Specification Volume 0

SPECIFICATION OF THE BLUETOOTH SYSTEM

Wireless Connections Made Easy



Master Table of Contents & Compliance Requirements

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TOYOTA EXHIBIT 1014

Revision History

The Revision History is shown in the “Appendix” on page 41[vol. 0].

Contributors

The persons who contributed to this specification are listed in the Appendix.

Web Site

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In addition each volume has a TOC and each part of a volume is preceded by a detailed TOC.

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- The TOC for each Volume starts at the top of a page.
- The Volume number in black is followed by the Volume name in red.

Note: Each Volume is a self contained book which is published and updated separately and is equipped with a TOC of its own. However, this Master TOC is also revised as soon as any of the other Volumes are updated.

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Red or blue text on the following pages indicates hypertext links that take you directly to the indicated section, on condition that you have access to a complete specification.

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BLUETOOTH COMPLIANCE REQUIREMENTS

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1 INTRODUCTION

The Bluetooth Qualification Program Reference Document (PRD) is the primary reference document for the Bluetooth Qualification Program and defines its requirements, functions, and policies. The PRD is available on the Bluetooth Web site.

Passing the Bluetooth Qualification Process demonstrates a certain measure of compliance and interoperability, but because products are not tested for every aspect of this Bluetooth Specification, qualification does not guarantee compliance. Passing the Bluetooth Qualification Process only satisfies one condition of the license grant. The Member has the ultimate responsibility to ensure that the qualified product complies with this Bluetooth Specification and interoperates with other products.

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2 SCOPE

This part of the specification defines some fundamental concepts used in the Bluetooth Qualification Program.

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3 DEFINITIONS

Bluetooth Qualification Process – The process defined in the Bluetooth Qualification Program Reference Document (PRD) to qualify a design used in implementations of Bluetooth wireless technology.

Bluetooth Qualification Program – The Bluetooth Qualification Process together with other related requirements and processes.

3.1 TYPES OF BLUETOOTH PRODUCTS

Bluetooth Product – Any product containing an implementation of Bluetooth wireless technology. All Bluetooth Products shall be one of the following:

- Bluetooth End Product
- Bluetooth Host Subsystem Product
- Bluetooth Controller Subsystem Product
- Bluetooth Profile Subsystem Product
- Bluetooth Component Product
- Bluetooth Development Tool
- Bluetooth Test Equipment

Bluetooth End Product - An implementation of Bluetooth wireless technology that implements, at a minimum, all mandatory requirements in Radio, Baseband, Link Manager, Logical Link Control and Adaptation Protocol, Service Discovery Protocol and Generic Access Profile parts of the Specification.

Bluetooth Subsystem Product - An implementation of Bluetooth wireless technology that implements only a portion of the Specification, in compliance with such portion of the Specification, and in accordance with the mandatory requirements as defined herein. Bluetooth Subsystem Products can be qualified solely for distribution and the use of Bluetooth wireless technology in Bluetooth Subsystem Products require such Bluetooth Subsystem Products to be combined with a complementary Bluetooth End Product or one or more complementary Bluetooth Subsystem Products such that the resulting combination satisfies the requirements of a Bluetooth End Product. There are three types of Bluetooth Subsystem Products as defined below:

- *Bluetooth Host Subsystem Product* – A Bluetooth Subsystem Product containing, at a minimum, all the mandatory requirements defined in the Host Controller Interface, Logical Link Control and Adaptation Protocol, Service Discovery Protocol and Generic Access Profile parts of this Specification, but none of the protocols below Host Controller Interface (HCI). In addition, a Bluetooth Host Subsystem Product may contain, at a minimum, all the mandatory requirements defined in one or more of the protocols and profiles above HCI.

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- *Bluetooth Controller Subsystem Product* – A Bluetooth Subsystem Product containing, at a minimum, all the mandatory requirements defined in the Bluetooth Radio, Baseband, Link Manager and HCI parts of this Specification, but none of the Protocols and Profiles above HCI.
- *Bluetooth Profile Subsystem Product* – A Bluetooth Subsystem Product containing, at a minimum, all the mandatory requirements defined in one or more of the profile specifications.

Bluetooth Component Product - An implementation of Bluetooth wireless technology, which does not meet the requirements of a Bluetooth End Product, but implements, at a minimum, all the mandatory requirements of either one or more of any of the protocol and profile parts of the Specification in compliance with such portion of the Specification. Bluetooth Component Products can be qualified solely for distribution and the use of the Bluetooth wireless technology in Bluetooth Component Products require such Bluetooth Component Products to be incorporated in Bluetooth End Products or Bluetooth Subsystem Products.

Bluetooth Development Tool - An implementation of Bluetooth wireless technology, intended to facilitate the development of new Bluetooth designs. Bluetooth Development Tools can be qualified solely for distribution and the use of the Bluetooth wireless technology in development of new Bluetooth Products.

Bluetooth Test Equipment - An implementation of Bluetooth wireless technology, intended to facilitate the testing of new Bluetooth Products. Bluetooth Test Equipment can be qualified solely for distribution and the use of the Bluetooth wireless technology in testing of new Bluetooth Products. Where necessary, Bluetooth Test Equipment may deviate from the Specification in order to fulfill the test purposes in the Bluetooth Test Specifications.

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4 CORE CONFIGURATIONS

This section defines the set of features that are required for a product to be qualified to a specification name. Each core configuration is defined by a set of LMP feature bits or L2CAP feature bits that shall be supported to allow the configuration name to be used.

The configuration requirements imposed on a device depends on the profiles that it supports.

4.1 SPECIFICATION NAMING CONVENTIONS

Each specification is named by its core specification version number, followed by a list of the core configuration names that are implemented and qualified.

A complete specification name shall be stated as the core specification version number followed by “+”, and then either a single core configuration name or a sequence of core configuration names separated by “+”.

Examples of complete specification names including the core configuration names:

- Bluetooth v2.1
- Bluetooth v2.1 + EDR

In this example, a product claiming “Bluetooth v2.1” may implement some of the EDR features, following the requirements in other parts of the specifications, and be qualified for those features. If the full set required in Section 4.2 are not supported the “+ EDR” configuration name shall not be used in product literature.

4.2 EDR CONFIGURATIONS

This section specifies additional compliance requirements that shall be followed if the configuration name “EDR” is used within the complete specification name. The configuration name “EDR” may only be used with core specification version number 2.0 or later versions of the specification.

Table 4.1 defines three categories of Transport Requirements that shall be satisfied subject to the following rules:

- A Bluetooth product shall support category 1 whenever it supports asynchronous transports for the profiles it incorporates.
- A Bluetooth product shall support category 2 whenever it supports asynchronous transports with multislot ACL packets for the profiles it incorporates.
- A Bluetooth product shall support category 3 whenever it supports eSCO synchronous transports for the profiles it incorporates.

.

No.	Transport Requirements	LMP Feature Bits Required	L2CAP Feature Bits Required
1	EDR for asynchronous transports (single slot)	Enhanced Data Rate ACL 2 Mbps mode (25) Enhanced Data Rate ACL 3 Mbps mode (26)	None
2	EDR for asynchronous transports (multi-slot)	3-slot Enhanced Data Rate ACL packets (39) 5-slot Enhanced Data Rate ACL packets (40)	None
3	EDR for synchronous transports	Enhanced Data Rate eSCO 2 Mbps mode (45) Enhanced Data Rate eSCO 3 Mbps mode (46)	None

Table 4.1: EDR configuration requirements

Note: No additional requirements are stated on the support of 2-EV5 and 3-EV5 packets.

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1 REVISION HISTORY

Public versions are marked with bold in the tables below.

Beginning with v1.2 of the Core System Package the core Bluetooth specification documents, protocols and profiles, were transferred to a new partitioning comprising volumes and individual profile specifications are each contained in an individual document instead of the two volumes (Core and Profiles) used in v1.1.

For more detailed information about changes between versions published before v1.2, please see the Appendix I “Revision History” in v1.1.

1.1 [VOL 0] MASTER TOC & COMPLIANCE REQUIREMENTS

1.1.1 Bluetooth Compliance Requirements

Rev	Date	Comments
v2.1 + EDR		No content changes. Updates to the Table of Contents.
v2.0 + EDR	Oct 15 2004	This version of the specification is intended to be a separate Bluetooth Specification that has all the functional characteristics of the v1.2 Bluetooth Specification that adds the Enhanced Data Rate (EDR) feature which required changes to Volume 0, Part A, Master Table of Contents.
v1.2	Nov 05 2003	This Part was moved from the Core volume. No content changes been made to this document since v1.1

1.2 [VOL 1] ARCHITECTURE & TERMINOLOGY OVERVIEW

Rev	Date	Comments
v2.1 + EDR		Added definitions for new features: Encryption Pause Resume, Erroneous Data reporting, Extended Inquiry Response, Link Supervision Timeout Event, Packet Boundary Flag, Secure Simple Pairing, Sniff Subrating
v2.0 + EDR	Oct 15 2004	This version of the specification is intended to be a separate Bluetooth Specification that has all the functional characteristics of the v1.2 Bluetooth Specification that adds the Enhanced Data Rate (EDR) feature which incorporates changes to Volume 1, Part B, Acronyms and Abbreviations.
v1.2	Nov 05 2003	New volume with informational content. This volume will always be updated in parallel with the Core System volumes

1.3 [VOL 2 & 3] CORE SYSTEM PACKAGE

Rev	Date	Comments
v2.1 + EDR		• New features added in 2.1 + EDR: -Encryption Pause and Resume -Erroneous Data Reporting -Extended Inquiry Response -Link Supervision Timeout Changed Event -Non-Flushable Packet Boundary Flag -Secure Simple Pairing -Sniff Subrating -Security Mode 4 • Updates to IEEE language in Volume 2, Part H, Security • Errata for v2.0 + EDR
v2.0 + EDR	Aug 01 2004	This version of the specification is intended to be a separate Bluetooth Specification. This specification was created by adding EDR and the errata.
v1.2	Nov 05 2003	New features added in v1.2: - Architectural overview - Faster connection - Adaptive frequency hopping - Extended SCO links - Enhanced error detection and flow control - Enhanced synchronization capability - Enhanced flow specification The Core System Package now comprises two volumes and the text has gone through a radical change both in terms of structure and nomenclature. The language is also more precise and is adapted to meet the IEEE standard. The following parts are moved from the Core System Package to other volumes or were deprecated: RFCOMM [vol 7], Object Exchange (IrDA Interoperability) [vol 8], TCS [vol 9], Interoperability Requirements for Bluetooth as a WAP Bearer [vol 6], HCI USB Transport Layer [vol4], HCI RS232 Transport Layer [vol 4], HCI UART Transport Layer [vol 4], Bluetooth Compliance Requirements [vol 0], Optional Paging Schemes [deprecated]
1.1	Feb 22 2001	The specification was updated with Errata items previously published on the web site. The Bluetooth Assigned Numbers appendix was lifted out from the specification to allow continuous maintenance on the web site.

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Rev	Date	Comments
1.0B	Dec. 1 1999	The specification was updated with Errata items previously published on the web site and was revised from a linguistic point of view. The following parts were added: Interoperability Requirements for Bluetooth as a WAP Bearer, Test Control Interface, Sample Data (appendix), Bluetooth Audio (appendix), Baseband Timers (appendix) and Optional Paging Scheme (appendix)
1.0a	July 26 1999	The first version of the Bluetooth Specification published on the public web site. Added part: Bluetooth Compliance Requirements.
1.0 draft	July 5th 1999	The following parts were added: Service Discovery Protocol (SDP), Telephony Control Specification (TCS), Bluetooth Assigned Numbers (appendix) and Message Sequence Charts (appendix)
0.9	April l30 1999	The following parts were added: IrDA Interoperability, HCI RS232 Transport Layer, HCI UART Transport Layer and Test Mode
0.8	Jan 21 1999	The following parts were added: Radio Specification, L2CAP, RFCOMM, HCI & HCI USB Transport Layer
0.7	Oct 19 1998	This first version only included Baseband and Link Manager Protocol

1.4 [VOL 4] TRANSPORT LAYERS

Rev	Date	Comments
v2.1 + EDR	Feb 22 2007	Added this volume to the specification

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Yoshimitsu Shimojo	Toshiba
Toshiki Kizu	Toshiba
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Sam Turner	TTPCom

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Toru Aihara	IBM Corporation
Chatschik Bisdikian	IBM Corporation
Kris Fleming	Intel Corporation
David E. Cypher	NIST
Thomas Busse	Nokia Corporation
Julien Corthial	Nokia Corporation
Olaf Joeressen	Nokia Corporation
Thomas Müller	Nokia Corporation
Dong Nguyen	Nokia Corporation
Harmke de Groot	Philips
Terry Bourk	Silicon Wave
Johannes Elg	Telefonaktiebolaget LM Ericsson
Jaap Haartsen	Telefonaktiebolaget LM Ericsson
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Mary A. DuVal	Texas Instruments
Onn Haran	Texas Instruments
John Mersh	TTPCom

2.3.4 Part D: Error Codes

Version 2.1 + EDR

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Shawn Ding	Broadcom
Robin Heydon	CSR
Simon Kingston	CSR
Steven Singer	CSR
Steven Wenham	CSR
Selim Aissi	Intel
Penny Chen	Intel
Mattias Edlund	Infineon
Josh Benaloh	Microsoft
Andy Glass	Microsoft
Peter Hauser	Microsoft
Joby Lafky	Microsoft
Kristin Lauter	Microsoft
Dan Simon	Microsoft
Don Stanwyck	Microsoft
Yacov Yacobi	Microsoft
Gideon Yuval	Microsoft
Greg Muchnik	Motorola
N Asokan	Nokia
Philip Ginzboorg	Nokia
Kanji Kerai	Nokia
Noel Lobo	Nokia
Kaisa Nyberg	Nokia
Arto Palin	Nokia
Päivi Ruuska	Nokia
Dominique Everaere	NXP
Reinhard Meindl	NXP
Joel Linsky	Qualcomm
Terry Bourk	Qualcomm
Guido Bertoni	ST
Eran Reuveni	Texas Instruments
Shimojo Yoshimitsu	Toshiba

Version 1.2

Robin Heydon (section owner)	CSR
Roland Hellfajer	Infineon
Joel Linsky	Silicon Wave
John Mersh	TTPCom

This part was earlier included in the LMP and HCI functional Specifications.

2.3.5 Part E: Bluetooth Host Controller Interface Functional Specification

Version 2.1 + EDR

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Shawn Ding	Broadcom
Henrik Hedlund	CSR
Robin Heydon	CSR
Simon Kingston	CSR
Steven Singer	CSR
Steven Wenham	CSR
Paul Wright	CSR
Mattias Edlund	Infineon
David Suvak	iAnywhere
Selim Aissi	Intel
Penny Chen	Intel
Josh Benaloh	Microsoft
Andy Glass	Microsoft
Peter Hauser	Microsoft
Joby Lafky	Microsoft
Kristin Lauter	Microsoft
Dan Simon	Microsoft
Don Stanwyck	Microsoft
Yacov Yacobi	Microsoft
Gideon Yuval	Microsoft
Greg Muchnik	Motorola
N Asokan	Nokia
Philip Ginzboorg	Nokia
Kanji Kerai	Nokia
Noel Lobo	Nokia
Kaisa Nyberg	Nokia
Arto Palin	Nokia
Päivi Ruuska	Nokia
Dominique Everaere	NXP
Javier del Prado Pavon	NXP
Reinhard Meindl	NXP
Joel Linsky	Qualcomm
Terry Bourk	Qualcomm
Tim Howes	Symbian
Jorgen van Parijs	ST
Guido Bertoni	ST
Amihai Kidron	Texas Instruments
Eran Reuveni	Texas Instruments
Shimojo Yoshimitsu	Toshiba

Version 2.0 + EDR

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Joel Linsky	RF Micro Devices

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Robin Heydon	CSR

Version 1.2

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Henrik Andersen	Digianswer/Motorola
Jan Åberg	Ericsson
Martin van der Zee	Ericsson
Don Liechty	Extended Systems
Kevin Marquess	Hyper Corp
Roland Hellfajer	Infineon
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Tod Sizer	Lucent Technologies
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Andy Glass	Microsoft
Adrian Stephens	Mobilian
Jürgen Schnitzler	Nokia
Thomas Müller	Nokia
Rene Tischer	Nokia
Rob Davies	Philips
Antonio Salloum	Philips
Joel Linsky	Silicon Wave
Terry Bourk	Silicon Wave
Len Ott	Socket Communications
Randy Erman	Taiyo Yuden
Yoshimitsu Shimojo	Toshiba
Toshiki Kizu	Toshiba
Katsuhiro Kinoshita	Toshiba
Sam Turner	TTPCom
John Mersh	TTPCom

Previous versions

Todor Cooklev	3Com Corporation
Toru Aihara	IBM Corporation
Chatschik Bisdikian	IBM Corporation
Nathan Lee	IBM Corporation
Akihiko Mizutani	IBM Corporation
Les Cline	Intel Corporation
Bailey Cross	Intel Corporation
Kris Fleming	Intel Corporation
Robert Hunter	Intel Corporation
Jon Inouye	Intel Corporation
Srikanth Kambhatla	Intel Corporation

Steve Lo	Intel Corporation
Vijay Suthar	Intel Corporation
Bruce P. Kraemer	Intersil
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David E. Cypher	NIST
Thomas Busse	Nokia Corporation
Julien Courthial	Nokia Corporation
Thomas Müller	Nokia Corporation
Dong Nguyen	Nokia Corporation
Jürgen Schnitzler	Nokia Corporation
Fujio Watanabe	Nokia Corporation
Christian Zechlin	Nokia Corporation
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Patrik Lundin	Telefonaktiebolaget LM Ericsson
Tobias Melin	Telefonaktiebolaget LM Ericsson
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Thomas M. Siep	Texas Instruments
Masahiro Tada	Toshiba Corporation
John Mersh	TTPCom

2.3.6 Part F: Message Sequence Charts

Version 2.1 + EDR

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Robert Hulvey	Broadcom
Shawn Ding	Broadcom
Henrik Hedlund	CSR
Robin Heydon	CSR
Simon Kingston	CSR
Steven Singer	CSR
Steven Wenham	CSR
Selim Aissi	Intel
Penny Chen	Intel
Mattias Edlund	Infineon
Josh Benaloh	Microsoft
Andy Glass	Microsoft
Peter Hauser	Microsoft
Joby Lafky	Microsoft
Kristin Lauter	Microsoft
Dan Simon	Microsoft
Don Stanwyck	Microsoft
Yacov Yacobi	Microsoft
Gideon Yuval	Microsoft
Greg Muchnik	Motorola

