022.01)
H04L 67/63	(2022.01)
H04L 67/1001	(2022.01)
H04L 67/1008	(2022.01)
H04L 67/1014	(2022.01)
H04L 67/1029	(2022.01)

(52) **U.S. Cl.**
CPC ***H04L 67/55*** (2022.05); ***H04L 45/306*** (2013.01); ***H04L 67/63*** (2022.05); *H04L 67/1001* (2022.05); *H04L 67/1008* (2013.01); *H04L 67/1014* (2013.01); *H04L 67/1029* (2013.01)

(58) **Field of Classification Search**
None
See application file for complete search history.

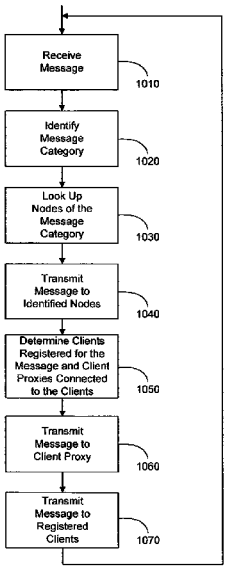
(56) **References Cited**

To view the complete listing of prior art documents cited during the proceeding for Reexamination Control Number 90/019,662, please refer to the USPTO's Patent Electronic System.

Primary Examiner — Woo H Choi

(57) **ABSTRACT**

A network routes update messages containing updates to properties of live objects from input sources to clients having the objects. When the clients receive live objects, the clients identify the object IDs associated with the objects and register the object IDs with the routing network. The routing network is adapted to selectively send update messages to nodes in the network and the nodes forward the messages to the clients. One implementation uses a hierarchy of registries to indicate which nodes and clients receive which update messages. Another implementation assigns update messages to one or more of N categories and nodes to one or more of M types, and the gateways maintain mapping between categories and types. To ensure that clients receive all of the update messages for which they register, the clients connect to client proxies that in turn connect to at least one node of each type.



EX PARTE REEXAMINATION CERTIFICATE

NO AMENDMENTS HAVE BEEN MADE TO THE PATENT

AS A RESULT OF REEXAMINATION, IT HAS BEEN DETERMINED THAT:

The patentability of claims **1**, **6-7**, **14-17**, **20-23**, **26-29** and **32-35** is confirmed.

Claims **2-5**, **8-13**, **18-19**, **24-25**, **30-31** and **36-37** were not reexamined.

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