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N Asokan	Nokia
Philip Ginzboorg	Nokia
Kanji Kerai	Nokia
Noel Lobo	Nokia
Kaisa Nyberg	Nokia
Päivi Ruuska	Nokia
Dominique Everaere	NXP
Reinhard Meindl	NXP
Joel Linsky	Qualcomm
Terry Bourk	Qualcomm
Guido Bertoni	ST
Tim Howes	Symbian
Amihai Kidron	Texas Instruments
Eran Reuveni	Texas Instruments
Shimojo Yoshimitsu	Toshiba

Version 1.2

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Simon Morris	CSR
Jan Åberg	Ericsson
Christian Gehrmann	Ericsson
Joel Linsky	Silicon Wave
John Mersh	TTPCom

Previous versions

Todor Cooklev	3Com Corporation
Toru Aihara	IBM Corporation
Chatschik Bisdikian	IBM Corporation
Nathan Lee	IBM Corporation
Kris Fleming	Intel Corporation
Greg Muchnik	Motorola, Inc.
David E. Cypher	NIST
Thomas Busse	Nokia Corporation
Dong Nguyen (Section Owner)	Nokia Corporation
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Tobias Melin	Telefonaktiebolaget LM Ericsson
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2.3.7 Part G: Sample Data

Version 2.1 + EDR

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Robin Heydon	CSR
Simon Kingston	CSR
Steven Singer	CSR
Steven Wenham	CSR
Selim Aissi	Intel
Penny Chen	Intel
Josh Benaloh	Microsoft
Andy Glass	Microsoft
Peter Hauser	Microsoft
Joby Lafky	Microsoft
Kristin Lauter	Microsoft
Dan Simon	Microsoft
Don Stanwyck	Microsoft
Yacov Yacobi	Microsoft
Gideon Yuval	Microsoft
Greg Muchnik	Microsoft
N Asokan	Nokia
Philip Ginzboorg	Nokia
Kanji Kerai	Nokia
Noel Lobo	Nokia
Kaisa Nyberg	Nokia
Arto Palin	Nokia
Päivi Ruuska	Nokia
Dominique Everaere	NXP
Reinhard Meindl	NXP
Guido Bertoni	ST
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2.3.8 Part H: Security Specification

Version 2.1 + EDR

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Henrik Hedlund	CSR

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Robin Heydon	CSR
Simon Kingston	CSR
Steven Singer	CSR
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Selim Aissi	Intel
Penny Chen	Intel
Josh Benaloh	Microsoft
Andy Glass	Microsoft
Peter Hauser	Microsoft
Joby Lafky	Microsoft
Kristin Lauter	Microsoft
Dan Simon	Microsoft
Don Stanwyck	Microsoft
Yacov Yacobi	Microsoft
Gideon Yuval	Microsoft
Greg Muchnik	Microsoft
Mattias Edlund	Infineon
N Asokan	Nokia
Philip Ginzboorg	Nokia
Kanji Kerai	Nokia
Noel Lobo	Nokia
Kaisa Nyberg	Nokia
Arto Palin	Nokia
Päivi Ruuska	Nokia
Dominique Everaere	NXP
Reinhard Meindl	NXP
Joel Linsky	QUALCOMM
Terry Bourk	QUALCOMM
Guido Bertoni	ST
Amihai Kidron	Texas Instruments
Eran Reuveni	Texas Instruments
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2.4 [VOL 3] CORE SYSTEM PACKAGE, HOST

2.4.1 Part A: Logical Link Control and Adaptation Protocol Specification

Version 2.1 + EDR

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Robin Heydon	CSR
Simon Kingston	CSR
Steven Singer	CSR
Steven Wenham	CSR
Selim Aissi	Intel
Penny Chen	Intel
Dave Suvak	iAnywhere
Josh Benaloh	Microsoft
Andy Glass	Microsoft
Peter Hauser	Microsoft
Joby Lafky	Microsoft
Kristin Lauter	Microsoft
Dan Simon	Microsoft
Don Stanwyck	Microsoft
Yacov Yacobi	Microsoft
Gideon Yuval	Microsoft
Greg Muchnik	Microsoft
N Asokan	Nokia
Philip Ginzboorg	Nokia
Kanji Kerai	Nokia
Noel Lobo	Nokia
Kaisa Nyberg	Nokia
Dominique Everaere	NXP
Javier del Prado Pavon	NXP
Reinhard Meindl	NXP
Joel Linsky	QUALCOMM
Terry Bourk	QUALCOMM
Tim Howes	Symbian
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Jennifer Bray	CSR
Jan Åberg	Ericsson

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Doron Holan	Microsoft
Andy Glass	Microsoft
Brian Redding	Motorola
Jürgen Schnitzler	Nokia
Thomas Müller	Nokia
Rob Davies	Philips
Terry Bourk	Silicon Wave
Michael Hasling	Tality

Previous versions

Jon Burgess	3Com Corporation
Paul Moran	3Com Corporation
Doug Kogan	Extended Systems
Kevin D. Marquess	Hyper Corporation
Toru Aihara	IBM Corporation
Chatschik Bisdikian	IBM Corporation
Kris Fleming	Intel Corporation
Uma Gadamsetty	Intel Corporation
Robert Hunter	Intel Corporation
Jon Inouye	Intel Corporation
Steve C. Lo	Intel Corporation
Chunrong Zhu	Intel Corporation
Sergey Solyanik	Microsoft Corporation
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Nada Golmie	NIST
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Rauno Makinen	Nokia Corporation
Thomas Müller	Nokia Corporation
Petri Nykänen	Nokia Corporation
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Kinoshita Katsuhiro	Toshiba Corporation

2.4.2 Part B: Service Discovery Protocol (SDP)

Version 2.1 + EDR

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Parviz Kermani	IBM Corporation
Brent Miller	IBM Corporation
Dick Osterman	IBM Corporation
Bob Pascoe	IBM Corporation
Jon Inouye	Intel Corporation
Srikanth Kambhatla	Intel Corporation
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Kati Rantala	Nokia Corporation
Thomas Müller	Nokia Corporation
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Kazuaki Iwamura	Toshiba Corporation

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2.4.3 Part C Generic Access Profile

Version 2.1 + EDR

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Mattias Edlund	Infineon
Päivi Ruuska	Nokia
Arto Palin	Nokia
Joel Linsky	QUALCOMM
Terry Bourk	QUALCOMM
Shimojo Yoshimitsu	Toshiba

Version 1.2

Jennifer Bray	CSR
Alexander Thoukydides	CSR
Christian Gehrmann	Ericsson
Henrik Hedlund	Ericsson
Jan Åberg	Ericsson
Stefan Agnani	Ericsson
Thomas Müller	Nokia
Joel Linsky	Silicon Wave
Terry Bourk	Silicon Wave
Katsuhiro Kinoshita	Toshiba

Previous versions

Ken Morley	3Com Corporation
Chatschik Bisdikian	IBM Corporation
Jon Inouye	Intel Corporation
Brian Redding	Motorola, Inc.
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Stephane Bouet	Nokia Corporation
Thomas Müller	Nokia Corporation
Martin Roter	Nokia Corporation
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Patric Lind (Section Owner)	Telefonaktiebolaget LM Ericsson
Erik Slotboom	Telefonaktiebolaget LM Ericsson
Johan Sörensen	Telefonaktiebolaget LM Ericsson

2.4.4 Part D: Test Support

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Stefan Agnani (section owner)	Ericsson
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Joel Linsky	Silicon Wave
Michael Hasling	Tality

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Arno Kefenbaum	Nokia Corporation
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Roland Schmale	Nokia Corporation
Fujio Watanabe	Nokia Corporation
Stefan Agnani	Telefonaktiebolaget LM Ericsson
Mårten Mattsson	Telefonaktiebolaget LM Ericsson
Tobias Melin	Telefonaktiebolaget LM Ericsson
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Mårten Mattsson	Telefonaktiebolaget LM Ericsson
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2.5 [VOL 4] HOST CONTROLLER INTERFACE [TRANSPORT LAYER]

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Robert Hunter	Intel Corporation
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Uwe Gondrum	Nokia Corporation
Thomas Müller	Nokia Corporation
Christian Zechlin	Nokia Corporation
Johannes Elg	Telefonaktiebolaget LM Ericsson
Sven Jerlhagen	Telefonaktiebolaget LM Ericsson
Christian Johansson	Telefonaktiebolaget LM Ericsson
Patrik Lundin	Telefonaktiebolaget LM Ericsson
Lars Novak	Telefonaktiebolaget LM Ericsson
Masahiro Tada	Toshiba Corporation
Steve Ross	Digianswer A/S
Chatschik Bisdikian	IBM Corporation
Les Cline	Intel Corporation
Brad Hosler	Intel Corporation
John Howard	Intel Corporation
Srikanth Kambhatla	Intel Corporation
Kosta Koeman	Intel Corporation
John McGrath	Intel Corporation
Patrik Lundin	Telefonaktiebolaget LM Ericsson
Leonard Ott	Socket Communications
Rebecca O'Dell	Signia Technologies
Tsuyoshi Okada	Matsushita Electric
Robin Heydon	CSR
Toru Aihara	IBM Corporation
Edgar Kerstan	IBM Corporation
Nathan Lee	IBM Corporation
Kris Fleming	Intel Corporation
Robert Hunter	Intel Corporation
Patrick Kane	Motorola, Inc.
Uwe Gondrum	Nokia Corporation
Thomas Müller	Nokia Corporation
Christian Zechlin	Nokia Corporation
Johannes Elg	Telefonaktiebolaget LM Ericsson
Sven Jerlhagen	Telefonaktiebolaget LM Ericsson
Christian Johansson	Telefonaktiebolaget LM Ericsson
Patrik Lundin	Telefonaktiebolaget LM Ericsson
Lars Novak	Telefonaktiebolaget LM Ericsson
Masahiro Tada	Toshiba Corporation

Steve Ross	Digianswer A/S
Chatschik Bisdikian	IBM Corporation
Les Cline	Intel Corporation
Brad Hosler	Intel Corporation
John Howard	Intel Corporation
Srikanth Kambhatla	Intel Corporation
Kosta Koeman	Intel Corporation
John McGrath	Intel Corporation
Patrik Lundin	Telefonaktiebolaget LM Ericsson
Leonard Ott	Socket Communications
Rebecca O'Dell	Signia Technologies
Tsuyoshi Okada	Matsushita Electric
Robin Heydon	CSR

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Specification Volume 1

SPECIFICATION OF THE *BLUETOOTH* SYSTEM

Wireless Connections Made Easy



Architecture & Terminology Overview

Covered Core Package version:
2.1 + EDR

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Revision History

The Revision History is shown in the “Appendix” on page 41 [vol. 0].

Contributors

The persons who contributed to this specification are listed in the Appendix.

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1 GENERAL DESCRIPTION

Bluetooth wireless technology is a short-range communications system intended to replace the cable(s) connecting portable and/or fixed electronic devices. The key features of Bluetooth wireless technology are robustness, low power, and low cost. Many features of the core specification are optional, allowing product differentiation.

The Bluetooth core system consists of an RF transceiver, baseband, and protocol stack. The system offers services that enable the connection of devices and the exchange of a variety of classes of data between these devices.

This chapter of the specification provides an overview of the Bluetooth system architecture, communication topologies and data transport features. The text in this chapter of the specification should be treated as informational and used as a background and for context-setting.

1.1 OVERVIEW OF OPERATION

The Bluetooth RF (physical layer) operates in the unlicensed ISM band at 2.4 GHz. The system employs a frequency hop transceiver to combat interference and fading and provides many FHSS carriers. RF operation uses a shaped, binary frequency modulation to minimize transceiver complexity. The symbol rate is 1 Megasymbol per second (Ms/s) supporting the bit rate of 1 Megabit per second (Mb/s) or, with Enhanced Data Rate, a gross air bit rate of 2 or 3Mb/s. These modes are known as Basic Rate and Enhanced Data Rate respectively.

During typical operation a physical radio channel is shared by a group of devices that are synchronized to a common clock and frequency hopping pattern. One device provides the synchronization reference and is known as the master. All other devices are known as slaves. A group of devices synchronized in this fashion form a piconet. This is the fundamental form of communication in the Bluetooth wireless technology.

Devices in a piconet use a specific frequency hopping pattern, which is algorithmically determined by certain fields in the Bluetooth address and clock of the master. The basic hopping pattern is a pseudo-random ordering of the 79 frequencies in the ISM band. The hopping pattern may be adapted to exclude a portion of the frequencies that are used by interfering devices. The adaptive hopping technique improves Bluetooth co-existence with static (non-hopping) ISM systems when these are co-located.

The physical channel is sub-divided into time units known as slots. Data is transmitted between Bluetooth devices in packets, that are positioned in these slots. When circumstances permit, a number of consecutive slots may be allocated to a single packet. Frequency hopping takes place between the transmis-

sion or reception of packets. Bluetooth technology provides the effect of full duplex transmission through the use of a Time-Division Duplex (TDD) scheme.

Above the physical channel there is a layering of links and channels and associated control protocols. The hierarchy of channels and links from the physical channel upwards is physical channel, physical link, logical transport, logical link and L2CAP channel. These are discussed in more detail in Section 3.3 on page 33 to Section 3.6 on page 49 but are introduced here to aid the understanding of the remainder of this section.

Within a physical channel, a physical link is formed between any two devices that transmit packets in either direction between them. In a piconet physical channel there are restrictions on which devices may form a physical link. There is a physical link between each slave and the master. Physical links are not formed directly between the slaves in a piconet.

The physical link is used as a transport for one or more logical links that support unicast synchronous, asynchronous and isochronous traffic, and broadcast traffic. Traffic on logical links is multiplexed onto the physical link by occupying slots assigned by a scheduling function in the resource manager.

A control protocol for the baseband and physical layers is carried over logical links in addition to user data. This is the link manager protocol (LMP). Devices that are active in a piconet have a default asynchronous connection-oriented logical transport that is used to transport the LMP protocol signaling. For historical reasons this is known as the ACL logical transport. The default ACL logical transport is the one that is created whenever a device joins a piconet. Additional logical transports may be created to transport synchronous data streams when this is required.

The Link Manager function uses LMP to control the operation of devices in the piconet and provide services to manage the lower architectural layers (radio layer and baseband layer). The LMP protocol is only carried on the default ACL logical transport and the default broadcast logical transport.

Above the baseband layer the L2CAP layer provides a channel-based abstraction to applications and services. It carries out segmentation and reassembly of application data and multiplexing and de-multiplexing of multiple channels over a shared logical link. L2CAP has a protocol control channel that is carried over the default ACL logical transport. Application data submitted to the L2CAP protocol may be carried on any logical link that supports the L2CAP protocol.

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1.2 NOMENCLATURE

Where the following terms appear in the specification they have the meaning given in Table 1.1 on page 15.

Ad Hoc Network	A network typically created in a spontaneous manner. An ad hoc network requires no formal infrastructure and is limited in temporal and spatial extent.
Active Slave Broadcast (ASB)	The Active Slave Broadcast logical transport that is used to transport L2CAP user traffic to all active devices in the piconet. See Section 3.5.7 on page 45
Beacon Train	A pattern of reserved slots within a basic or adapted piconet physical channel. Transmissions starting in these slots are used to resynchronize parked devices.
Bluetooth	Bluetooth is a wireless communication link, operating in the unlicensed ISM band at 2.4 GHz using a frequency hopping transceiver. It allows real-time AV and data communications between Bluetooth Hosts. The link protocol is based on time slots.
Bluetooth Baseband	The part of the Bluetooth system that specifies or implements the medium access and physical layer procedures to support the exchange of real-time voice, data information streams, and ad hoc networking between Bluetooth Devices.
Bluetooth Clock	A 28 bit clock internal to a Bluetooth controller sub-system that ticks every 312.5µs. The value of this clock defines the slot numbering and timing in the various physical channels.
Bluetooth Controller	A sub-system containing the Bluetooth RF, baseband, resource controller, link manager, device manager and a Bluetooth HCI.
Bluetooth Device	A Bluetooth Device is a device that is capable of short-range wireless communications using the Bluetooth system.
Bluetooth Device Address	A 48 bit address used to identify each Bluetooth device.
BD_ADDR	The Bluetooth Device Address, BD_ADDR, is used to identify a Bluetooth device.
Bluetooth HCI	The Bluetooth Host Controller Interface (HCI) provides a command interface to the baseband controller and link manager and access to hardware status and control registers. This interface provides a uniform method of accessing the Bluetooth baseband capabilities.

Table 1.1: Nomenclature.

Bluetooth Host	A Bluetooth Host is a computing device, peripheral, cellular telephone, access point to PSTN network or LAN, etc. A Bluetooth Host attached to a Bluetooth Controller may communicate with other Bluetooth Hosts attached to their Bluetooth Controllers as well.
Channel	Either a physical channel or an L2CAP channel, depending on the context.
Connect (to service)	The establishment of a connection to a service. If not already done, this also includes establishment of a physical link, logical transport, logical link and L2CAP channel.
Connectable device	A Bluetooth device in range that periodically listens on its page scan physical channel and will respond to a page on that channel.
Connected devices	Two Bluetooth devices in the same piconet and with a physical link between them.
Connecting	A phase in the communication between devices when a connection between them is being established. (Connecting phase follows after the link establishment phase is completed.)
Connection	A connection between two peer applications or higher layer protocols mapped onto an L2CAP channel.
Connection establishment	A procedure for creating a connection mapped onto a channel.
Coverage area	The area where two Bluetooth devices can exchange messages with acceptable quality and performance.
Creation of a secure connection	A procedure of establishing a connection, including authentication and encryption.
Creation of a trusted relationship	A procedure where the remote device is marked as a trusted device. This includes storing a common link key for future authentication and pairing (if the link key is not available).
Device discovery	A procedure for retrieving the Bluetooth device address, clock, class-of-device field and used page scan mode from discoverable devices.
Discoverable device	A Bluetooth device in range that periodically listens on an inquiry scan physical channel and will respond to an inquiry on that channel. Discoverable device are normally also connectable.
Inquiring device	A Bluetooth device that is carrying out the inquiry procedure.
Inquiry	A procedure where a Bluetooth device transmits inquiry messages and listens for responses in order to discover the other Bluetooth devices that are within the coverage area.

Table 1.1: Nomenclature.

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Inquiry scan	A procedure where a Bluetooth device listens for inquiry messages received on its inquiry scan physical channel.
Interoperability	The ability of two or more systems or components to exchange information and to use the information that has been exchanged.
Isochronous data	Information in a stream where each information entity in the stream is bound by a time relationship to previous and successive entities.
Known device	A Bluetooth device for which at least the BD_ADDR is stored.
L2CAP Channel	A logical connection on L2CAP level between two devices serving a single application or higher layer protocol.
L2CAP Channel establishment	A procedure for establishing a logical connection on L2CAP level.
Link establishment	A procedure for establishing the default ACL link and hierarchy of links and channels between devices.
Link	Shorthand for a logical link.
Link key	A secret key that is known by two devices and is used in order to authenticate each device to the other
LMP authentication	An LMP level procedure for verifying the identity of a remote device.
LMP pairing	A procedure that authenticates two devices and creates a common link key that can be used as a basis for a trusted relationship or a (single) secure connection.
Logical link	The lowest architectural level used to offer independent data transport services to clients of the Bluetooth system.
Logical transport	Used in Bluetooth to represent commonality between different logical links due to shared acknowledgement protocol and link identifiers.
Name discovery	A procedure for retrieving the user-friendly name (the Bluetooth device name) of a connectable device.
Packet	Format of aggregated bits that are transmitted on a physical channel.
Page	The initial phase of the connection procedure where a device transmits a train of page messages until a response is received from the target device or a time-out occurs.
Page scan	A procedure where a device listens for page messages received on its page scan physical channel.
Paging device	A Bluetooth device that is carrying out the page procedure.

Table 1.1: Nomenclature.

Paired device	A Bluetooth device with which a link key has been exchanged (either before connection establishment was requested or during connecting phase).
Parked device	A device operating in a basic mode piconet that is synchronized to the master but has given up its default ACL logical transport.
Passkey	A 6-digit number used to authenticate connections when Secure Simple Pairing is used.
Physical Channel	Characterized by synchronized occupancy of a sequence of RF carriers by one or more devices. A number of physical channel types exist with characteristics defined for their different purposes.
Physical Link	A Baseband-level connection between two devices established using paging.
Piconet	A collection of devices occupying a shared physical channel where one of the devices is the Piconet Master and the remaining devices are connected to it.
Piconet Physical Channel	A Channel that is divided into time slots in which each slot is related to an RF hop frequency. Consecutive hops normally correspond to different RF hop frequencies and occur at a standard hop rate of 1600 hops/s. These consecutive hops follow a pseudo-random hopping sequence, hopping through a 79 RF channel set, or optionally fewer channels when Adaptive Frequency Hopping (AFH) is in use.
Piconet Master	The device in a piconet whose Bluetooth Clock and Bluetooth Device Address are used to define the piconet physical channel characteristics.
Piconet Slave	Any device in a piconet that is not the Piconet Master, but is connected to the Piconet Master.
PIN	A user-friendly number that can be used to authenticate connections to a device before pairing has taken place.
PMP	A Participant in Multiple Piconets. A device that is concurrently a member of more than one piconet, which it achieves using time division multiplexing (TDM) to interleave its activity on each piconet physical channel.
The Parked Slave Broadcast (PSB)	The Parked Slave Broadcast logical transport that is used for communications between the master and parked devices. Section 3.5.8 on page 46.
Scatternet	Two or more piconets that include one or more devices acting as PMPs.
Service Layer Protocol	A protocol that uses an L2CAP channel for transporting PDUs.

Table 1.1: Nomenclature.

Service discovery	Procedures for querying and browsing for services offered by or through another Bluetooth device.
Silent device	A Bluetooth device appears as silent to a remote device if it does not respond to inquiries made by the remote device.
Unknown device	A Bluetooth device for which no information (Bluetooth Device Address, link key or other) is stored.

Table 1.1: Nomenclature.

2 CORE SYSTEM ARCHITECTURE

The Bluetooth core system covers the four lowest layers and associated protocols defined by the Bluetooth specification as well as one common service layer protocol, the Service Discovery Protocol (SDP) and the overall profile requirements are specified in the Generic Access Profile (GAP). A complete Bluetooth application requires a number of additional service and higher layer protocols that are defined in the Bluetooth specification, but are not described here. The core system architecture is shown in Figure 2.1 on page 20 except for SDP that is not shown for clarity.

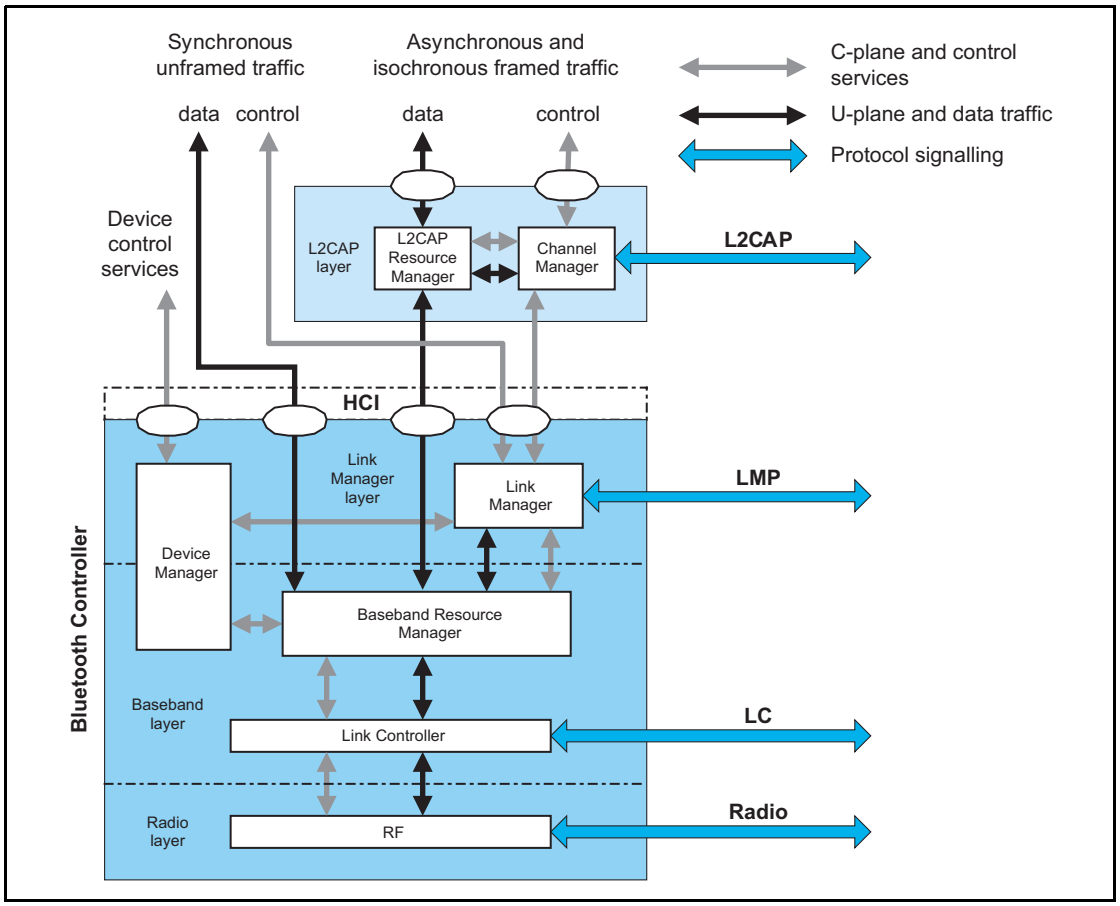


Figure 2.1: Bluetooth core system architecture

Figure 2.1 on page 20 shows the four lowest layers, each with its associated communication protocol. The lowest three layers are sometimes grouped into a subsystem known as the Bluetooth controller. This is a common implementation involving a standard physical communications interface between the Bluetooth controller and remainder of the Bluetooth system including the L2CAP, service and higher layers (known as the Bluetooth host.) Although this interface is optional the architecture is designed to allow for its existence and characteristics. The Bluetooth specification enables interoperability between independent Bluetooth systems by defining the protocol messages exchanged between equivalent layers, and also interoperability between independent

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Bluetooth subsystems by defining a common interface between Bluetooth controllers and Bluetooth hosts.

A number of functional blocks and the path of services and data between them are shown. The functional blocks shown in the diagram are informative; in general the Bluetooth specification does not define the details of implementations except where this is required for interoperability. Thus the functional blocks in Figure 2.1 on page 20 are shown in order to aid description of the system behavior. An implementation may be different from the system shown in Figure 2.1 on page 20.

Standard interactions are defined for all inter-device operation, where Bluetooth devices exchange protocol signaling according to the Bluetooth specification. The Bluetooth core system protocols are the Radio (RF) protocol, Link Control (LC) protocol, Link Manager (LM) protocol and Logical Link Control and Adaptation protocol (L2CAP), all of which are fully defined in subsequent parts of the Bluetooth specification. In addition the Service Discovery Protocol (SDP) is a service layer protocol required by all Bluetooth applications.

The Bluetooth core system offers services through a number of service access points that are shown in the diagram as ellipses. These services consist of the basic primitives that control the Bluetooth core system. The services can be split into three types. There are device control services that modify the behavior and modes of a Bluetooth device, transport control services that create, modify and release traffic bearers (channels and links), and data services that are used to submit data for transmission over traffic bearers. It is common to consider the first two as belonging to the C-plane and the last as belonging to the U-plane.

A service interface to the Bluetooth controller subsystem is defined such that the Bluetooth controller may be considered a standard part. In this configuration the Bluetooth controller operates the lowest three layers and the L2CAP layer is contained with the rest of the Bluetooth application in a host system. The standard interface is called the Host to Controller Interface (HCI) and its service access points are represented by the ellipses on the upper edge of the Bluetooth controller subsystem in Figure 2.1 on page 20. Implementation of this standard service interface is optional.

As the Bluetooth architecture is defined with the possibility of separate host and controller communicating through an HCI, a number of general assumptions are made. The Bluetooth controller is assumed to have limited data buffering capabilities in comparison with the host. Therefore the L2CAP layer is expected to carry out some simple resource management when submitting L2CAP PDUs to the controller for transport to a peer device. This includes segmentation of L2CAP SDUs into more manageable PDUs and then the fragmentation of PDUs into start and continuation packets of a size suitable for the controller buffers, and management of the use of controller buffers to ensure availability for channels with Quality of Service (QoS) commitments.

The Baseband layer provides the basic ARQ protocol in Bluetooth. The L2CAP layer can optionally provide a further error detection and retransmission to the L2CAP PDUs. This feature is recommended for applications with requirements for a low probability of undetected errors in the user data. A further optional feature of L2CAP is a window-based flow control that can be used to manage buffer allocation in the receiving device. Both of these optional features augment the QoS performance in certain scenarios.

Although these assumptions may not be required for embedded Bluetooth implementations that combine all layers in a single system, the general architectural and QoS models are defined with these assumptions in mind, in effect a lowest common denominator.

Automated conformance testing of implementations of the Bluetooth core system is required. This is achieved by allowing the tester to control the implementation through the RF interface, which is common to all Bluetooth systems, and through the Test Control Interface (TCI), which is only required for conformance testing.

The tester uses exchanges with the Implementation Under Test (IUT) through the RF interface to ensure the correct responses to requests from remote devices. The tester controls the IUT through the TCI to cause the IUT to originate exchanges through the RF interface so that these can also be verified as conformant.

The TCI uses a different command-set (service interface) for the testing of each architectural layer and protocol. A subset of the HCI command-set is used as the TCI service interface for each of the layers and protocols within the Bluetooth Controller subsystem. A separate service interface is used for testing the L2CAP layer and protocol. As an L2CAP service interface is not defined in the Bluetooth core specification it is defined separately in the Test Control Interface specification. Implementation of the L2CAP service interface is only required for conformance testing.

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2.1 CORE ARCHITECTURAL BLOCKS

This section describes the function and responsibility of each of the blocks shown in Figure 2.1 on page 20. An implementation is not required to follow the architecture described above, though every implementation shall conform to the protocol specifications described in subsequent parts of the Bluetooth specification, and shall implement the behavioral aspects of the system outlined below and specified in subsequent parts of the Bluetooth specification.

2.1.1 Channel manager

The channel manager is responsible for creating, managing and destroying L2CAP channels for the transport of service protocols and application data streams. The channel manager uses the L2CAP protocol to interact with a channel manager on a remote (peer) device to create these L2CAP channels and connect their endpoints to the appropriate entities. The channel manager interacts with its local link manager to create new logical links (if necessary) and to configure these links to provide the required quality of service for the type of data being transported.

2.1.2 L2CAP resource manager

The L2CAP resource manager block is responsible for managing the ordering of submission of PDU fragments to the baseband and some relative scheduling between channels to ensure that L2CAP channels with QoS commitments are not denied access to the physical channel due to Bluetooth controller resource exhaustion. This is required because the architectural model does not assume that the Bluetooth controller has limitless buffering, or that the HCI is a pipe of infinite bandwidth.

L2CAP Resource Managers may also carry out traffic conformance policing to ensure that applications are submitting L2CAP SDUs within the bounds of their negotiated QoS settings. The general Bluetooth data transport model assumes well-behaved applications, and does not define how an implementation is expected to deal with this problem.

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2.1.3 Device manager

The device manager is the functional block in the baseband that controls the general behavior of the Bluetooth device. It is responsible for all operation of the Bluetooth system that is not directly related to data transport, such as inquiring for the presence of other nearby Bluetooth devices, connecting to other Bluetooth devices, or making the local Bluetooth device discoverable or connectable by other devices.

The device manager requests access to the transport medium from the baseband resource controller in order to carry out its functions.

The device manager also controls local device behavior implied by a number of the HCI commands, such as managing the device local name, any stored link keys, and other functionality.

2.1.4 Link manager

The link manager is responsible for the creation, modification and release of logical links (and, if required, their associated logical transports), as well as the update of parameters related to physical links between devices. The link manager achieves this by communicating with the link manager in remote Bluetooth devices using the Link Management Protocol (LMP.)

The LM protocol allows the creation of new logical links and logical transports between devices when required, as well as the general control of link and transport attributes such as the enabling of encryption on the logical transport, the adapting of transmit power on the physical link, or the adjustment of QoS settings for a logical link.

2.1.5 Baseband resource manager

The baseband resource manager is responsible for all access to the radio medium. It has two main functions. At its heart is a scheduler that grants time on the physical channels to all of the entities that have negotiated an access contract. The other main function is to negotiate access contracts with these entities. An access contract is effectively a commitment to deliver a certain QoS that is required in order to provide a user application with an expected performance.

The access contract and scheduling function must take account of any behavior that requires use of the Bluetooth radio. This includes (for example) the normal exchange of data between connected devices over logical links, and logical transports, as well as the use of the radio medium to carry out inquiries, make connections, be discoverable or connectable, or to take readings from unused carriers during the use of adaptive frequency hopping mode.

In some cases the scheduling of a logical link results in changing to a different physical channel from the one that was previously used. This may be (for example) due to involvement in scatternet, a periodic inquiry function, or page

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scanning. When the physical channels are not time slot aligned, then the resource manager also accounts for the realignment time between slots on the original physical channel and slots on the new physical channel. In some cases the slots will be naturally aligned due to the same device clock being used as a reference for both physical channels.

2.1.6 Link controller

The link controller is responsible for the encoding and decoding of Bluetooth packets from the data payload and parameters related to the physical channel, logical transport and logical link.

The link controller carries out the link control protocol signaling (in close conjunction with the scheduling function of the resource manager), which is used to communicate flow control and acknowledgement and retransmission request signals. The interpretation of these signals is a characteristic of the logical transport associated with the baseband packet. Interpretation and control of the link control signaling is normally associated with the resource manager's scheduler.

2.1.7 RF

The RF block is responsible for transmitting and receiving packets of information on the physical channel. A control path between the baseband and the RF block allows the baseband block to control the timing and frequency carrier of the RF block. The RF block transforms a stream of data to and from the physical channel and the baseband into required formats.

3 DATA TRANSPORT ARCHITECTURE

The Bluetooth data transport system follows a layered architecture. This description of the Bluetooth system describes the Bluetooth core transport layers up to and including L2CAP channels. All Bluetooth operational modes follow the same generic transport architecture, which is shown in Figure 3.1 on page 26.

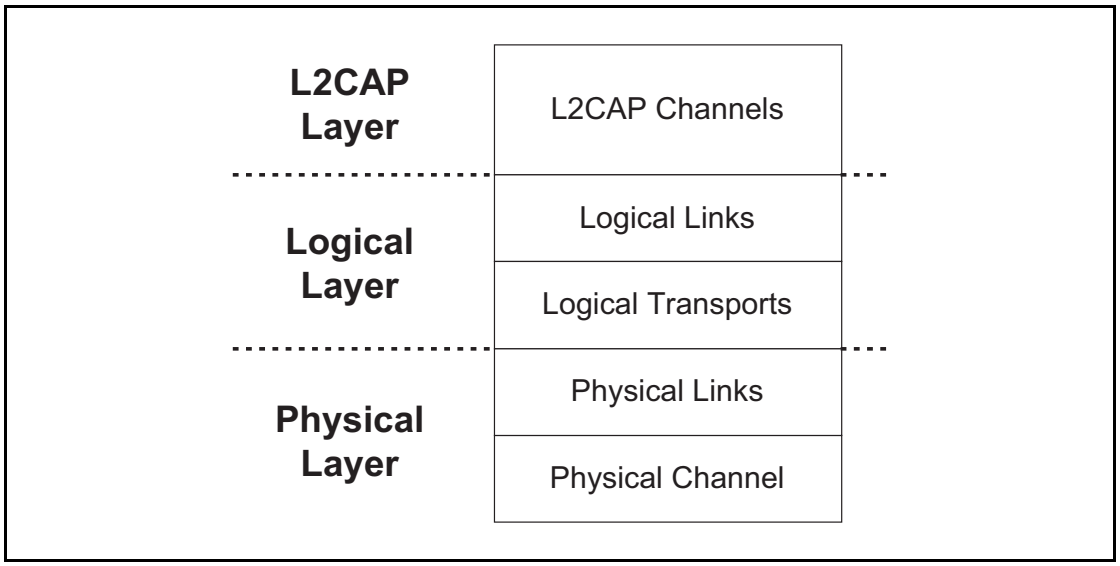


Figure 3.1: Bluetooth generic data transport architecture

For efficiency and legacy reasons, the Bluetooth transport architecture includes a sub-division of the logical layer, distinguishing between logical links and logical transports. This sub-division provides a general (and commonly understood) concept of a logical link that provides an independent transport between two or more devices. The logical transport sub-layer is required to describe the inter-dependence between some of the logical link types (mainly for reasons of legacy behavior.)

Previous versions of the Bluetooth specification described the ACL and SCO links as physical links. With the addition of Extended SCO (eSCO) and for future expansion it is better to consider these as logical transport types, which more accurately encapsulates their purpose. However, they are not as independent as might be desired, due to their shared use of resources such as the LT_ADDR and acknowledgement/repeat request (ARQ) scheme. Hence the architecture is incapable of representing these logical transports with a single transport layer. The additional logical transport layer goes some way towards describing this behavior.

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3.1 CORE TRAFFIC BEARERS

The Bluetooth core system provides a number of standard traffic bearers for the transport of service protocol and application data. These are shown in Figure 3.2 on page 27 below (for ease of representation this is shown with higher layers to the left and lower layers to the right).

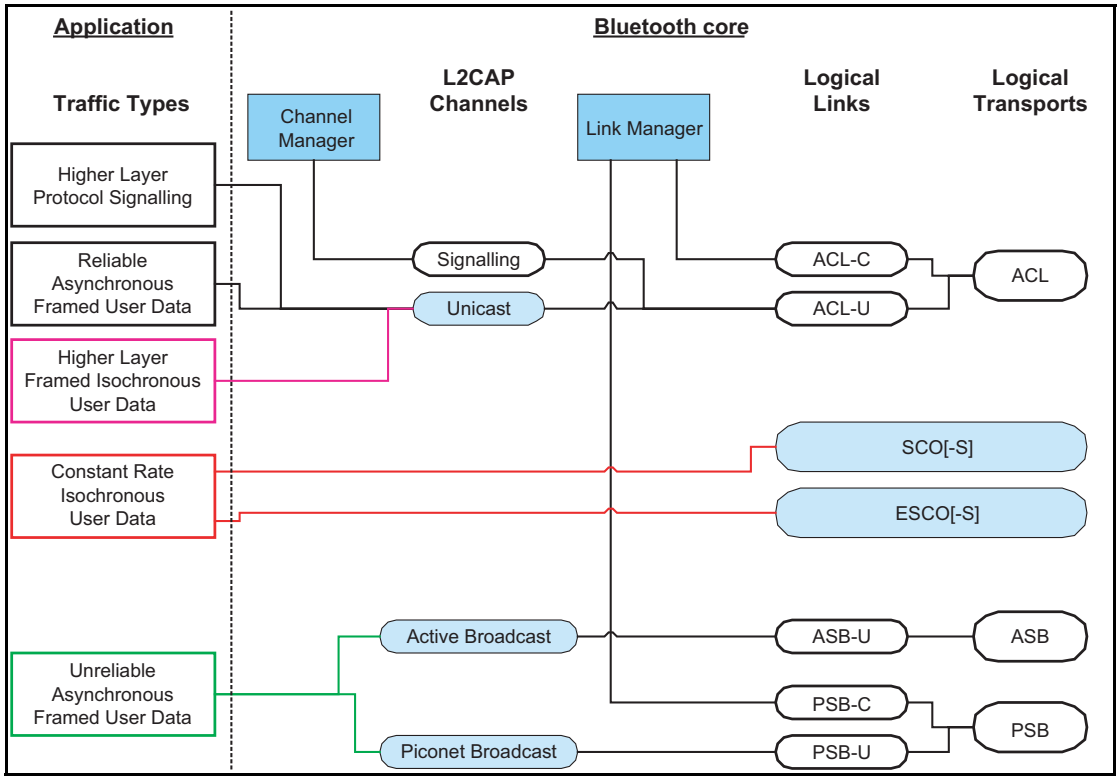


Figure 3.2: Bluetooth traffic bearers

The core traffic bearers that are available to applications are shown in Figure 3.2 on page 27 as the shaded rounded rectangles. The architectural layers that are defined to provide these services are described in Section 2 on page 20. A number of data traffic types are shown on the left of the diagram linked to the traffic bearers that are typically suitable for transporting that type of data traffic.

The logical links are named using the names of the associated logical transport and a suffix that indicates the type of data that is transported. (C for control links carrying LMP messages, U for L2CAP links carrying user data (L2CAP PDUs) and S for stream links carrying unformatted synchronous or isochronous data.) It is common for the suffix to be removed from the logical link without introducing ambiguity, thus a reference to the default ACL logical transport can be resolved to mean the ACL-C logical link in cases where the LMP protocol is being discussed, or the ACL-U logical link when the L2CAP layer is being discussed.

The mapping of application traffic types to Bluetooth core traffic bearers in Figure 3.2 on page 27 is based on matching the traffic characteristics with the

bearer characteristics. It is recommended to use these mappings as they provide the most natural and efficient method of transporting the data with its given characteristics.

However, an application (or an implementation of the Bluetooth core system) may choose to use a different traffic bearer, or a different mapping to achieve a similar result. For example, in a piconet with only one slave, the master may choose to transport L2CAP broadcasts over the ACL-U logical link rather than over the ASB-U or PSB-U logical links. This will probably be more efficient in terms of bandwidth (if the physical channel quality is not too degraded.) Use of alternative transport paths to those in Figure 3.2 on page 27 is only acceptable if the characteristics of the application traffic type are preserved.

Figure 3.2 on page 27 shows a number of application traffic types. These are used to classify the types of data that may be submitted to the Bluetooth core system. The original data traffic type may not be the same as the type that is submitted to the Bluetooth core system if an intervening process modifies it. For example, video data is generated at a constant rate but an intermediate coding process may alter this to variable rate, e.g. by MPEG4 encoding. For the purposes of the Bluetooth core system, only the characteristic of the submitted data is of interest.

3.1.1 Framed data traffic

The L2CAP layer services provide a frame-oriented transport for asynchronous and isochronous user data. The application submits data to this service in variable-sized frames (up to a negotiated maximum for the channel) and these frames are delivered in the same form to the corresponding application on the remote device. There is no requirement for the application to insert additional framing information into the data, although it may do so if this is required (such framing is invisible to the Bluetooth core system.)

Connection-oriented L2CAP channels may be created for transport of unicast (point-to-point) data between two Bluetooth devices. A connectionless L2CAP channel exists for broadcasting data. In the case of piconet topologies the master device is always the source of broadcast data and the slave device(s) are the recipients. Traffic on the broadcast L2CAP channel is uni-directional. Unicast L2CAP channels may be uni-directional or bi-directional.

L2CAP channels have an associated QoS setting that defines constraints on the delivery of the frames of data. These QoS settings may be used to indicate (for example) that the data is isochronous, and therefore has a limited lifetime after which it becomes invalid, or that the data should be delivered within a given time period, or that the data is reliable and should be delivered without error, however long this takes.

The L2CAP channel manager is responsible for arranging to transport the L2CAP channel data frames on an appropriate baseband logical link, possibly multiplexing this onto the baseband logical link with other L2CAP channels with similar characteristics.

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3.1.2 Unframed data traffic

If the application does not require delivery of data in frames, possibly because it includes in-stream framing, or because the data is a pure stream, then it may avoid the use of L2CAP channels and make direct use of a baseband logical link.

The Bluetooth core system supports the direct transport of application data that is isochronous and of a constant rate (either bit-rate, or frame-rate for pre-framed data), using a SCO-S or eSCO-S logical link. These logical links reserve physical channel bandwidth and provide a constant rate transport locked to the piconet clock. Data is transported in fixed size packets at fixed intervals with both of these parameters negotiated during channel establishment. eSCO links provide a greater choice of bit-rates and also provide greater reliability by using limited retransmission in case of error. Enhanced Data Rate operation is supported for eSCO, but not for SCO logical transports. SCO and eSCO logical transports do not support multiplexed logical links or any further layering within the Bluetooth core. An application may choose to layer a number of streams within the submitted SCO/eSCO stream, provided that the submitted stream is, or has the appearance of being, a constant rate stream.

The application chooses the most appropriate type of logical link from those available at the baseband, and creates and configures it to transport the data stream, and releases it when completed. (The application will normally also use a framed L2CAP unicast channel to transport its C-plane information to the peer application on the remote device.)

If the application data is isochronous and of a variable rate, then this may only be carried by the L2CAP unicast channel, and hence will be treated as framed data.

3.1.3 Reliability of traffic bearers

Bluetooth is a wireless communications system. In poor RF environments, this system should be considered inherently unreliable. To counteract this the system provides levels of protection at each layer. The baseband packet header uses forward error correcting (FEC) coding to allow error correction by the receiver and a header error check (HEC) to detect errors remaining after correction. Certain Baseband packet types include FEC for the payload. Furthermore, some Baseband packet types include a cyclic redundancy error check (CRC).

On ACL logical transports the results of the error detection algorithm are used to drive a simple acknowledgement/repeat request (ARQ) protocol. This provides an enhanced reliability by re-transmitting packets that do not pass the receiver’s error checking algorithm. It is possible to modify this scheme to support latency-sensitive packets by discarding an unsuccessfully transmitted packet at the transmitter if the packet’s useful life has expired. eSCO links use

a modified version of this scheme to improve reliability by allowing a limited number of retransmissions.

The resulting reliability gained by this ARQ scheme is only as dependable as the ability of the HEC and CRC codes to detect errors. In most cases this is sufficient, however it has been shown that for the longer packet types the probability of an undetected error is too high to support typical applications, especially those with a large amount of data being transferred.

The L2CAP layer provides an additional level of error control that is designed to detect the occasional undetected errors in the baseband layer and request retransmission of the affected data. This provides the level of reliability required by typical Bluetooth applications.

Broadcast links have no feedback route, and are unable to use the ARQ scheme (although the receiver is still able to detect errors in received packets.) Instead each packet is transmitted several times in the hope that the receiver is able to receive at least one of the copies successfully. Despite this approach there are still no guarantees of successful receipt, and so these links are considered unreliable.

In summary, if a link or channel is characterized as reliable this means that the receiver is capable of detecting errors in received packets and requesting retransmission until the errors are removed. Due to the error detection system used some residual (undetected) errors may still remain in the received data. For L2CAP channels the level of these is comparable to other communication systems, although for logical links the residual error level is somewhat higher.

The transmitter may remove packets from the transmit queue such that the receiver does not receive all the packets in the sequence. If this happens detection of the missing packets is delegated to the L2CAP layer.

On an unreliable link the receiver is capable of detecting errors in received packets but cannot request retransmission. The packets passed on by the receiver may be without error, but there is no guarantee that all packets in the sequence are received. Hence the link is considered fundamentally unreliable. There are limited uses for such links, and these uses are normally dependent on the continuous repetition of data from the higher layers while it is valid.

Stream links have a reliability characteristic somewhere between a reliable and an unreliable link, depending on the current operating conditions.

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3.2 TRANSPORT ARCHITECTURE ENTITIES

The Bluetooth transport architecture entities are shown in Figure 3.3 on page 31 and are described from the lowest layer upwards in the subsequent sections.

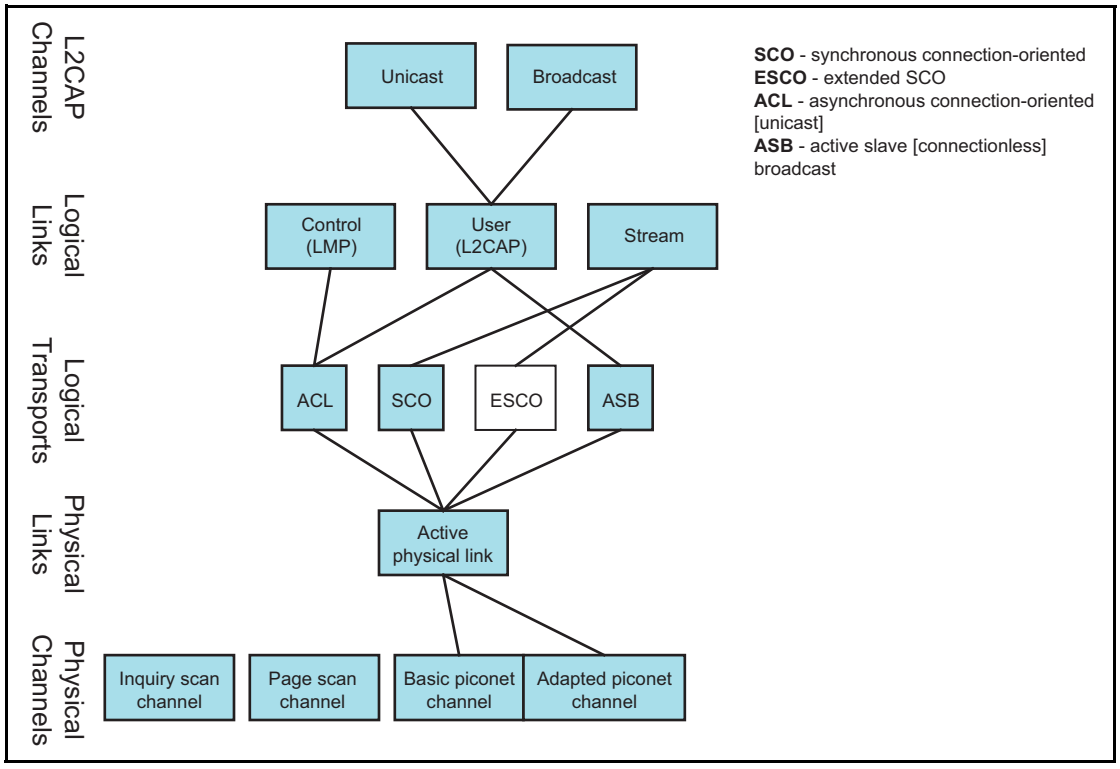


Figure 3.3: Overview of transport architecture entities and hierarchy

3.2.1 Bluetooth generic packet structure

The general packet structure nearly reflects the architectural layers found in the Bluetooth system. The packet structure is designed for optimal use in normal operation. It is shown in Figure 3.4 on page 31.

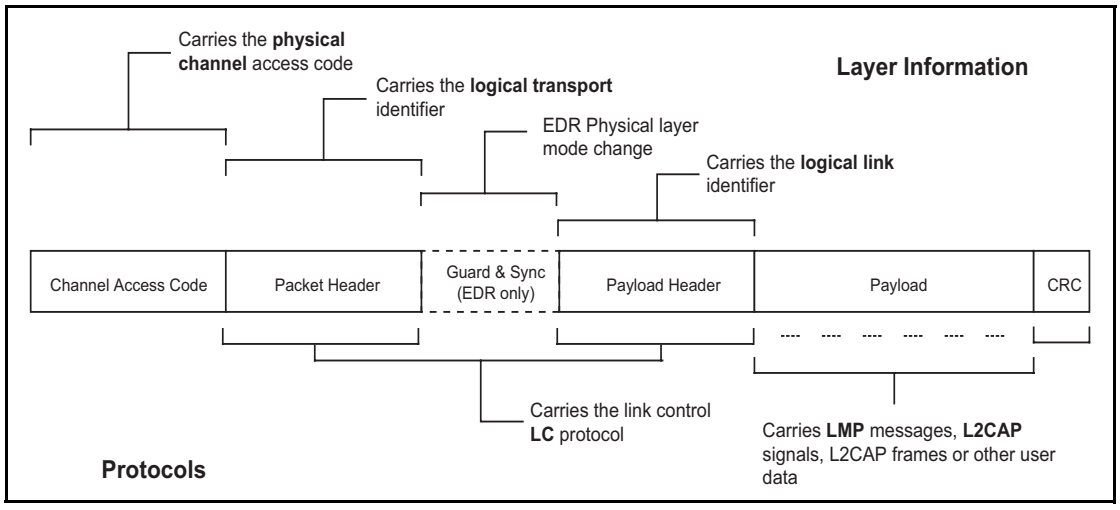


Figure 3.4: Bluetooth packet structure

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Packets normally only include the fields that are necessary to represent the layers required by the transaction. Thus a simple inquiry request over an inquiry scan physical channel does not create or require a logical link or higher layer and therefore consists only of the channel access code (associated with the physical channel.) General communication within a piconet uses packets that include all of the fields, as all of the architectural layers are used.

All packets include the channel access code. This is used to identify communications on a particular physical channel, and to exclude or ignore packets on a different physical channel that happens to be using the same RF carrier in physical proximity.

There is no direct field within the Bluetooth packet structure that represents or contains information relating to physical links. This information is implied in the logical transport address (LT_ADDR) carried in the packet header.

Most packets include a packet header. The packet header is always present in packets transmitted on physical channels that support physical links, logical transports and logical links. The packet header carries the LT_ADDR, which is used by each receiving device to determine if the packet is addressed to the device and is used to route the packet internally.

The packet header also carries part of the link control (LC) protocol that is operated per logical transport (except for ACL and SCO transports that operate a shared LC protocol carried on either logical transport.)

The Enhanced Data Rate (EDR) packets have a guard time and synchronization sequence before the payload. This is a field used for physical layer change of modulation scheme.

The payload header is present in all packets on logical transports that support multiple logical links. The payload header includes a logical link identifier field used for routing the payload, and a field indicating the length of the payload. Some packet types also include a CRC after the packet payload that is used to detect most errors in received packets.

EDR packets have a trailer after the CRC.

The packet payload is used to transport the user data. The interpretation of this data is dependent on the logical transport and logical link identifiers. For ACL logical transports Link Manager Protocol (LMP) messages and L2CAP signals are transported in the packet payload, along with general user data from applications. For SCO and eSCO logical transports the payload contains the user data for the logical link.

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3.3 PHYSICAL CHANNELS

The lowest architectural layer in the Bluetooth system is the physical channel. A number of types of physical channel are defined. All Bluetooth physical channels are characterized by an RF frequency combined with temporal parameters and restricted by spatial considerations. For the basic and adapted piconet physical channels frequency hopping is used to change frequency periodically to reduce the effects of interference and for regulatory reasons.

Two Bluetooth devices use a shared physical channel for communication. To achieve this their transceivers need to be tuned to the same RF frequency at the same time, and they need to be within a nominal range of each other.

Given that the number of RF carriers is limited and that many Bluetooth devices may be operating independently within the same spatial and temporal area there is a strong likelihood of two independent Bluetooth devices having their transceivers tuned to the same RF carrier, resulting in a physical channel collision. To mitigate the unwanted effects of this collision each transmission on a physical channel starts with an access code that is used as a correlation code by devices tuned to the physical channel. This channel access code is a property of the physical channel. The access code is always present at the start of every transmitted packet.

Four Bluetooth physical channels are defined. Each is optimized and used for a different purpose. Two of these physical channels (the basic piconet channel and adapted piconet channel) are used for communication between connected devices and are associated with a specific piconet. The remaining physical channels are used for discovering Bluetooth devices (the inquiry scan channel) and for connecting Bluetooth devices (the page scan channel.)

A Bluetooth device can only use one of these physical channels at any given time. In order to support multiple concurrent operations the device uses time-division multiplexing between the channels. In this way a Bluetooth device can appear to operate simultaneously in several piconets, as well as being discoverable and connectable.

Whenever a Bluetooth device is synchronized to the timing, frequency and access code of a physical channel it is said to be ‘connected’ to this channel (whether or not it is actively involved in communications over the channel.) The Bluetooth specification assumes that a device is only capable of connecting to one physical channel at any time. Advanced devices be may capable of connecting simultaneously to more than one physical channel, but the specification does not assume that this is possible.

3.3.1 Basic piconet channel

3.3.1.1 Overview

The basic piconet channel is used for communication between connected devices during normal operation.

3.3.1.2 Characteristics

The basic piconet channel is characterized by a pseudo-random sequence hopping through the RF channels. The hopping sequence is unique for the piconet and is determined by the Bluetooth device address of the master. The phase in the hopping sequence is determined by the Bluetooth clock of the master. All Bluetooth devices participating in the piconet are time- and hop-synchronized to the channel.

The channel is divided into time slots where each slot corresponds to an RF hop frequency. Consecutive hops correspond to different RF hop frequencies. The time slots are numbered according to the Bluetooth clock of the piconet master. Packets are transmitted by Bluetooth devices participating in the piconet aligned to start at a slot boundary. Each packet starts with the channel's access code, which is derived from the Bluetooth device address of the piconet.

On the basic piconet channel the master controls access to the channel. The master starts its transmission in even-numbered time slots only. Packets transmitted by the master are aligned with the slot start and define the piconet timing. Packets transmitted by the master may occupy up to five time slots depending on the packet type.

Each master transmission is a packet carrying information on one of the logical transports. Slave devices may transmit on the physical channel in response. The characteristics of the response are defined by the logical transport that is addressed.

For example on the asynchronous connection-oriented logical transport the addressed slave device responds by transmitting a packet containing information for the same logical transport that is nominally aligned with the next (odd-numbered) slot start. Such a packet may occupy up to five time slots, depending on the packet type. On a broadcast logical transport no slaves are allowed to respond.

A special characteristic of the basic piconet physical channel is the use of some reserved slots to transmit a beacon train. The beacon train is only used if the piconet physical channel has parked slaves connected to it. In this situation the master transmits a packet in the reserved beacon train slots (these packets are used by the slave to resynchronize to the piconet physical channel.) The master may transmit packets from any logical transport onto these slots, providing there is a transmission starting in each of the slots. In the case where

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there is information from the parked slave broadcast (PSB) logical transport to be transmitted then this is transmitted in the beacon train slots and takes priority over any other logical transport.

3.3.1.3 Topology

A basic piconet channel may be shared by any number of Bluetooth devices, limited only by the resources available on the piconet master device. Only one device is the piconet master, all others being piconet slaves. All communication is between the master and slave devices. There is no direct communication between slave devices on the piconet channel.

There is, however, a limitation on the number of logical transports that can be supported within a piconet. This means that although there is no theoretical limit to the number of Bluetooth devices that share a channel there is a limit to the number of these devices that can be actively involved in exchanging data with the master.

3.3.1.4 Supported layers

The basic piconet channel supports a number of physical links, logical transports, logical links and L2CAP channels used for general purpose communications.

3.3.2 Adapted piconet channel

3.3.2.1 Overview

The adapted piconet channel differs from the basic piconet channel in two ways. First the frequencies on which the slaves transmit are the same as the preceding master transmit frequency. In other words the frequency is not recomputed between master and subsequent slave packets. The second way in which the adapted piconet channel differs from the basic piconet channel is that the adapted type can be based on fewer than the full 79 frequencies. A number of frequencies may be excluded from the hopping pattern by being marked as “unused”. The remainder of the 79 frequencies are included. The two sequences are the same except that whenever the basic pseudo-random hopping sequence would have selected an unused frequency it is replaced with an alternative chosen from the used set.

Because the adapted piconet channel uses the same timing and access code as the basic piconet channel, the two channels are often coincident. This provides a deliberate benefit as it allows slaves in either the basic piconet channel or the adapted piconet channel to adjust their synchronization to the master.

The topology and supported layers of the adapted piconet physical channel are identical to the basic piconet physical channel.

3.3.3 Inquiry scan channel

3.3.3.1 Overview

In order for a device to be discovered an inquiry scan channel is used. A discoverable device listens for inquiry requests on its inquiry scan channel and then sends responses to these requests. In order for a device to discover other devices, it iterates (hops) through all possible inquiry scan channel frequencies in a pseudo-random fashion, sending an inquiry request on each frequency and listening for any response.

3.3.3.2 Characteristics

Inquiry scan channels follow a slower hopping pattern and use an access code to distinguish between occasional occupancy of the same radio frequency by two co-located devices using different physical channels.

The access code used on the inquiry scan channel is taken from a reserved set of inquiry access codes that are shared by all Bluetooth devices. One access code is used for general inquiries, and a number of additional access codes are reserved for limited inquiries. Each device has access to a number of different inquiry scan channels. As all of these channels share an identical hopping pattern, a device may concurrently occupy more than one inquiry scan channel if it is capable of concurrently correlating more than one access code.

A device using one of its inquiry scan channels remains passive until it receives an inquiry message on this channel from another Bluetooth device. This is identified by the appropriate inquiry access code. The inquiry scanning device will then follow the inquiry response procedure to return a response to the inquiring device.

In order for a device to discover other Bluetooth devices it uses the inquiry scan channel of these devices in order to send inquiry requests. As it has no prior knowledge of the devices to discover, it cannot know the exact characteristics of the inquiry scan channel.

The device takes advantage of the fact that inquiry scan channels have a reduced number of hop frequencies and a slower rate of hopping. The inquiring device transmits inquiry requests on each of the inquiry scan hop frequencies and listens for an inquiry response. This is done at a faster rate, allowing the inquiring device to cover all inquiry scan frequencies in a reasonably short time period.

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3.3.3.3 Topology

Inquiring and discoverable devices use a simple exchange of packets to fulfill the inquiring function. The topology formed during this transaction is a simple and transient point-to-point connection.

3.3.3.4 Supported layers

During the exchange of packets between an inquiring and discoverable device it may be considered that a temporary physical link exists between these devices. However, the concept is quite irrelevant as it has no physical representation but is only implied by the brief transaction between the devices. No further architectural layers are considered to be supported.

3.3.4 Page scan channel

3.3.4.1 Overview

A connectable device (one that is prepared to accept connections) does so using a page scan channel. A connectable device listens for page requests on its page scan channel and enters into a sequence of exchanges with this device. In order for a device to connect to another device, it iterates (hops) through all page scan channel frequencies in a pseudo-random fashion, sending a page request on each frequency and listening for any response.

3.3.4.2 Characteristics

The page scan channel uses an access code derived from the scanning device's Bluetooth device address to identify communications on the channel. The page scan channel uses a slower hopping rate than the hop rate of the basic and adapted piconet channels. The hop selection algorithm uses the Bluetooth device clock of the scanning device as an input.

A device using its page scan channel remains passive until it receives a page request from another Bluetooth device. This is identified by the page scan channel access code. The two devices will then follow the page procedure to form a connection. Following a successful conclusion of the page procedure both devices switch to the basic piconet channel that is characterized by having the paging device as master.

In order for a device to connect to another Bluetooth device it uses the page scan channel of the target device in order to send page requests. If the paging device does not know the phase of the target device's page scan channel it therefore does not know the current hop frequency of the target device. The paging device transmits page requests on each of the page scan hop frequencies and listens for a page response. This is done at a faster hop rate, allowing

the paging device to cover all page scan frequencies in a reasonably short time period.

The paging device may have some knowledge of the target device's Bluetooth clock (indicated during a previous inquiry transaction between the two devices, or as a result of a previous involvement in a piconet with the device), in which case it is able to predict the phase of the target device's page scan channel. It may use this information to optimize the synchronization of the paging and page scanning process and speed up the formation of the connection.

3.3.4.3 Topology

Paging and connectable devices use a simple exchange of packets to fulfill the paging function. The topology formed during this transaction is a simple and transient point-to-point connection.

3.3.4.4 Supported layers

During the exchange of packets between a paging and connectable device it may be considered that a temporary physical link exists between these devices. However, the concept is quite irrelevant as it has no physical representation but is only implied by the brief transaction between the devices. No further architectural layers are considered to be supported.

3.4 PHYSICAL LINKS

A physical link represents a baseband connection between Bluetooth devices. A physical link is always associated with exactly one physical channel (although a physical channel may support more than one physical link.)

Within the Bluetooth system a physical link is a virtual concept that has no direct representation within the structure of a transmitted packet. The access code packet field, together with the clock and address of the master Bluetooth device are used to identify a physical channel. However there is no subsequent part of the packet that directly identifies the physical link. Instead the physical link may be identified by association with the logical transport, as each logical transport is only received on one physical link.

Some physical link types have properties that may be modified. An example of this is the transmit power for the link. Other physical link types have no such properties. In the case of physical links with modifiable properties the LM protocol is used to adapt these properties. As the LM protocol is supported at a higher layer (by a logical link) the appropriate physical link is identified by implication from the logical link that transports the LM signaling.

In the situation where a transmission is broadcast over a number of different physical links, then the transmission parameters are selected to be suitable for all of the physical links.

3.4.1 Links supported by the basic and adapted piconet physical channel

The basic and adapted piconet physical channels support a physical link which may be active or parked. The physical link is a point-to-point link between the master and a slave. It is always present when the slave is synchronized in the piconet.

3.4.1.1 Active physical link

The physical link between a master and a slave device is active if a default ACL logical transport exists between the devices. Active physical links have no direct identification of their own, but are identified by association with the default ACL logical transport ID with which there is a one-to-one correspondence.

An active physical link has the associated properties of radio transmit power in each direction. Transmissions from slave devices are always directed over the active physical link to the master, and use the transmit power that is a property of this link in the slave to master direction. Transmissions from the master may be directed over a single active physical link (to a specific slave) or over a number of physical links (to a group of slaves in the piconet.) In the case of point-to-point transmissions the master uses the appropriate transmit power for the physical link in question. (In the case of point-to-multipoint transmissions the master uses a transmit power appropriate for the set of devices addressed.)

Active physical links may be placed into Hold or Sniff mode. The effect of these modes is to modify the periods when the physical link is active and may carry traffic. Logical transports that have defined scheduling characteristics are not affected by these modes and continue according to their pre-defined scheduling behavior. The default ACL logical transport and other links with undefined scheduling characteristics are subject to the mode of the active physical link.

3.4.1.2 Parked physical link

The physical link between a master and a slave device is parked when the slave remains synchronized in the piconet, but has no default ACL logical transport. Such a slave is also said to be parked. A beacon train is used to provide regular synchronization to all parked slaves connected to the piconet physical channel. A parked slave broadcast (PSB) logical transport is used to allow communication of a subset of LMP signaling and broadcast L2CAP to parked slaves. The PSB logical transport is closely associated with the beacon train.

A slave is parked (its active link is changed to a parked link) using the park procedure. The master is not allowed to park a slave that has any user created logical transport supported by the physical link. These logical transports are first removed, and any L2CAP channels that are built on these logical transports are also removed. The broadcast logical transport and default ACL logical transports are not considered as user created and are not explicitly removed. When the active link is replaced with a parked link the default ACL

logical transport is implicitly removed. The supported logical links and L2CAP channels remain in existence, but become suspended. It is not possible to use these links and L2CAP channels to transport signaling or data while the active link is absent.

A parked slave may become active using the unpark procedure. This procedure is requested by the slave at an access window and initiated by the master. Following the unpark procedure the parked physical link is changed to an active physical link and the default ACL logical transport is re-created. L2CAP channels that were suspended during the most recent park procedure are associated with the new default ACL logical transport and become active again.

Parked links do not support radio power control, as there is no feedback path from parked slaves to the piconet master that can be used to signal received signal strength at the slave or for the master to measure received signal strength from the slave. Transmissions are carried out at nominal power on parked links.

Parked links use the same physical channel as their associated active link. If a master manages a piconet that contains parked slaves using the basic piconet physical channel and also parked slaves using the adapted piconet physical channel then it must create a parked slave broadcast logical transport (and associated transport) for each of these physical channels.

A parked slave may use the inactive periods of the parked slave broadcast logical transport to save power, or it may carry out activities on other physical channels unrelated to the piconet within which it is parked.

3.4.2 Links supported by the scanning physical channels

In the case of the inquiry scan and page scan channels the physical link exists for a relatively short time and cannot be controlled or modified in any way. These types of physical link are not further elaborated.

3.5 LOGICAL LINKS AND LOGICAL TRANSPORTS

A variety of logical links are available to support different application data transport requirements. Each logical link is associated with a logical transport, which has a number of characteristics. These characteristics include flow control, acknowledgement/repeat mechanisms, sequence numbering and scheduling behavior. Logical transports are able to carry different types of logical links (depending on the type of the logical transport). In the case of some of the Bluetooth logical links these are multiplexed onto the same logical transport. Logical transports may be carried by active physical links on either the basic or the adapted piconet physical channel.

Logical transport identification and real-time (link control) signaling are carried in the packet header, and for some logical links identification is carried in the

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payload header. Control signaling that does not require single slot response times is carried out using the LMP protocol.

Table 3.1 on page 41 lists all of the logical transport types, the supported logical link types, which type of physical links and physical channels can support them, and a brief description of the purpose of the logical transport.

Logical transport	Links supported	Supported by	Overview
Asynchronous Connection-Oriented (ACL)	Control (LMP) ACL-C User (L2CAP) ACL-U	Active physical link, basic or adapted physical channel	Reliable or time-bounded, bi-directional, point-to-point.
Synchronous Connection-Oriented (SCO)	Stream (unframed) SCO-S	Active physical link, basic or adapted physical channel	Bi-directional, symmetric, point-to-point, AV channels. Used for 64Kb/s constant rate data.
Extended Synchronous Connection-Oriented (eSCO)	Stream (unframed) eSCO-S	Active physical link, basic or adapted physical channel	Bi-directional, symmetric or asymmetric, point-to-point, general regular data, limited retransmission. Used for constant rate data synchronized to the master Bluetooth clock.
Active slave broadcast (ASB)	User (L2CAP) ASB-U	Active physical link, basic or adapted physical channel.	Unreliable, uni-directional broadcast to any devices synchronized with the physical channel. Used for broadcast L2CAP groups.
Parked slave broadcast (PSB)	Control (LMP) PSB-C, User (L2CAP) PSB-U	Parked physical link, basic or adapted physical channel.	Unreliable, uni-directional broadcast to all piconet devices. Used for LMP and L2CAP traffic to parked devices, and for access requests from parked devices.

Table 3.1: Logical transport types.

The classification of each link type follows from a selection procedure within three categories.

3.5.1 Casting

The first category is that of casting. This may be either unicast or broadcast.

- *Unicast links.* Unicast links exist between exactly two endpoints. Traffic may be sent in either direction on unicast links. All unicast links are connection-oriented, meaning that a connection procedure takes place before the link may be used. In the case of the default ACL links, the connection procedure is an implicit step within the general paging procedure used to form ad-hoc piconets.
- *Broadcast links.* Broadcast links exist between one source device and zero or more receiver devices. Traffic is unidirectional, i e only sent from the source devices to the receiver devices. Broadcast links are connectionless, meaning that there is no procedure to create these links, and data may be sent over them at any time. Broadcast links are unreliable, and there is no guarantee that the data will be received.

3.5.2 Scheduling and acknowledgement scheme

The second category relates to the scheduling and acknowledgement scheme of the link, and implies the type of traffic that is supported by the link. These are synchronous, isochronous or asynchronous. There are no specific isochronous links defined, though the default ACL link can be configured to operate in this fashion.

- *Synchronous links.* Synchronous links provide a method of associating the Bluetooth piconet clock with the transported data. This is achieved by reserving regular slots on the physical channel, and transmitting fixed size packets at these regular intervals. Such links are suitable for constant rate isochronous data.
- *Asynchronous links.* Asynchronous links provide a method for transporting data that has no time-based characteristics. The data is normally expected to be retransmitted until successfully received, and each data entity can be processed at any time after receipt, without reference to the time of receipt of any previous or successive entity in the stream (providing the ordering of data entities is preserved.)
- *Isochronous links.* Isochronous links provide a method for transporting data that has time-based characteristics. The data may be retransmitted until received or expired. The data rate on the link need not be constant (this being the main difference from synchronous links.)

3.5.3 Class of data

The final category is related to the class of data that is carried by the link. This is either control (LMP) data or user data. The user data category is sub-divided into L2CAP (or framed) data and stream (or unframed) data.

- *Control links.* Control links are only used for transporting LMP messages between two link managers. These links are invisible above the baseband layer, and cannot be directly instantiated, configured or released by applications, other than by the use of the connection and disconnection services that have this effect implicitly. Control links are always multiplexed with an equivalent L2CAP link onto an ACL logical transport. Subject to the rules

defining the ARQ scheme, the control link traffic always takes priority over the L2CAP link traffic.

- *L2CAP links.* L2CAP links are used to transport L2CAP PDUs, which may carry the L2CAP signaling channel (on the default ACL-U logical link only) or framed user data submitted to user-instantiated L2CAP channels. L2CAP frames submitted to the baseband may be larger than the available baseband packets. A link control protocol embedded within the LLID field preserves the frame-start and frame-continuation semantics when the frame is transmitted in a number of fragments to the receiver.
- *Stream links.* Stream links are used to transport user data that has no inherent framing that should be preserved when delivering the data. Lost data may be replaced by padding at the receiver.

3.5.4 Asynchronous connection-oriented (ACL)

The asynchronous connection-oriented (ACL) logical transport is used to carry LMP and L2CAP control signaling and best effort asynchronous user data. The ACL logical transport uses a simple 1-bit ARQN/SEQN scheme to provide simple channel reliability. Every active slave device within a piconet has one ACL logical transport to the piconet master, known as the default ACL.

The default ACL is created between the master and the slave when a device joins a piconet (connects to the basic piconet physical channel). This default ACL is assigned a logical transport address (LT_ADDR) by the piconet master. This LT_ADDR is also used to identify the active physical link when required (or as a piconet active member identifier, effectively for the same purpose.)

The LT_ADDR for the default ACL is reused for synchronous connection-oriented logical transports between the same master and slave. (This is for reasons of compatibility with earlier Bluetooth specifications.) Thus the LT_ADDR is not sufficient on its own to identify the default ACL. However the packet types used on the ACL are different from those used on the synchronous connection-oriented logical transport. Therefore, the ACL logical transport can be identified by the LT_ADDR field in the packet header in combination with the packet type field.

The default ACL may be used for isochronous data transport by configuring it to automatically flush packets after the packets have expired. Asynchronous traffic can be sent over an ACL logical transport configured for isochronous traffic by marking the asynchronous packets as non-automatically-flushable. This allows both isochronous and asynchronous traffic to be transferred at the same time to a single device.

If the default ACL is removed from the active physical link then all other logical transports that exist between the master and the slave are also removed. In the case of unexpected loss of synchronization to the piconet physical channel the physical link and all logical transports and logical links cease to exist at the time that this synchronization loss is detected.

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A device may remove its default ACL (and by implication its active physical link) but remain synchronized to the piconet. This procedure is known as parking, and a device that is synchronized to the piconet, but has no active physical link is parked within that piconet.

When the device transitions to the parked state the default ACL logical links that are transported on the default ACL logical transport remain in existence, but become suspended. No data may be transferred across a suspended logical link. When the device transitions from the parked state back into active state, a new default ACL logical transport is created (it may have a different LT_ADDR from the previous one) and the suspended logical links are attached to this default ACL and become active once again.

3.5.5 Synchronous connection-oriented (SCO)

The synchronous connection-oriented (SCO) logical transport is a symmetric, point-to-point channel between the master and a specific slave. The SCO logical transport reserves slots on the physical channel and can therefore be considered as a circuit-switched connection between the master and the slave. SCO logical transports carry 64 kb/s of information synchronized with the piconet clock. Typically this information is an encoded voice stream. Three different SCO configurations exist, offering a balance between robustness, delay and bandwidth consumption.

Each SCO-S logical link is supported by a single SCO logical transport, which is assigned the same LT_ADDR as the default ACL logical transport between the devices. Therefore the LT_ADDR field is not sufficient to identify the destination of a received packet. Because the SCO links use reserved slots, a device uses a combination of the LT_ADDR, the slot numbers (a property of the physical channel) and the packet type to identify transmissions on the SCO link.

Although slots are reserved for the SCO, it is permissible to use a reserved slot for traffic from another channel that has a higher priority. This may be required as a result of QoS commitments, or to send LMP signaling on the default ACL when the physical channel bandwidth is fully occupied by SCOs. As SCOs carry different packet types to ACLs, the packet type is used to identify SCO traffic (in addition to the slot number and LT_ADDR.)

There are no further architectural layers defined by the Bluetooth core specification that are transported over an SCO link. A number of standard formats are defined for the 64 kb/s stream that is transported, or an unformatted stream is allowed where the application is responsible for interpreting the encoding of the stream.

3.5.6 Extended synchronous connection-oriented (eSCO)

The extended synchronous connection-oriented (eSCO) logical transport is a symmetric or asymmetric, point-to-point link between the master and a specific

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slave. The eSCO reserves slots on the physical channel and can therefore be considered as a circuit-switched connection between the master and the slave. eSCO links offer a number of extensions over the standard SCO links, in that they support a more flexible combination of packet types and selectable data contents in the packets and selectable slot periods, allowing a range of synchronous bit rates to be supported.

eSCO links also can offer limited retransmission of packets (unlike SCO links where there is no retransmission.) If these retransmissions are required they take place in the slots that follow the reserved slots, otherwise the slots may be used for other traffic.

Each eSCO-S logical link is supported by a single eSCO logical transport, identified by a LT_ADDR that is unique within the piconet for the duration of the eSCO. eSCO-S links are created using LM signaling and follow scheduling rules similar to SCO-S links.

There are no further architectural layers defined by the Bluetooth core specification that are transported over an eSCO-S link. Instead applications may use the data stream for whatever purpose they require, subject to the transport characteristics of the stream being suitable for the data being transported.

3.5.7 Active slave broadcast (ASB)

The active slave broadcast logical transport is used to transport L2CAP user traffic to all devices in the piconet that are currently connected to the physical channel that is used by the ASB. There is no acknowledgement protocol and the traffic is uni-directional from the piconet master to the slaves. The ASB channel may be used for L2CAP group traffic (a legacy of the 1.1 specification), and is never used for L2CAP connection-oriented channels, L2CAP control signaling or LMP control signaling.

The ASB logical transport is inherently unreliable because of the lack of acknowledgement. To improve the reliability, each packet is transmitted a number of times. An identical sequence number is used to assist with filtering retransmissions at the slave device.

The ASB logical transport is identified by a reserved LT_ADDR. (The reserved LT_ADDR address is also used by the PSB logical transport.) An active slave will receive traffic on both logical transports, and cannot readily distinguish between them. As the ASB logical transport does not carry LMP traffic an active slave can ignore packets received over the LMP logical link on the ASB logical transport. However L2CAP traffic transmitted over the PSB logical transport is also received by active slaves on the ASB logical transport and cannot be distinguished from L2CAP traffic sent on the ASB transport.

An ASB is implicitly created whenever a piconet exists, and there is always one ASB associated with each of the basic and adapted piconet physical channels that exist within the piconet. Because the basic and adapted piconet physical

channels are mostly coincident a slave device cannot distinguish which of the ASB channels is being used to transmit the packets. This adds to the general unreliability of the ASB channel. (Although it is, perhaps, no more unreliable than general missed packets.)

A master device may decide to use only one of its two possible ASBs (when it has both a basic and adapted piconet physical channel), as with sufficient retransmissions it is possible to address both groups of slaves on the same ASB channel.

The ASB channel is never used to carry LMP or L2CAP control signals.

3.5.8 Parked slave broadcast (PSB)

The parked slave broadcast logical transport is used for communications between the master and slaves that are parked (have given up their default ACL logical transport.) The parked slave broadcast link is the only logical transport that exists between the piconet master and parked slaves.

The PSB logical transport is more complex than the other logical transports as it consists of a number of phases, each having a different purpose. These phases are the control information phase (used to carry the LMP logical link), the user information phase (used to carry the L2CAP logical link), and the access phase (carrying baseband signaling.) The control information and broadcast information phases are usually mutually exclusive as only one of them can be supported in a single beacon interval. (Even if there is no control or user information phase, the master is still required to transmit a packet in the beacon slots so that the parked slaves can resynchronize.) The access phase is normally present unless cancelled in a control information message.

The control information phase is used for the master to send information to the parked slaves containing modifications to the PSB transport attributes, modifications to the beacon train attributes, or a request for a parked slave to become active in the piconet (known as unparking). This control information is carried in LMP messages on the LMP logical link. (The control information phase is also present in the case of a user information phase where the user information requires more than one baseband packet.)

Packets in the control information phase are always transmitted in the physical channel beacon train slots, and cannot be transmitted on any other slots. The control information occupies a single DM1 packet and is repeated in every beacon train slot within a single beacon interval. (If there is no control information then there may be a user information phase that uses the beacon slots. If neither phase is used then the beacon slots are used for other logical transport traffic or for NULL packets.)

The user information phase is used for the master to send L2CAP packets that are destined for all piconet slaves. User information may occupy one or more

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baseband packets. If the user information occupies a single packet then the user information packet is repeated in each of the piconet channel beacon train slots.

If the user information occupies more than one baseband packet then it is transmitted in slots after the beacon train (the broadcast scan window) and the beacon slots are used to transmit a control information phase message that contains the timing attributes of this broadcast scan window. This is required so that the parked slaves remain connected to the piconet physical channel to receive the user information.

The access phase is normally present unless temporarily cancelled by a control message carried in the control information broadcast phase. The access window consists of a sequence of slots that follow the beacon train. In order for a parked slave to become active in the piconet, it may send such an access request to the piconet master during the access window. Each parked slave is allocated an access request address (not necessarily unique) that controls when during the access window the slave requests access.

The PSB logical transport is identified by the reserved LT_ADDR of 0. This reserved LT_ADDR address is also used by the ASB logical transport. Parked slaves are not normally confused by the duplicated use of the LT_ADDR as they are only connected to the piconet physical channel during the time that the PSB transport is being used.

3.5.9 Logical links

Some logical transports are capable of supporting different logical links, either concurrently multiplexed, or one of the choice. Within such logical transports, the logical link is identified by the logical link identifier (LLID) bits in the payload header of baseband packets that carry a data payload. The logical links distinguish between a limited set of core protocols that are able to transmit and receive data on the logical transports. Not all of the logical transports are able to carry all of the logical links (the supported mapping is shown in Figure 3.2 on page 27.) In particular the SCO and eSCO logical transports are only able to carry constant data rate streams, and these are uniquely identified by the LT_ADDR. Such logical transports only use packets that do not contain a payload header, as their length is known in advance, and no LLID is necessary.

3.5.10 ACL Control Logical Link (ACL-C)

The ACL Control Logical Link (ACL-C) is used to carry LMP signaling between devices in the piconet. The control link is only carried on the default ACL logical transport and on the PSB logical transport (in the control information phase). The ACL-C link is always given priority over the ACL-U (see below) link when carried on the same logical transport. The ACL-C link is always given priority over the ACL-U (see below) link when carried on the same logical transport.

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3.5.10 User Asynchronous/Isochronous Logical Link (ACL-U)

The user asynchronous/isochronous logical link (ACL-U) is used to carry all asynchronous and isochronous framed user data. The ACL-U link is carried on all but the synchronous logical transports. Packets on the ACL-U link are identified by one of two reserved LLID values. One value is used to indicate whether the baseband packet contains the start of an L2CAP frame and the other indicates a continuation of a previous frame. This ensures correct synchronization of the L2CAP reassembler following flushed packets. The use of this technique removes the need for a more complex L2CAP header in every baseband packet (the header is only required in the L2CAP start packets), but adds the requirement that a complete L2CAP frame shall be transmitted before a new one is transmitted. (An exception to this rule being the ability to flush a partially transmitted L2CAP frame in favor of another L2CAP frame.)

3.5.11 User Synchronous/Extended Synchronous Logical Links (SCO-S/ eSCO-S)

Synchronous (SCO-S) and extended synchronous (eSCO-S) logical links are used to support isochronous data delivered in a stream without framing. These links are associated with a single logical transport, where data is delivered in constant sized units at a constant rate. There is no LLID within the packets on these transports, as only a single logical link can be supported, and the packet length and scheduling period are pre-defined and remain fixed during the lifetime of the link.

Variable rate isochronous data cannot be carried by the SCO-S or eSCO-S logical links. In this case the data must be carried on ACL-U logical links, which use packets with a payload header.

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3.6 L2CAP CHANNELS

L2CAP provides a multiplexing role allowing many different applications to share the resources of an ACL-U logical link between two devices. Applications and service protocols interface with L2CAP using a channel-oriented interface to create connections to equivalent entities on other devices.

L2CAP channel endpoints are identified to their clients by a Channel Identifier (CID). This is assigned by L2CAP, and each L2CAP channel endpoint on any device has a different CID.

L2CAP channels may be configured to provide an appropriate Quality of Service (QoS) to the application. L2CAP maps the channel onto the ACL-U logical link.

L2CAP supports channels that are connection-oriented and others that are group-oriented. Group-oriented channels may be mapped onto the ASB-U logical link, or implemented as iterated transmission to each member in turn over an ACL-U logical link.

Apart from the creation, configuration and dismantling of channels, the main role of L2CAP is to multiplex service data units (SDUs) from the channel clients onto the ACL-U logical links, and to carry out a simple level of scheduling, selecting SDUs according to relative priority.

L2CAP can provide a per channel flow control with the peer L2CAP layer. This option is selected by the application when the channel is established. L2CAP can also provide enhanced error detection and retransmission to (a) reduce the probability of undetected errors being passed to the application and (b) recover from loss of portions of the user data when the Baseband layer performs a flush on the ACL-U logical link.

In the case where an HCI is present, the L2CAP is also required to segment L2CAP SDUs into fragments that will fit into the baseband buffers, and also to operate a token based flow control procedure over the HCI, submitting fragments to the baseband only when allowed to do so. This may affect the scheduling algorithm.

4 COMMUNICATION TOPOLOGY

4.1 PICONET TOPOLOGY

Any time a Bluetooth link is formed it is within the context of a piconet. A piconet consists of two or more devices that occupy the same physical channel (which means that they are synchronized to a common clock and hopping sequence.) The common (piconet) clock is identical to the Bluetooth clock of one of the devices in the piconet, known as the master of the piconet, and the hopping sequence is derived from the master's clock and the master's Bluetooth device address. All other synchronized devices are referred to as slaves in the piconet. The terms master and slave are only used when describing these roles in a piconet.

Within a common location a number of independent piconets may exist. Each piconet has a different physical channel (that is a different master device and an independent piconet clock and hopping sequence.)

A Bluetooth device may participate concurrently in two or more piconets. It does this on a time-division multiplexing basis. A Bluetooth device can never be a master of more than one piconet. (Since the piconet is defined by synchronization to the master's Bluetooth clock it is impossible to be the master of two or more piconets.) A Bluetooth device may be a slave in many independent piconets.

A Bluetooth device that is a member of two or more piconets is said to be involved in a scatternet. Involvement in a scatternet does not necessarily imply any network routing capability or function in the Bluetooth device. The Bluetooth core protocols do not, and are not intended to offer such functionality, which is the responsibility of higher level protocols and is outside the scope of the Bluetooth core specification.

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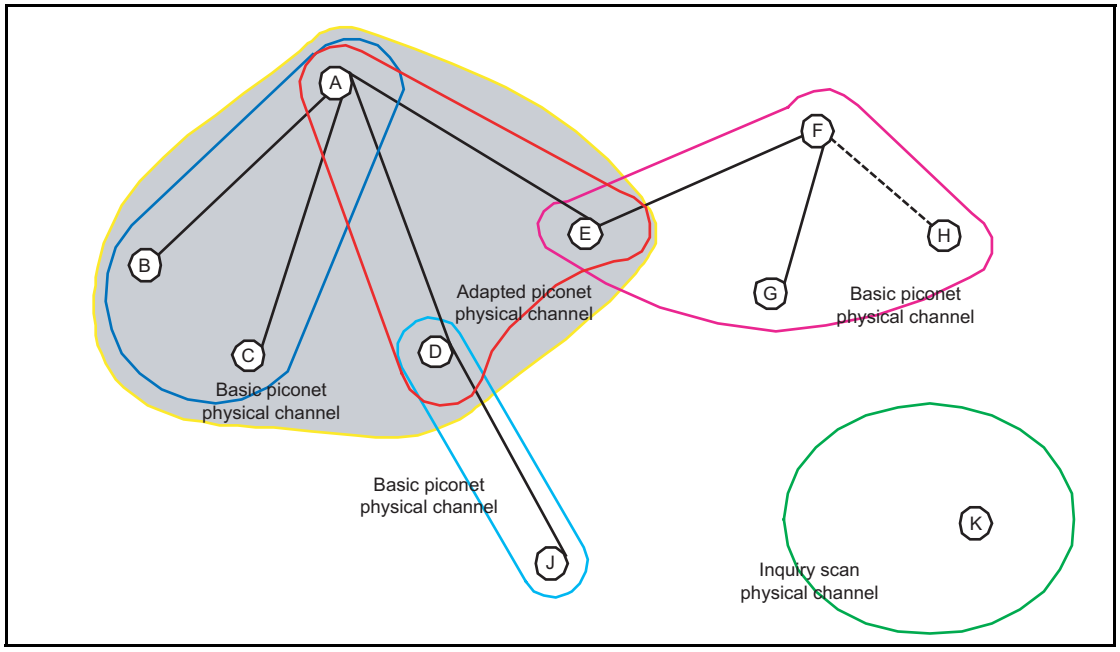


Figure 4.1: Example Bluetooth topology

In Figure 4.1 on page 51 an example topology is shown that demonstrates a number of the architectural features described below. Device A is a master in a piconet (represented by the shaded area, and known as piconet A) with devices B, C, D and E as slaves. Two other piconets are shown: a) one piconet with device F as master (known as piconet F) and devices E, G and H as slaves and b) one piconet with device D as master (known as piconet D) and device J as slave.

In piconet A there are two physical channels. Devices B and C are using the basic piconet physical channel (represented by the blue enclosure) as they do not support adaptive frequency hopping. Devices D and E are capable of supporting adaptive frequency hopping, and are using the adapted piconet physical channel (represented by the red enclosure.) Device A is capable of adaptive frequency hopping, and operates in a TDM basis on both physical channels according to which slave is being addressed.

Piconet D and piconet F are both using only a basic piconet physical channel (represented by the cyan and magenta enclosures respectively.) In the case of piconet D this is because device J does not support the adaptive hopping mode. Although device D supports adaptive hopping it cannot use it in this piconet. In piconet F device F does not support adaptive hopping, and therefore it cannot be used in this piconet.

Device K is shown in the same locality as the other devices. It is not currently a member of a piconet, but has services that it offers to other Bluetooth devices. It is currently listening on its inquiry scan physical channel (represented by the green enclosure), awaiting an inquiry request from another device.

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4.2 OPERATIONAL PROCEDURES AND MODES

The typical operational mode of a Bluetooth device is to be connected to other Bluetooth devices (in a piconet) and exchanging data with that Bluetooth device. As Bluetooth is an ad-hoc wireless communications technology there are also a number of operational procedures that enable piconets to be formed so that the subsequent communications can take place. Procedures and modes are applied at different layers in the architecture and therefore a device may be engaged in a number of these procedures and modes concurrently.

4.2.1 Inquiry (Discovering) Procedure

Bluetooth devices use the inquiry procedure to discover nearby devices, or to be discovered by devices in their locality.

The inquiry procedure is asymmetrical. A Bluetooth device that tries to find other nearby devices is known as an inquiring device and actively sends inquiry requests. Bluetooth devices that are available to be found are known as discoverable devices and listen for these inquiry requests and send responses. The inquiry procedure uses a special physical channel for the inquiry requests and responses.

Both inquiring and discoverable devices may already be connected to other Bluetooth devices in a piconet. Any time spent inquiring or occupying the inquiry scan physical channel needs to be balanced with the demands of the QoS commitments on existing logical transports.

The inquiry procedure does not make use of any of the architectural layers above the physical channel, although a transient physical link may be considered to be present during the exchange of inquiry and inquiry response information.

4.2.1.1 Extended Inquiry Response

An Extended Inquiry Response can be used to provide miscellaneous information during the inquiry response procedure. Data types are defined for such things as local name and supported services, information that otherwise would have to be obtained during a connection. A device that receives a local name and a list of supported services in an extended inquiry response does not have to connect to do a remote name request and an SDP service search, thereby shortening the time to useful information. It is recommended that a device includes all supported services and a significant minority part of its local name in the extended inquiry response.

The extended inquiry response procedure is backwards compatible with the standard inquiry response procedure.

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4.2.2 Paging (Connecting) Procedure

The procedure for forming connections is asymmetrical and requires that one Bluetooth device carries out the page (connection) procedure while the other Bluetooth device is connectable (page scanning.) The procedure is targeted, so that the page procedure is only responded to by one specified Bluetooth device.

The connectable device uses a special physical channel to listen for connection request packets from the paging (connecting) device. This physical channel has attributes that are specific to the connectable device, hence only a paging device with knowledge of the connectable device is able to communicate on this channel.

Both paging and connectable devices may already be connected to other Bluetooth devices in a piconet. Any time spent paging or occupying the page scan physical channel needs to be balanced with the demands of the QoS commitments on existing logical transports.

4.2.3 Connected mode

After a successful connection procedure, the devices are physically connected to each other within a piconet. This means that there is a piconet physical channel to which they are both connected, there is a physical link between the devices, and there are default ACL-C and ACL-U logical links. When in the connected mode it is possible to create and release additional logical links, and to change the modes of the physical and logical links while remaining connected to the piconet physical channel. It is also possible for the device to carry out inquiry, paging or scanning procedures or to be connected to other piconets without needing to disconnect from the original piconet physical channel.

Additional logical links are created using the Link Manager that exchanges Link Manager Protocol messages with the remote Bluetooth device to negotiate the creation and settings for these links. Default ACL-C and ACL-U logical links are always created during the connection process, and these are used for LMP messages and the L2CAP signaling channel respectively.

It is noted that two default logical links are created when two units are initially connected. One of these links (ACL-C) transports the LMP control protocol and is invisible to the layers above the Link Manager. The other link (ACL-U) transports the L2CAP signaling protocol and any multiplexed L2CAP best-effort channels. It is common to refer to a default ACL logical transport, which can be resolved by context, but typically refers to the default ACL-U logical link. Also note that these two logical links share a logical transport.

During the time that a slave device is actively connected to a piconet there is always a default ACL logical transport between the slave and the master device. There are two methods of deleting the default ACL logical transport. The first method is to detach the device from the piconet physical channel, at

which time the entire hierarchy of L2CAP channels, logical links, and logical transports between the devices is deleted.

The second method is to place the physical link to the slave device in the park state, at which time it gives up its default ACL logical transport. This is only allowed if all other logical transports have been deleted (except for the ASB logical transport that cannot be explicitly created or deleted.) It is not allowed to park a device while it has any logical transports other than the default ACL and ASB logical transports.

When the slave device physical link is parked, its default ACL logical transport is released and the ASB logical transport is replaced with a PSB logical transport. The ACL-C and ACL-U logical links that are multiplexed onto the default ACL logical transport remain in existence but cannot be used for the transport of data. The Link Manager on the master device restricts itself to the use of LMP messages that are allowed to be transported over the PSB-C logical link. The Channel Manager and L2CAP Resource Manager ensure that no L2CAP unicast data traffic is submitted to the controller while the device is parked. The Channel Manager may decide to manage the parking and unparking of the device as necessary to allow data to be transported.

4.2.4 Hold mode

Hold mode is not a general device mode, but applies to unreserved slots on the physical link. When in this mode, the physical link is only active during slots that are reserved for the operation of the synchronous link types SCO and eSCO. All asynchronous links are inactive. Hold modes operate once for each invocation and are then exited when complete, returning to the previous mode.

4.2.5 Sniff mode

Sniff mode is not a general device mode, but applies to the default ACL logical transports. When in this mode the availability of these logical transports is modified by defining a duty cycle consisting of periods of presence and absence. Devices that have their default ACL logical transports in sniff mode may use the absent periods to engage in activity on another physical channel, or to enter reduced power mode. Sniff mode only affects the default ACL logical transports (i.e. their shared ACL logical transport), and does not apply to any additional SCO or eSCO logical transports that may be active. The periods of presence and absence of the physical link on the piconet physical channel is derived as a union of all logical transports that are built on the physical link.

Sniff subrating provides a mechanism for further reducing the active duty cycle, thereby enhancing the power-saving capability of sniff mode. Sniff subrating allows a Host to create a guaranteed access-like connection by specifying maximum transmit and receive latencies. This allows the basebands to optimize the low power performance without having to exit and re-enter sniff mode using Link Manager commands.

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Note that broadcast logical transports have no defined expectations for presence or absence. A master device should aim to schedule broadcasts to coincide with periods of physical link presence within the piconet physical channel, but this may not always be possible or practical. Repetition of broadcasts is defined to improve the possibilities for reaching multiple slaves without overlapping presence periods. However, broadcast logical transports cannot be considered to be reliable.

4.2.6 Parked state

A slave device may remain connected to a piconet but have its physical link in the parked state. In this state the device cannot support any logical links to the master with the exception of the PSB-C and PSB-U logical links that are used for all communication between the piconet master and the parked slave.

When the physical link to a slave device is parked this means that there are restrictions on when the master and slave may communicate, defined by the PSB logical transport parameters. During times when the PSB logical transport is inactive (or absent) then the devices may engage in activity on other physical channels, or enter reduced power mode.

4.2.7 Role switch procedure

The role switch procedure is a method for swapping the roles of two devices connected in a piconet. The procedure involves moving from the physical channel that is defined by the original master device to the physical channel that is defined by the new master device. In the process of swapping from one physical channel to the next, the hierarchy of physical links and logical transports are removed and rebuilt, with the exception of the ASB and PSB logical transports that are implied by the topology and are not preserved. After the role switch, the original piconet physical channel may cease to exist or may be continued if the original master had other slaves that are still connected to the channel.

The procedure only copies the default ACL logical links and supporting layers to the new physical channel. Any additional logical transports are not copied by this procedure, and if required this must be carried out by higher layers. The LT_ADDRs of any affected transports may not be preserved as the values may already be in use on the new physical channel.

If there are any QoS commitments or modes such as sniff mode on the original logical transports, then these are not preserved after a role switch. These must be renegotiated after the role switch has completed.

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4.2.8 Enhanced Data Rate

Enhanced Data Rate is a method of extending the capacity and types of Bluetooth packets for the purposes of increasing the maximum throughput, providing better support for multiple connections, and lowering power consumption, while the remainder of the architecture is unchanged.

Enhanced Data Rate may be selected as a mode that operates independently on each logical transport. Once enabled, the packet type bits in the packet header are interpreted differently from their meaning in Basic Rate mode. This different interpretation is clarified in conjunction with the logical transport address field in the header. The result of this interpretation allows the packet payload header and payload to be received and demodulated according to the packet type. Enhanced Data Rate can be enabled only for ACL-U, eSCO-S logical transports and cannot be enabled for ACL-C, SCO-S, and the broadcast logical transports.

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5 SECURE SIMPLE PAIRING OVERVIEW

The primary goal of Secure Simple Pairing is to simplify the pairing procedure for the user. Secondary goals are to maintain or improve the security in Bluetooth wireless technology. Since high levels of security and ease-of-use are often at opposite ends of the spectrum in many technologies and products, much care has been taken to maximize security while minimizing complexity from the end user's point of view.

5.1 SECURITY GOALS

Secure Simple Pairing has two security goals: protection against passive eavesdropping and protection against man-in-the-middle (MITM) attacks (active eavesdropping). It is a goal of Secure Simple Pairing to exceed the maximum security level provided by the use of a 16 alphanumeric PIN with the pairing algorithm used in Bluetooth Core Specification version 2.0 + EDR and earlier versions. Note that many Bluetooth devices compliant with Bluetooth Core Specification 2.0 + EDR and earlier versions use a 4-digit PIN or a fixed PIN of commonly known values significantly limiting the security on the link.

5.2 PASSIVE EAVESDROPPING PROTECTION

A strong link key coupled with a strong encryption algorithm is necessary to give the user protection against passive eavesdropping. The strength of the link key is based on the amount of entropy (or randomness) in its generation process which would not be known by an attacker. Using legacy pairing, the only source of entropy is the PIN which, in many use cases, is typically four digits either selected by the user or fixed for a given product. Therefore, if the pairing procedure and one authentication exchange is recorded one can run an exhaustive search to find the PIN in a very short amount of time on commonly available computing hardware. With Secure Simple Pairing, the recording attack becomes much harder as the attacker must have solved a hard problem in public key cryptography in order to derive the link key from the recorded information. This protection is independent of the length of the passkey or other numeric values that the user must handle. Secure Simple Pairing gives the same resistance against the recording and passive eavesdropping attacks even when the user is not required to do anything.

Secure Simple Pairing uses Elliptic Curve Diffie Hellman (ECDH) public key cryptography as a means to thwart passive eavesdropping attacks. ECDH provides a very high degree of strength against passive eavesdropping attacks but it may be subject to MITM attacks, which however, are much harder to perform in practice than the passive eavesdropping attack (see Section 5.3, “Man-In-The-Middle Protection,” on page 58).

Using the security protocols in the Bluetooth Core Specification version 2.0 + EDR and earlier with a 16 numeric digit PIN achieves about 53 bits of entropy

whereas a 16 character alphanumeric, case sensitive PIN yields about 95 bits of entropy when the entire 62 character set is used ([0, … 9, 'A', … 'Z', 'a', … 'z']). Secure Simple Pairing has approximately 95 bits of entropy using the FIPS approved P192 elliptic curve which is at least as good as the entropy in Bluetooth Core Specification 2.0 + EDR and earlier using a 16 character, alphanumeric, case sensitive PIN. Secure Simple Pairing, therefore, exceeds the security requirements of the Bluetooth SIM Access Profile (SAP) which is the profile with the most stringent security requirements. ECDH cryptography was selected over standard Diffie Hellman (often referred to as DH76) since it is computationally less complex and less likely to exceed the low computational capacity in common Bluetooth Controllers.

5.3 MAN-IN-THE-MIDDLE PROTECTION

A man-in-the-middle (MITM) attack occurs when a user wants to connect two devices but instead of connecting directly with each other they unknowingly connect to a third (attacking) device that plays the role of the device they are attempting to pair with. The third device then relays information between the two devices giving the illusion that they are directly connected. The attacking device may even eavesdrop on communication between the two devices (known as active eavesdropping) and is able to insert and modify information on the connection. In this type of attack, all of the information exchanged between the two devices are compromised and the attacker may inject commands and information into each of the devices thus potentially damaging the function of the devices. Devices falling victim to the attack are capable of communicating only when the attacker is present. If the attacker is not active or out range, the two victim devices will not be able to communicate directly with each other and the user will notice it.

To prevent MITM attacks, Secure Simple Pairing offers two user assisted numeric methods: numerical comparison or passkey entry. If Secure Simple Pairing would use 16 decimal digit numbers, then the usability would be the same as using legacy pairing with 16 decimal digit PIN. The chance for a MITM to succeed inserting its own link keys in this case is a 1 in $10^{16} = 2^{53}$ pairing instances, which is an unnecessarily low probability.

Secure Simple Pairing protects the user from MITM attacks with a goal of offering a 1 in 1,000,000 chance that a MITM could mount a successful attack. The strength of the MITM protections was selected to minimize the user impact by using a six digit number for numerical comparison and Passkey entry. This level of MITM protection was selected since, in most cases, users can be alerted to the potential presence of a MITM attacker when the connection process fails as a result of a failed MITM attack. While most users feel that provided that they have not compromised their passkey, a 4-digit key is sufficient for authentication (i.e. bank card PIN codes), the use of six digits allows Secure Simple Pairing to be FIPS compliant and this was deemed to have little perceivable usability impact.

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5.4 ASSOCIATION MODELS

Secure Simple Pairing uses four association models referred to as Numeric Comparison, Just Works, Out Of Band, and Passkey Entry. Each of these association models are described in more detail in the following sections.

The association model used is deterministic based on the I/O capabilities of the two devices.

5.4.1 Numeric Comparison

The Numeric Comparison association model is designed for scenarios where both devices are capable of displaying a six digit number and both are capable of having the user enter "yes" or "no". A good example of this model is the cell phone / PC scenario.

The user is shown a six digit number (from "000000" to "999999") on both displays and then asked whether the numbers are the same on both devices. If "yes" is entered on both devices, the pairing is successful.

The numeric comparison serves two purposes. First, since many devices do not have unique names, it provides confirmation to the user that the correct devices are connected with each other. Second, the numeric comparison provides protection against MITM attacks (see Section 5.3 on page 58).

Note that there is a significant difference from a cryptographic point of view between Numeric Comparison and the PIN entry model used by Bluetooth Core Specification and earlier versions. In the Numeric Comparison association model, the six digit number is an artifact of the security algorithm and not an input to it, as is the case in the Bluetooth security model. Knowing the displayed number is of no benefit in decrypting the encoded data exchanged between the two devices.

5.4.2 Just Works

The Just Works association model is primarily designed for scenarios where at least one of the devices does not have a display capable of displaying a six digit number nor does it have a keyboard capable of entering six decimal digits. A good example of this model is the cell phone/mono headset scenario where most headsets do not have a display.

The Just Works association model uses the Numeric Comparison protocol but the user is never shown a number and the application may simply ask the user to accept the connection (exact implementation is up to the end product manufacturer).

The Just Works association model provides the same protection as the Numeric Comparison association model against passive eavesdropping but offers no protection against the MITM attack.

When compared against today's experience of a headset with a fixed PIN, the security level of the Just Works association model is considerably higher since a high degree of protection against passive eavesdropping is realized.

5.4.3 Out of Band

The Out of Band (OOB) association model is primarily designed for scenarios where an Out of Band mechanism is used to both discover the devices as well as to exchange or transfer cryptographic numbers used in the pairing process. In order to be effective from a security point of view, the Out of Band channel should provide different properties in terms of security compared to the Bluetooth radio channel. The Out of Band channel should be resistant to MITM attacks. If it is not, security may be compromised during authentication.

The user's experience differs a bit depending on the Out of Band mechanism. As an example, with a Near Field Communication (NFC) solution, the user(s) will initially touch the two devices together, and is given the option to pair the first device with the other device. If "yes" is entered, the pairing is successful. This is a single touch experience where the exchanged information is used in both devices. The information exchanged includes discovery information (such as the Bluetooth Device Address) as well as cryptographic information. One of the devices will use a Bluetooth Device Address to establish a connection with the other device. The rest of the exchanged information is used during authentication.

The OOB mechanism may be implemented as either read only or read/write. If one side is read only, a one-way authentication is performed. If both sides are read/write, a two-way authentication is performed.

The OOB protocol is selected only when the pairing process has been activated by previous OOB exchange of information and one (or both) of the device(s) gives OOB as the IO capabilities. The protocol uses the information which has been exchanged and simply asks the user to confirm connection.

The OOB association model supports any OOB mechanism where cryptographic information and the Bluetooth Device Address can be exchanged. The OOB association model does not support a solution where the user has activated a Bluetooth connection and would like to use OOB for authentication only.

5.4.4 Passkey Entry

The Passkey Entry association model is primarily designed for the scenario where one device has input capability but does not have the capability to display six digits and the other device has output capabilities. A good example of this model is the PC and keyboard scenario.

The user is shown a six digit number (from "000000" to "999999") on the device with a display, and is then asked to enter the number on the other

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device. If the value entered on the second device is correct, the pairing is successful. Note that there is a significant difference from a cryptographic point of view between Passkey Entry and the PIN entry model used by Bluetooth Core Specification 2.0 + EDR and earlier versions. In the Passkey Entry association model, the six digit number is independent of the security algorithm and not an input to it, as is the case in the 2.0 + EDR security model. Knowing the entered number is of no benefit in decrypting the encoded data exchanged between the two devices.

5.4.5 Association Model Overview

The following diagram shows Secure Simple Pairing from the point of view of the technology used for discovery and then the different association possibilities.

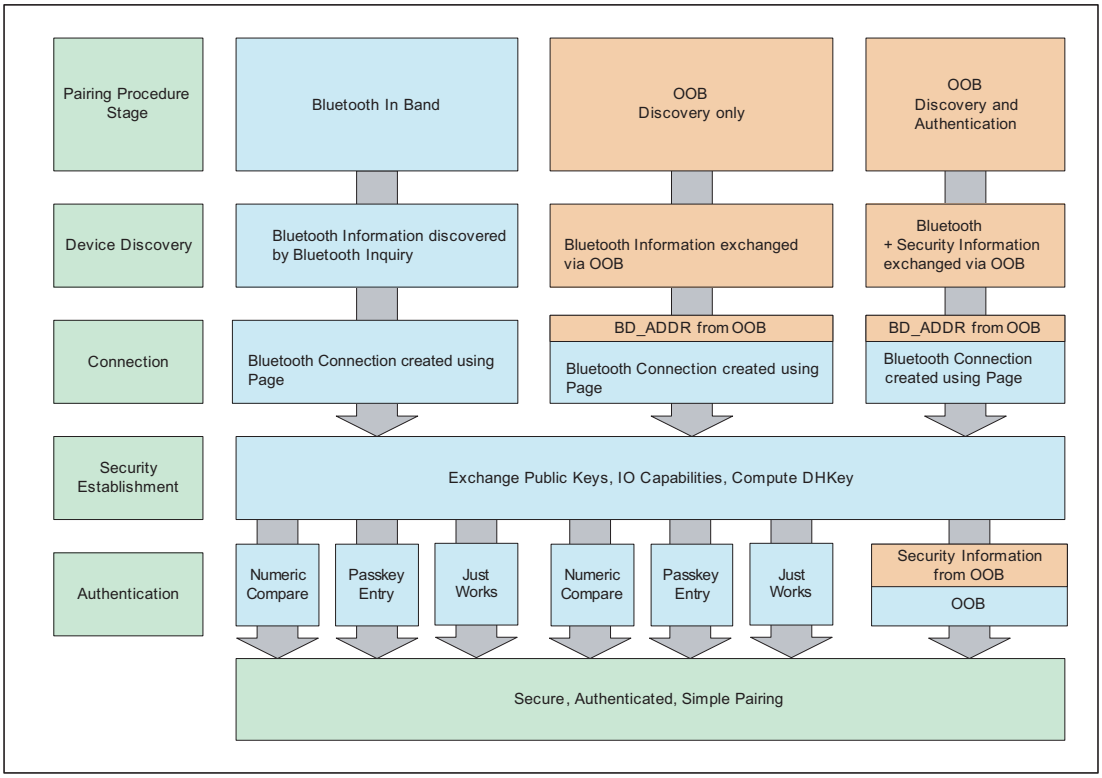


Figure 5.1: Secure Simple Pairing Association Models

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ACRONYMS & ABBREVIATIONS

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1 LIST OF ACRONYMS AND ABBREVIATIONS

Acronym or abbreviation	Writing out in full	Which means
8DPSK	8 phase Differential Phase Shift Keying	3 Mbps modulation type used by Enhanced Data rate
π/4 DQPSK	pi/4 Rotated Differential Quaternary Phase Shift Keying	2 Mbps modulation type used by Enhanced Data Rate
A		
ACK	Acknowledge	
ACL	Asynchronous Connection-oriented [logical transport]	Reliable or time-bounded, bi-directional, point-to-point.
ACL-C	ACL Control [logical link] (LMP)	
ACL-U	ACL User [logical link] (L2CAP)	
ACO	Authenticated Ciphering Offset	
AFH	Adaptive Frequency Hopping	
AHS	Adapted Hop Sequence	
AR_ADDR	Access Request Address	
ARQ	Automatic Repeat Request	
ASB	Active Slave Broadcast [logical transport]	Unreliable, uni-directional broadcast to any devices synchronized with the physical channel.
ASB-U	ASB User [logical link] (L2CAP)	
B		
BB	Baseband	
BCH	Bose, Chaudhuri & Hocquenghem	Type of code The persons who discovered these codes in 1959 (H) and 1960 (B&C)
BD_ADDR	Bluetooth Device Address	
BER	Bit Error Rate	
BT	Bandwidth Time	
C		
CAC	Channel Access Code	
CL	Connectionless	
CODEC	COder DECoder	

Table 1.1: Acronyms and Abbreviations

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Acronym or abbreviation	Writing out in full	Which means
COF	Ciphering Offset	
CRC	Cyclic Redundancy Check	
CVSD	Continuous Variable Slope Delta Modulation	
D		
DAC	Device Access Code	
DCE	Data Communication Equipment	
DCE	Data Circuit-Terminating Equipment	In serial communications, DCE refers to a device between the communication endpoints whose sole task is to facilitate the communications process; typically a modem
DCI	Default Check Initialization	
DEVM	Differential Error Vector Magnitude	Measure of modulation error used for Enhanced Data Rate transmitter testing
DH	Data-High Rate	Data packet type for high rate data
DIAC	Dedicated Inquiry Access Code	
DM	Data - Medium Rate	Data packet type for medium rate data
DPSK	Differential Phase Shift Keying	Generic description of Enhanced Data Rate modulation
DQPSK	Differential Quaternary Phase Shift Keying	Modulation type used by Enhanced Data Rate
DTE	Data Terminal Equipment	In serial communications, DTE refers to a device at the endpoint of the communications path; typically a computer or terminal.
DUT	Device Under Test	
DV	Data Voice	Data packet type for data and voice
E		
ED	Erroneous Data	
EDR	Enhanced Data Rate	
EIR	Extended Inquiry Response	Host supplied information transmitted in the inquiry responses substate.

Table 1.1: Acronyms and Abbreviations

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Acronym or abbreviation	Writing out in full	Which means
EIRP	Effective Isotropic Radiated Power	Equivalent power that an isotropic antenna must transmit to provide the same field power density
EPR	Encryption Pause Resume	Mode allowing the device to initiate the pause and resume encryption sequence
eSCO	Extended Synchronous Connection Oriented [logical transport]	Bi-directional, symmetric or asymmetric, point-to-point, general regular data, limited retransmission.
eSCO-S	Stream eSCO (unframed)	used to support isochronous data delivered in a stream without framing
ETSI	European Telecommunications Standards Institute	
F		
FCC	Federal Communications Commission	
FEC	Forward Error Correction code	
FH	Frequency Hopping	
FHS	Frequency Hop Synchronization	
FIFO	First In First Out	
FM	Frequency Modulation	Modulation Type
G		
GFSK	Gaussian Frequency Shift Keying	
GIAC	General Inquiry Access Code	
H		
HCI	Host Controller Interface	
HEC	Header-Error-Check	
HID	Human Interface Device	
HV	High quality Voice	e.g. HV1 packet
HW	Hardware	
I		
IAC	Inquiry Access Code	

Table 1.1: Acronyms and Abbreviations

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Acronym or abbreviation	Writing out in full	Which means
IEEE	Institute of Electronic and Electrical Engineering	
IETF	Internet Engineering Task Force	
IP	Internet Protocol	
IrDA	Infra-red Data Association	
ISM	Industrial, Scientific, Medical	
IUT	Implementation Under Test	
L		
L2CAP	Logical Link Control and Adaptation Protocol	
LAP	Lower Address Part	
LC	Link Controller	Link Controller (or baseband) part of the Bluetooth protocol stack. Low level Baseband protocol handler
LC	Link Control [logical link]	The control logical links LC and ACL-C are used at the link control level and link manager level, respectively.
LCP	Link Control Protocol	
LFSR	Linear Feedback Shift Register	
LLID	Logical Link Identifier	
LM	Link Manager	
LMP	Link Manager Protocol	For LM peer to peer communication
LSB	Least Significant Bit	
LSTO	Link Supervision Timeout Event	Controller can send LSTO event to Host
LT_ADDR	Logical Transport ADDRess	
M		
M	Master or Mandatory	
MAC	Medium Access Control	
Mbps	Million (Mega) bits per second	
MITM	Man-in-the-middle	
MMI	Man Machine Interface	

Table 1.1: Acronyms and Abbreviations

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Acronym or abbreviation	Writing out in full	Which means
MS	Mobile Station	
MSB	Most Significant Bit	
MSC	Message Sequence Chart	
MTU	Maximum Transmission Unit	
N		
NAK	Negative Acknowledge	
NAP	Non-significant Address Part	
O		
O	Optional	
OBEX	OBject EXchange protocol	
OCF	OpCode Command Field	
OGF	OpCode Group Field	
P		
PBF	Packet Boundary Flag	The device supports the capability to correctly handle HCI ACL Data Packets
PCM	Pulse Coded Modulation	
PDU	Protocol Data Unit	a message
PIN	Personal Identification Number	
PM_ADDR	Parked Member Address	
PN	Pseudo-random Noise	
PPM	Part Per Million	
PPP	Point-to-Point Protocol	
PRBS	Pseudo Random Bit Sequence	
PRNG	Pseudo Random Noise Generation	
PSB	Parked Slave Broadcast [logical transport]	Unreliable, uni-directional broadcast to all piconet devices.
PSB-C	PSB Control [logical link] (LMP)	
PSB-U	PSB User [logical link] (L2CAP)	
PSK	Phase Shift Keying	Class of modulation types
PSTN	Public Switched Telephone Network	

Table 1.1: Acronyms and Abbreviations

Acronym or abbreviation	Writing out in full	Which means
ptt	Packet Type Table	The ptt parameter is used to select the logical transport types via LMP.
Q		
QoS	Quality of Service	
R		
RAND	Random number	
RF	Radio Frequency	
RFC	Request For Comments	
RFCMode	Retransmission and Flow Control Mode	
RFCOMM		Serial cable emulation protocol based on ETSI TS 07.10
RMS	Root Mean Square	
RSSI	Received Signal Strength Indication	
RX	Receive	
S		
S	Slave	
SAP	Service Access Points	
SAR	Segmentation and Reassembly	
SCO	Synchronous Connection-Oriented [logical transport]	Bi-directional, symmetric, point-to-point, AV channels.
SCO-S	Stream SCO (unframed)	
SCO-S	Synchronous logical link	used to support isochronous data delivered in a stream without framing
SD	Service Discovery	
SDP	Service Discovery Protocol	
SDU	Service Data Unit	
SEQN	Sequential Numbering scheme	
SRES	Signed Response	
SS	Supplementary Services	
SSI	Signal Strength Indication	

Table 1.1: Acronyms and Abbreviations

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Acronym or abbreviation	Writing out in full	Which means
SSP	Secure Simple Pairing	Simplifies the pairing process and improves Bluetooth security.
SSR	Sniff Subrating	A mode that defines the anchor points at which the master transmits to the slave.
SW	Software	
T		
TBD	To Be Defined	
TC	Test Control	Test Control layer for the test interface
TCI	Test Control Interface	
TCP/IP	Transport Control Protocol/Internet Protocol	
TCS	Telephony Control protocol Specification	
TDD	Time-Division Duplex	
TX	Transmit	
U		
UAP	Upper Address Part	
UART	Universal Asynchronous receiver Transmitter	
UI	User Interface	
UI	Unnumbered Information	
ULAP	Upper and Lower Address Parts	
USB	Universal Serial Bus	
UUID	Universal Unique Identifier	
W		
WAP	Wireless Application Protocol	

Table 1.1: Acronyms and Abbreviations

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CORE SPECIFICATION CHANGE HISTORY

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1 DEPRECATED FEATURES

Some features have been deemed no longer useful and have been deprecated. The term *deprecation* does not mean that they are no longer allowed, but that they are no longer recommended as the best way of performing a given function.

Features that have been deprecated are not included in this specification or corresponding test documentation. Deprecated features may be implemented using the latest qualifiable version of the Bluetooth core specification that specifies the deprecated features.

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2 CHANGES FROM V1.1 TO V1.2

2.1 NEW FEATURES

Several new features were introduced in the Bluetooth Core Specification 1.2. The major areas of improvement are:

- Architectural overview
- Faster connection
- Adaptive frequency hopping
- Extended SCO links
- Enhanced error detection and flow control
- Enhanced synchronization capability
- Enhanced flow specification

The feature descriptions are incorporated into the existing text in different core parts, described in Volumes 2 and 3.

2.2 STRUCTURE CHANGES

The Bluetooth Core Specification 1.2 was significantly restructured for better consistency and readability. The most important structure changes have been performed in Baseband, LMP, HCI and L2CAP. The text in these sections has been rearranged to provide:

- Presentation of the information in a more logical progression
- Removal of redundant text and requirements
- Consolidation of baseband related requirements (for example, the *Baseband Timers* and *Bluetooth Audio* sections into the Baseband Specification)
- Alignment of the specification with the new architecture and terminology presented in the Architecture Overview (see “IEEE Language” on page 87 [vol. 1]).

In addition, new text and requirements were added for the new features as well as many changes throughout the specification to update the text to use IEEE language (see “IEEE Language” on page 87 [vol. 1]).

2.3 DEPRECATED FEATURES LIST

Features deprecated in version 1.2 are:

- The use of Unit Keys for security
- Optional Paging schemes
- 23 channel hopping sequence

- Page scan period mode and associated commands

2.4 CHANGES IN WORDING

Two general classes of changes to the wording of the Bluetooth Specification have been done for version 1.2. They are a formalization of the language by using conventions established by the Institute of Electrical and Electronic Engineers (IEEE), and a regularization of Bluetooth wireless technology-specific terms. Many portions of the version 1.1 specification use imprecise or inaccurate terms to describe attributes of the protocol. A more accurate terminology described in Part E was introduced into the version 1.2 specification and have been applied in future versions.

2.5 NOMENCLATURE CHANGES

The nomenclature in Bluetooth 1.2 was also updated due to new concepts that are introduced together with the new features and the new architecture. See “Architecture” on page 9.

3 CHANGES FROM V1.2 TO V2.0 + EDR

3.1 NEW FEATURES

The Bluetooth Core Specification version 2.0 + EDR introduces Enhanced Data Rate (EDR). EDR provides a set of additional packet types that use the new 2 Mbps and 3 Mbps modes.

In addition to EDR a set of errata provided in ESR02 has been incorporated into this version and revised to include changes caused by the addition of EDR.

These additions are incorporated into the existing text in different core parts described in Volumes 2 and 3.

3.2 DEPRECATED FEATURES

The only feature deprecated in version 2.0 + EDR is the Page Scan Period Mode and associated commands (based on Erratum 694 which is also included in ESR02).

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4 CHANGES FROM V2.0 + EDR TO V2.1 + EDR

4.1 NEW FEATURES

Several new features are introduced in Bluetooth Core Specification 2.1 + EDR. The major areas of improvement are:

- Erroneous Data Reporting
- Encryption Pause and Resume
- Extended Inquiry Response
- Link Supervision Timeout Changed Event
- Non-Automatically-Flushable Packet Boundary Flag
- Secure Simple Pairing
- Sniff Subrating
- Security Mode 4

4.2 DEPRECATED FEATURES

No features were deprecated in v2.1 + EDR.

MIXING OF SPECIFICATION VERSIONS

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1 MIXING OF SPECIFICATION VERSIONS

This part describes how different versions of the Core System Packages can be mixed in Bluetooth implementations. The Core System Packages consist of a Controller Package (see volume 2) and a Host Package (see volume 3).

In order to describe how these packages can be mixed, one needs to distinguish between four categories of features specified in the different specification versions. The four categories are:

Category	Description
Type 1	Feature that exists below HCI and cannot be configured via HCI
Type 2	Feature that exists below HCI and can be configured/enabled via HCI
Type 3	Feature that exists below and above HCI and requires HCI command/ events to function
Type 4	Feature that exists only above HCI

The outcome of mixing different core system packages are derived from the feature definitions in the table above:

- If an implementation contains features of type 1 or type 4, these features can function with any combination of Host Package and Controller Package versions.
- If an implementation contains features of type 2, these features can only be used under a default condition if a Host Package of an older version is mixed with a Controller Package of this version.
- In order to fully use the feature under all conditions, the Host Package and Controller Package must be of the same version.
- If an implementation contains features of type 3, these features can only function with a Host Package and a Controller Package both in the same version.

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1.1 FEATURES AND THEIR TYPES

The following table lists the features and their types.

Feature	Version	Type
Basic AFH operation	V1.2	1
Enhanced inquiry	V1.2	1
Configuration of AFH (setting channels and enabling/disabling channel assessment)	V1.2	2
Enhanced synchronization capability	V1.2	2
Interlaced inquiry scan	V1.2	2
Interlaced page scan	V1.2	2
Broadcast encryption	V1.2	2
Enhanced flow specification and flush time-out	V1.2	3
Extended SCO links	V1.2	3
Inquiry Result with RSSI	V1.2	3
L2CAP flow and error control	V1.2	4
2 Mbps EDR	V2.0 + EDR	2
3 Mbps EDR	V2.0 + EDR	2
3 slot packets in EDR	V2.0 + EDR	2
5 slot packets in EDR	V2.0 + EDR	2
2 Mbps eSCO	V2.0 + EDR	2*
3 Mbps eSCO	V2.0 + EDR	2*
3 slot packets for EDR eSCO	V2.0 + EDR	2*
Erroneous Data Reporting	V2.1 + EDR	3
Extended Inquiry Response	V2.1 + EDR	3
Encryption Pause and Resume	V2.1 + EDR	1
Link Supervision Timeout Changed Event	V2.1 + EDR	3
Non-Flushable Packet Boundary Flag	V2.1 + EDR	3
Sniff subrating	V2.1+ EDR	3
Secure Simple Pairing	V21.+ EDR	3

The EDR eSCO options are marked as 2* because eSCO requires profile support, but if a product includes the eSCO option from V1.2, then EDR eSCO will be supported without any new support above HCI.

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IEEE LANGUAGE

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1 USE OF IEEE LANGUAGE

One of the purposes of this terminology is to make it easy for the reader to identify text that describes requirements as opposed to background information. The general term for text that describes attributes that are required for proper implementation of Bluetooth wireless technology is normative. The general term for language that provides background and context for normative text is informative. These terms are used in various sections to clarify implementation requirements.

Many portions of the Bluetooth Specification use imprecise or inaccurate terms to describe attributes of the protocol. This subsection describes the correct usage of key terms that indicate degree of requirements for processes and data structures. The information here was derived from the Institute of Electrical and Electronic Engineers (IEEE) Style Guide, see http://standards.ieee.org/guides/style/.

The following list is a summary of the terms to be discussed in more detail below:

shall	is required to – used to define requirements
must	is a natural consequence of -- used only to describe unavoidable situations
will	it is true that -- only used in statements of fact
should	is recommended that – used to indicate that among several possibilities one is recommended as particularly suitable, but not required
may	is permitted to – used to allow options
can	is able to – used to relate statements in a causal fashion
is	is defined as – used to further explain elements that are previously required or allowed
note	<informational text ONLY>

For clarity of the definition of those terms, the following sections document why and how they are used. For these sections only, the IEEE terms are italicized to indicate their use as a noun. Uses and examples of the use of the terms in this section are underlined.

1.1 SHALL

The word *shall* is used to indicate mandatory requirements that shall be followed in order to conform to the specification and from which no deviation is permitted.

There is a strong implication that the presence of the word *shall* indicates a testable requirement. All testable requirements shall be reflected in the Protocol Implementation Conformance Statement (PICS). In turn, all PICS indicators should be reflected in the Test Cases (TCs) either directly or indirectly.

A direct reference is a specific test for the attribute cited in the text. For example, a minimum value for a given parameter may be an entry in the TCs. Indirect test coverage may be appropriate if the existence of the attribute is requisite for passing a higher level test.

1.2 MUST

The word *must* shall not be used when stating mandatory requirements. *Must* is used only to describe unavoidable situations and is seldom appropriate for the text of a Specification.

An example of an appropriate use of the term *must* is: “the Bluetooth radios must be in range of each other to communicate”.

1.3 WILL

The use of the word *will* shall not be used when stating mandatory requirements. The term *will* is only used in statements of fact. As with the term *must*, *will* is not generally applicable to the description of a protocol. An example of appropriate use of will is: “when power is removed from the radio, it can be assumed that communications will fail”

1.4 SHOULD

Should equals *is recommended that*. The word *should* is used to indicate that among several possibilities one is recommended as particularly suitable without mentioning or excluding others. Alternatively it may indicate that a certain course of action is preferred but not necessarily required. Finally, in the negative form, it indicates a certain course of action is deprecated but not prohibited.

In the Bluetooth Specification the term designates an optional attribute that may require an entry in the PICS.

Explicit specification of alternatives should be done when using *should*.

1.5 MAY

The word *may* is used to indicate a course of action permissible within the limits of the specification. The term *may* equals *is permitted*. This is generally used when there is a single, optional attribute described, but multiple alternatives may be cited.

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The use of *may* implies an optional condition in the PICS and therefore may need to be reflected in the corresponding test cases.

1.6 CAN

The word *can* is used for statements of possibility and capability, whether material, physical, or causal. The term *can* equals *is able to*.

The term *can* shall be used only in informative text. It describes capabilities by virtue of the rules established by normative text.

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